Phrase Structure and Parsing as Search
Informatics 2A: Lecture 17

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Part-of-speech tagging and its applications

The use of hidden Markov models for POS tagging

The Viterbi algorithm

Heads and Phrases

> Noun (N): Noun Phrase (NP)
> Adjective (A): Adjective Phrase (AP)
> Verb (V): Verb Phrase (VP)
> Preposition (P): Prepositional Phrase (PP)

- So far we have looked at terminals (words or POS tags).
- Today, we'll look at non-terminals, which correspond to phrases.
- The class that a word belongs to is closely linked to the name of the phrase it customarily appears in.
- In a X-phrase (eg NP), the key occurrence of X (eg N) is called the head.
- In English, the head tends to appear in the middle of a phrase.

Heads and Phrases

English NPs are commonly of the form:

(Det) Adj* **Noun** (PP | RelClause)*
NP: *the angry duck that tried to bite me*, head: *duck*.

VPs are commonly of the form:

(Aux) Adv* **Verb** Arg* Adjunct*
Arg → NP | PP
Adjunct → PP | AdvP | . . .
VP: *usually eats artichokes for dinner*, head: *eat*.

In Japanese, Korean, Hindi, Urdu, and other head-final languages, the head is at the end of its associated phrase.

In Irish, Welsh, Scots Gaelic and other head-initial languages, the head is at the beginning of its associated phrase.

Desirable Properties of a Grammar

Chomsky specified two properties that make a grammar "interesting and satisfying":

- It should be a finite specification of the strings of the language, rather than a list of its sentences.
- It should be revealing, in allowing strings to be associated with meaning (semantics) in a systematic way.

We can add another desirable property:

- It should capture structural and distributional properties of the language. (E.g. where heads of phrases are located; how a sentence transforms into a question; which phrases can float around the sentence.)

Desirable Properties of a Grammar

- Context-free grammars (CFGs) provide a pretty good approximation.
- Some features of NLs are more easily captured using mildly context-sensitive grammars, as well see later in the course.
- There are also more modern grammar formalisms that better capture structural and distributional properties of human languages. (E.g. combinatory categorial grammar.)
- But LL(1) grammars and the like definitely aren't enough for NLs. Even if we could make a NL grammar LL(1), we wouldn't want to: this would artificially suppress ambiguities, and would often mutilate the 'natural' structure of sentences.

Let's say we want to capture in a grammar the structural and distributional properties that give rise to sentences like:

A duck walked in the park.	NP,V,PP
The man walked with a duck.	NP,V,PP
You made a duck.	Pro,V,NP
You made her duck.	? Pro,V,NP
A man with a telescope saw you.	NP,PP,V,Pro
A man saw you with a telescope.	NP,V,Pro,PP
You saw a man with a telescope.	Pro,V,NP,PP

We want to write grammatical rules that generate these phrase structures, and lexical rules that generate the words appearing in them.

Grammar G1 generates the sentences on the previous slide:

Grammatical rules	Lexical rules
S → NP VP	Det → a | the | her (determiners)
NP → Det N	N → man | park | duck | telescope (nouns)
NP → Det N PP	Pro → you (pronoun)
NP → Pro	V → saw | walked | made (verbs)
VP → V NP PP	Prep → in | with | for (prepositions)
VP → V NP	
VP → V	
PP → Prep NP	

Does G1 produce a finite or an infinite number of sentences?

Recursion

Recursion in a grammar makes it possible to generate an infinite number of sentences.

