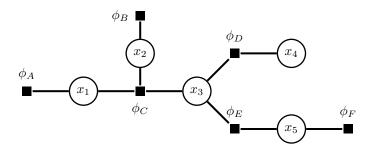


# Exercise 1. Sum-product message passing

We here re-consider the factor tree from the lecture on exact inference.

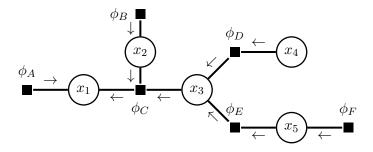


Let all variables be binary,  $x_i \in \{0,1\}$ , and the factors be defined as follows:

	$x_1$	$x_2$	$x_3$	$\phi_C$								
	0 1	0	0	4 2	$\overline{x_3}$	$x_4$	$\phi_D$	$\overline{x_3}$	$x_5$	$\phi_E$		
$x_1  \phi_A \qquad x_2  \phi_B$	0	1	0	2	0	0	8	0	0	3	$x_5$	$\phi_F$
0 2 0 4	1	1	0	6	1	0	2	1	0	6	$\theta$	1
1 4 1 4	0	0	1	2	0	1	2	0	1	6	1	8
	1	0	1	6	1	1	6	1	1	3		
	0	1	1	6								
	1	1	1	4								

(a) Mark the graph with arrows indicating all messages that need to be computed for the computation of  $p(x_1)$ .

#### Solution.



(b) Compute the messages that you have identified.

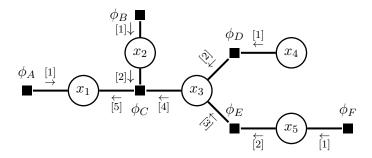
Assuming that the computation of the messages is scheduled according to a common clock, group the messages together so that all messages in the same group can be computed in parallel during a clock cycle.

**Solution.** Since the variables are binary, each message can be represented as a two-dimensional vector. We use the convention that the first element of the vector corresponds to the message for  $x_i = 0$  and the second element to the message for  $x_i = 1$ . For example,

$$\mu_{\phi_A \to x_1} = \begin{pmatrix} 2\\4 \end{pmatrix} \tag{S.1}$$

means that the message  $\mu_{\phi_A \to x_1}(x_1)$  equals 2 for  $x_1 = 0$ , i.e.  $\mu_{\phi_A \to x_1}(0) = 2$ .

The following figure shows a grouping (scheduling) of the computation of the messages.



### Clock cycle 1:

$$\mu_{\phi_{A} \to x_{1}} = \begin{pmatrix} 2 \\ 4 \end{pmatrix}$$
 $\mu_{\phi_{B} \to x_{2}} = \begin{pmatrix} 4 \\ 4 \end{pmatrix}$ 
 $\mu_{x_{4} \to \phi_{D}} = \begin{pmatrix} 1 \\ 1 \end{pmatrix}$ 
 $\mu_{\phi_{F} \to x_{5}} = \begin{pmatrix} 1 \\ 8 \end{pmatrix}$ 
(S.2)

Clock cycle 2:

$$\mu_{x_2 \to \phi_C} = \mu_{\phi_B \to x_2} = \begin{pmatrix} 4 \\ 4 \end{pmatrix}$$

$$\mu_{x_5 \to \phi_E} = \mu_{\phi_F \to x_5} = \begin{pmatrix} 1 \\ 8 \end{pmatrix}$$
(S.3)

Message  $\mu_{\phi_D \to x_3}$  is defined as

$$\mu_{\phi_D \to x_3}(x_3) = \sum_{x_4} \phi_D(x_3, x_4) \mu_{x_4 \to \phi_D}(x_4)$$
 (S.4)

so that

$$\mu_{\phi_D \to x_3}(0) = \sum_{x_4=0}^{1} \phi_D(0, x_4) \mu_{x_4 \to \phi_D}(x_4)$$
(S.5)

$$= \phi_D(0,0)\mu_{x_4 \to \phi_D}(0) + \phi_D(0,1)\mu_{x_4 \to \phi_D}(1)$$
 (S.6)

$$= 8 \cdot 1 + 2 \cdot 1 \tag{S.7}$$

$$= 10 \tag{S.8}$$

$$\mu_{\phi_D \to x_3}(1) = \sum_{x_4=0}^{1} \phi_D(1, x_4) \mu_{x_4 \to \phi_D}(x_4)$$
 (S.9)

$$= \phi_D(1,0)\mu_{x_4 \to \phi_D}(0) + \phi_D(1,1)\mu_{x_4 \to \phi_D}(1)$$
 (S.10)