In direct recursion, a non-terminal on the LHS of a rule also appears on its RHS. The following rules add direct recursion to G1:
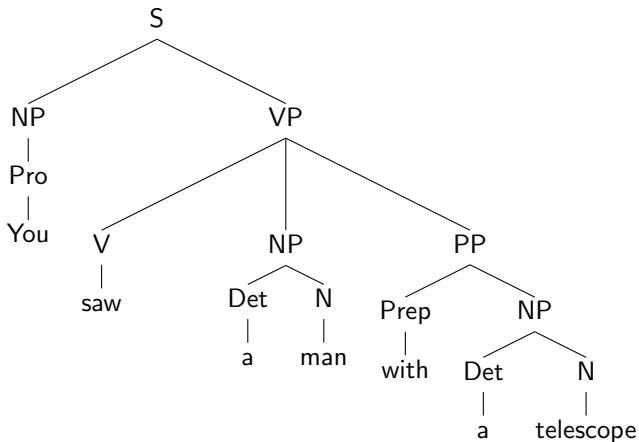
> VP → VP Conj VP
> Conj → and | or

In indirect recursion, some non-terminal can be expanded (via several steps) to a sequence of symbols containing that non-terminal:
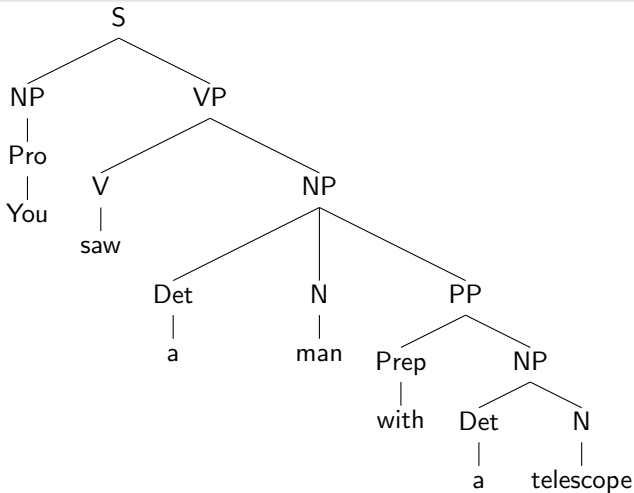
> NP → Det N PP
> PP → Prep NP

You saw a man with a telescope.

Structural Ambiguity

You saw a man with a telescope.

```
                    S
         ┌──────────┴──────────┐
        NP                     VP
         │              ┌───────┴────────┐
        Pro             V                NP
         │              │        ┌───────┼────────┐
        You            saw      Det      N        PP
                                 │       │     ┌───┴────┐
                                 a      man   Prep      NP
                                             with   ┌────┴────┐
                                                   Det        N
                                                    │         │
                                                    a     telescope
```

You saw a man with a telescope.
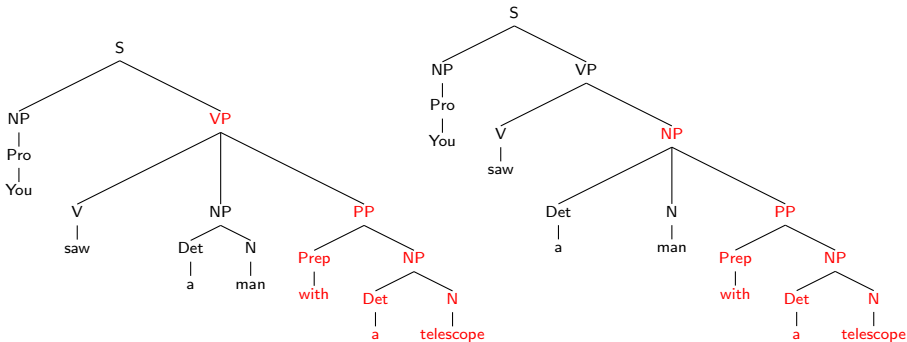
This illustrates attachment ambiguity: the PP can be a part of the VP or of the NP. Note that there's no POS ambiguity here.