$$= 2 \cdot 1 + 6 \cdot 1 \tag{S.11}$$

$$= 8 \tag{S.12}$$

and thus

$$\mu_{\phi_D \to x_3} = \begin{pmatrix} 10 \\ 8 \end{pmatrix}. \tag{S.13}$$

The above computations can be written more compactly in matrix notation. Let  $\phi_{\mathbf{D}}$  be the matrix that contains the outputs of  $\phi_D(x_3, x_4)$ 

$$\phi_{\mathbf{D}} = \begin{pmatrix} \phi_D(x_3 = 0, x_4 = 0) & \phi_D(x_3 = 0, x_4 = 1) \\ \phi_D(x_3 = 1, x_4 = 0) & \phi_D(x_3 = 1, x_4 = 1) \end{pmatrix} = \begin{pmatrix} 8 & 2 \\ 2 & 6 \end{pmatrix}.$$
 (S.14)

We can then write  $\mu_{\phi_D \to x_3}$  in terms of a matrix vector product,

$$\mu_{\phi_D \to x_3} = \phi_D \mu_{x_4 \to \phi_D}. \tag{S.15}$$

## Clock cycle 3:

Representing the factor  $\phi_E$  as matrix  $\phi_E$ ,

$$\phi_{E} = \begin{pmatrix} \phi_{E}(x_{3} = 0, x_{5} = 0) & \phi_{E}(x_{3} = 0, x_{5} = 1) \\ \phi_{E}(x_{3} = 1, x_{5} = 0) & \phi_{E}(x_{3} = 1, x_{5} = 1) \end{pmatrix} = \begin{pmatrix} 3 & 6 \\ 6 & 3 \end{pmatrix},$$
 (S.16)

we can write

$$\mu_{\phi_E \to x_3}(x_3) = \sum_{x_5} \phi_E(x_3, x_5) \mu_{x_5 \to \phi_E}(x_5)$$
 (S.17)

as a matrix vector product,

$$\mu_{\phi_E \to x_3} = \phi_E \mu_{x_5 \to \phi_E} \tag{S.18}$$

$$= \begin{pmatrix} 3 & 6 \\ 6 & 3 \end{pmatrix} \begin{pmatrix} 1 \\ 8 \end{pmatrix} \tag{S.19}$$

$$= \begin{pmatrix} 51\\30 \end{pmatrix}. \tag{S.20}$$

#### Clock cycle 4:

Variable node  $x_3$  has received all incoming messages, and can thus output  $\mu_{x_3 \to \phi_C}$ ,

$$\mu_{x_3 \to \phi_C}(x_3) = \mu_{\phi_D \to x_3}(x_3)\mu_{\phi_E \to x_3}(x_3). \tag{S.21}$$

Using  $\odot$  to denote element-wise multiplication of two vectors, we have

$$\mu_{\mathbf{x_3} \to \phi_{\mathbf{C}}} = \mu_{\phi_{\mathbf{D}} \to \mathbf{x_3}} \odot \mu_{\phi_{\mathbf{E}} \to \mathbf{x_3}}$$
 (S.22)

$$= \begin{pmatrix} 10\\8 \end{pmatrix} \odot \begin{pmatrix} 51\\30 \end{pmatrix} \tag{S.23}$$

$$= \begin{pmatrix} 510\\240 \end{pmatrix}. \tag{S.24}$$

#### Clock cycle 5:

Factor node  $\phi_C$  has received all incoming messages, and can thus output  $\mu_{\phi_C \to x_1}$ ,

$$\mu_{\phi_C \to x_1}(x_1) = \sum_{x_2, x_3} \phi_C(x_1, x_2, x_3) \mu_{x_2 \to \phi_C}(x_2) \mu_{x_3 \to \phi_C}(x_3).$$
 (S.25)

Writing out the sum for  $x_1 = 0$  and  $x_1 = 1$  gives

$$\mu_{\phi_C \to x_1}(0) = \sum_{x_2, x_3} \phi_C(0, x_2, x_3) \mu_{x_2 \to \phi_C}(x_2) \mu_{x_3 \to \phi_C}(x_3)$$
 (S.26)

$$= \phi_C(0, x_2, x_3) \mu_{x_2 \to \phi_C}(x_2) \mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (0, 0)} +$$
 (S.27)

$$\phi_C(0, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (1.0)} +$$
 (S.28)

$$\phi_C(0, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (0, 1)} +$$
 (S.29)

$$\phi_C(0, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (1.1)}$$
 (S.30)

$$=4 \cdot 4 \cdot 510+$$
 (S.31)

$$2 \cdot 4 \cdot 510 +$$
 (S.32)

$$2 \cdot 4 \cdot 240 +$$
 (S.33)

$$6 \cdot 4 \cdot 240 \tag{S.34}$$

$$=19920$$
 (S.35)