Structural Ambiguity

Grammar G1 only gives us one analysis of *you made her duck*.
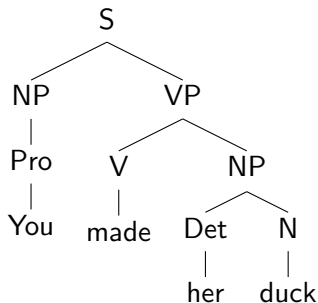
There is another, ditransitive (i.e., two-object) analysis of this sentence – one that underlies the pair:
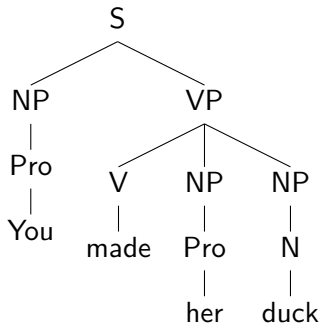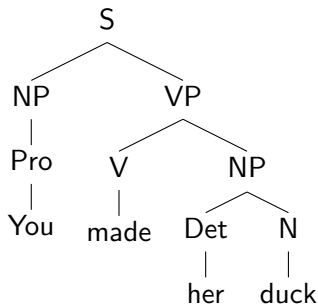
What did you make for her?
You made her duck.

For this alternative, G1 also needs rules like:

NP → N
VP → V NP NP
Pro → her

In this case, the structural ambiguity is rooted in POS ambiguity.

Structural Ambiguity

There is a third analysis as well, one that underlies the pair:

> What did you make her do?
> You made her duck. (move head or body quickly downwards)

Here, the small clause (*her duck*) is the direct object of a verb.

Similar small clauses are possible with verbs like *see*, *hear* and *notice*, but not *ask*, *want*, *persuade*, etc.

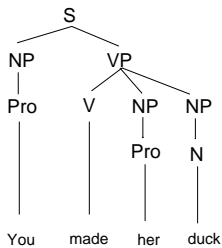G1 needs a rule that requires accusative case-marking on the subject of a small clause and no tense on its verb.:
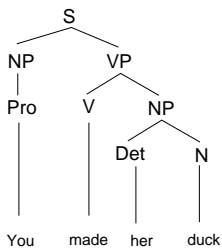
VP → V S1
S1 → NP(acc) VP(untensed)
NP(acc) → her | him | them

Now we have three analyses for *you made her duck:*
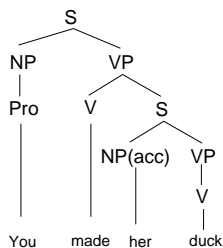
How can we compute these analyses automatically?

A parser is an algorithm that computes a structure for an input string given a grammar. All parsers have two fundamental properties:

- Directionality: the sequence in which the structures are constructed (e.g., top-down or bottom-up).
- Search strategy: the order in which the search space of possible analyses is explored (e.g., depth-first, breadth-first).

For instance, LL(1) parsing is top-down and depth-first.

As we've noted, LL(1) isn't good enough for NL. We'll be looking at other parsing algorithms that work for more general CFGs.

- Recursive descent parsers (top-down). Simple and very general, but inefficient. Other problems
- Shift-reduce parsers (bottom-up).
- The Cocke-Younger-Kasami algorithm (bottom up). Works for any CFG with reasonable efficiency.
- The Earley algorithm (top down). Chart parsing enhanced with prediction.

Recursive Descent Parsing

A recursive descent parser treats a grammar as a specification of how to break down a top-level goal into subgoals. Therefore:

- Parser searches through the trees licensed by the grammar to find the one that has the required sentence along its yield.
- **Directionality** = top-down: It starts from the start symbol of the grammar, and works its way down to the terminals.
- **Search strategy** = depth-first: It expands a given terminal as far as possible before proceeding to the next one.

Algorithm Sketch: Recursive Descent Parsing

1. The top-level goal is to derive the start symbol (S).

2. Choose a grammatical rule with S as its LHS
   (e.g, S $\rightarrow$ NP VP), and replace S with the RHS of the rule
   (the subgoals; e.g., NP and VP).

3. Choose a rule with the leftmost subgoal as its LHS (e.g.,
   NP $\rightarrow$ Det N). Replace the subgoal with the RHS of the rule.