$$\mu_{\phi_C \to x_1}(1) = \sum_{x_2, x_3} \phi_C(1, x_2, x_3) \mu_{x_2 \to \phi_C}(x_2) \mu_{x_3 \to \phi_C}(x_3)$$
 (S.36)

$$= \phi_C(1, x_2, x_3) \mu_{x_2 \to \phi_C}(x_2) \mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (0, 0)} +$$
 (S.37)

$$\phi_C(1, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3)|_{(x_2, x_3) = (1.0)} +$$
 (S.38)

$$\phi_C(1, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (0, 1)} +$$
 (S.39)

$$\phi_C(1, x_2, x_3)\mu_{x_2 \to \phi_C}(x_2)\mu_{x_3 \to \phi_C}(x_3) \mid_{(x_2, x_3) = (1, 1)}$$
 (S.40)

$$=2 \cdot 4 \cdot 510 +$$
 (S.41)

$$6 \cdot 4 \cdot 510 + \tag{S.42}$$

$$6 \cdot 4 \cdot 240 + \tag{S.43}$$

$$4 \cdot 4 \cdot 240 \tag{S.44}$$

$$=25920$$
 (S.45)

and hence

$$\mu_{\phi_C \to x_1} = \begin{pmatrix} 19920 \\ 25920 \end{pmatrix} \tag{S.46}$$

After step 5, variable node  $x_1$  has received all incoming messages and the marginal can be computed.

In addition to the messages needed for computation of  $p(x_1)$  one can compute *all* messages in the graph in five clock cycles, see Figure 1. This means that *all* marginals, as well as the joints of those variables sharing a factor node, are available after five clock cycles.

(c) What is  $p(x_1 = 1)$ ?

**Solution.** We compute the marginal  $p(x_1)$  as

$$p(x_1) \propto \mu_{\phi_A \to x_1}(x_1) \mu_{\phi_C \to x_1}(x_1)$$
 (S.47)

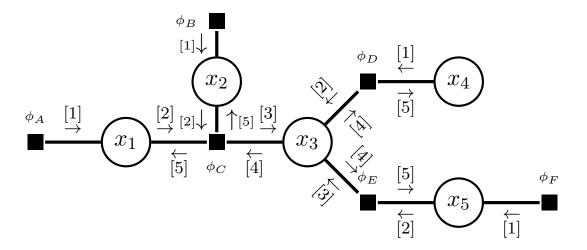


Figure 1: Answer to Exercise 1 Question (b): Computing all messages in five clock cycles. If we also computed the messages toward the leaf factor nodes, we needed six cycles, but they are not necessary for computation of the marginals so they are omitted.

which is in vector notation

$$\begin{pmatrix}
p(x_1 = 0) \\
p(x_1 = 1)
\end{pmatrix} \propto \boldsymbol{\mu}_{\boldsymbol{\phi}_{\boldsymbol{A}} \to \boldsymbol{x}_1} \odot \boldsymbol{\mu}_{\boldsymbol{\phi}_{\boldsymbol{C}} \to \boldsymbol{x}_1} \tag{S.48}$$

$$\propto \binom{2}{4} \odot \binom{19920}{25920} \tag{S.49}$$

$$\propto \begin{pmatrix} 39840\\103680 \end{pmatrix}. \tag{S.50}$$

Normalisation gives

$$\begin{pmatrix} p(x_1=0) \\ p(x_1=1) \end{pmatrix} = \frac{1}{39840 + 103680} \begin{pmatrix} 39840 \\ 103680 \end{pmatrix}$$
(S.51)

$$= \begin{pmatrix} 0.2776 \\ 0.7224 \end{pmatrix} \tag{S.52}$$

so that  $p(x_1 = 1) = 0.7224$ .

(d) Draw the factor graph corresponding to  $p(x_1, x_3, x_4, x_5|x_2 = 1)$  and provide the numerical values for all factors.

The pmf represented by the original factor graph is

$$p(x_1,\ldots,x_5) \propto \phi_A(x_1)\phi_B(x_2)\phi_C(x_1,x_2,x_3)\phi_D(x_3,x_4)\phi_E(x_3,x_5)\phi_F(x_5)$$

The conditional  $p(x_1, x_3, x_4, x_5 | x_2 = 1)$  is proportional to  $p(x_1, \ldots, x_5)$  with  $x_2$  fixed to  $x_2 = 1$ , i.e.

$$p(x_1, x_3, x_4, x_5 | x_2 = 1) \propto p(x_1, x_2 = 1, x_3, x_4, x_5)$$
 (S.53)

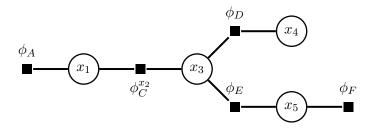
$$\propto \phi_A(x_1)\phi_B(x_2=1)\phi_C(x_1,x_2=1,x_3)\phi_D(x_3,x_4)\phi_E(x_3,x_5)\phi_F(x_5)$$
(S.54)

$$\propto \phi_A(x_1)\phi_C^{x_2}(x_1, x_3)\phi_D(x_3, x_4)\phi_E(x_3, x_5)\phi_F(x_5) \tag{S.55}$$

where  $\phi_C^{x_2}(x_1, x_3) = \phi_C(x_1, x_2 = 1, x_3)$ . The numerical values of  $\phi_C^{x_2}(x_1, x_3)$  can be read from the table defining  $\phi_C(x_1, x_2, x_3)$ , extracting those rows where  $x_2 = 1$ ,

	$x_1$	$x_2$	$x_3$	$\phi_C$
	0	0	0	4
$\rightarrow$	$\frac{1}{0}$	$0 \\ 1$	$0 \\ 0$	$\frac{2}{2}$
$\rightarrow$	1	1	0	6
	0	0	1	2
	1	0	1	6
$\rightarrow$	0	1	1	6

The factor graph for  $p(x_1, x_3, x_4, x_5 | x_2 = 1)$  is shown below. Factor  $\phi_B$  has disappeared since it only depended on  $x_2$  and thus became a constant. Factor  $\phi_C$  is replaced by  $\phi_C^{x_2}$  defined above. The remaining factors are the same as in the original factor graph.



(e) Compute  $p(x_1 = 1|x_2 = 1)$ , re-using messages that you have already computed for the evaluation of  $p(x_1 = 1)$ .

**Solution.** The message  $\mu_{\phi_A \to x_1}$  is the same as in the original factor graph and  $\mu_{x_3 \to \phi_C^{x_2}} = \mu_{x_3 \to \phi_C}$ . This is because the outgoing message from  $x_3$  corresponds to the effective factor obtained by summing out all variables in the sub-trees attached to  $x_3$  (without the  $\phi_C^{x_2}$  branch), and these sub-trees do not depend on  $x_2$ .

The message  $\mu_{\phi_C^{x_2} \to x_1}$  needs to be newly computed. We have

$$\mu_{\phi_C^{x_2} \to x_1}(x_1) = \sum_{x_3} \phi_C^{x_2}(x_1, x_3) \mu_{x_3 \to \phi_C^{x_2}}$$
 (S.56)

or in vector notation

$$\mu_{\phi_C^{x_2} \to x_1} = \phi_C^{x_2} \mu_{x_3 \to \phi_C^{x_2}} \tag{S.57}$$

$$= \begin{pmatrix} \phi_C^{x_2}(x_1 = 0, x_3 = 0) & \phi_C^{x_2}(x_1 = 0, x_3 = 1) \\ \phi_C^{x_2}(x_1 = 1, x_3 = 0) & \phi_C^{x_2}(x_1 = 1, x_3 = 1) \end{pmatrix} \boldsymbol{\mu}_{\boldsymbol{x_3} \to \boldsymbol{\phi}_{\boldsymbol{C}}^{\boldsymbol{x_2}}}$$
 (S.58)

$$= \begin{pmatrix} 2 & 6 \\ 6 & 4 \end{pmatrix} \begin{pmatrix} 510 \\ 240 \end{pmatrix} \tag{S.59}$$

$$= \begin{pmatrix} 2460\\4020 \end{pmatrix} \tag{S.60}$$

We thus obtain for the marginal posterior of  $x_1$  given  $x_2 = 1$ :

$$\begin{pmatrix}
p(x_1 = 0 | x_2 = 1) \\
p(x_1 = 1 | x_2 = 1)
\end{pmatrix} \propto \boldsymbol{\mu_{\phi_A \to x_1}} \odot \boldsymbol{\mu_{\phi_C^{x_2} \to x_1}}$$
(S.61)

$$\propto \begin{pmatrix} 2\\4 \end{pmatrix} \odot \begin{pmatrix} 2460\\4020 \end{pmatrix} \tag{S.62}$$

$$\propto \begin{pmatrix} 4920\\16080 \end{pmatrix}. \tag{S.63}$$

Normalisation gives

$$\begin{pmatrix}
p(x_1 = 0 | x_2 = 1) \\
p(x_1 = 1 | x_2 = 1)
\end{pmatrix} = \begin{pmatrix}
0.2343 \\
0.7657
\end{pmatrix}$$
(S.64)

and thus  $p(x_1 = 1|x_2 = 1) = 0.7657$ . The posterior probability is slightly larger than the prior probability,  $p(x_1 = 1) = 0.7224$ .