4. Whenever you reach a lexical rule (e.g., Det $\rightarrow$ the), match
   its RHS against the current position in the input string.
   - If it matches, move on to next position in the input.
   - If it doesn't, try next lexical rule with the same LHS.
   - If no rules with same LHS, backtrack to most recent choice of
     grammatical rule and choose another rule with the same LHS.
   - If no more grammatical rules, back up to the previous subgoal.

5. Iterate until the whole input string is consumed, or you fail to
   match one of the positions in the input. Backtrack on failure.
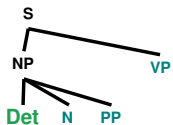
S

the   dog   saw   a   man   in   the   park

S

NP **VP**

the   dog   saw   a   man   in   the   park

the   dog   saw   a   man   in   the   park

S
NP   VP
Det   N   PP
the

. . . .
the   dog   saw   a   man   in   the   park

```
S
├── NP
│    ├── Det
│    ├── N
│    └── PP
└── VP

the
the   dog   saw   a   man   in   the   park
```

```
        S
       /  \
      NP    VP
     / \
   Det  N  PP
         |
        man



   the
.....................................................
   the  dog  saw  a  man  in  the  park
```

S
NP   VP
Det   N   PP

the   dog

the   dog   saw   a   man   in   the   park

S
NP   VP
Det   N   PP
P   NP

the   dog

the   dog   saw   a   man   in   the   park

```
        S
       ╱╲
     NP    VP
     ╱╲
  Det  N  PP
            ╱╲
           P   NP
           │
          in


  the  dog
..............................................
  the   dog   saw   a   man   in   the   park
```

S

NP   VP

Det   N

the

the   dog   saw   a   man   in   the   park

S

NP VP

Det N

the dog

the dog saw a man in the park

```
S
├── NP
│   ├── Det ── the
│   └── N ── dog
└── VP
    ├── V ── saw
    ├── NP
    └── PP
```

the   dog   saw   a   man   in   the   park

```
        S
   NP        VP
 Det   N    V    NP        PP
                Det  N  PP

 the   dog  saw  a
. . . . . . . . . . . . . . . . . . . . . . . . . . . . . . .
 the   dog   saw   a   man   in   the   park
```

S
NP   VP
Det   N   V   NP   PP
Det   N   PP
P   NP
Det   N   PP

the   dog   saw   a   man   in
the   dog   saw   a   man   in   the   park

S
NP   VP
Det   N   V   NP   PP
              Det N

the   dog   saw

the   dog   saw   a   man   in   the   park

Recursive Descent Parsing

```
                    S
          NP              VP
      Det      N      V    NP        PP
                          Det    N


      the   dog   saw   a   man
. . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . . .
      the   dog   saw   a   man   in   the   park
```
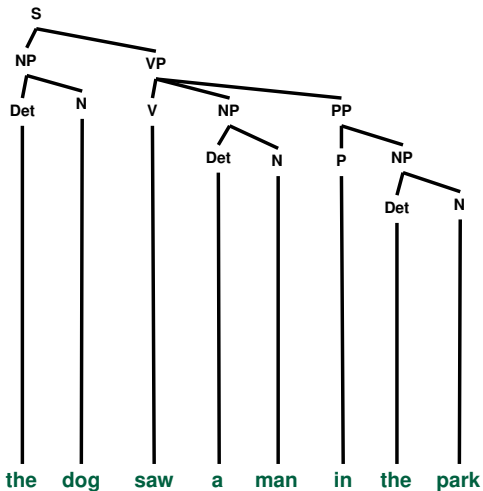
the dog saw a man in the park

A Shift-Reduce parser tries to find sequences of words and phrases that correspond to the righthand side of a grammar production and replace them with the lefthand side:

- **Directionality** = bottom-up: starts with the words of the input and tries to build trees from the words up.
- **Search strategy** = breadth-first: starts with the words, then applies rules with matching right hand sides, and so on until the whole sentence is reduced to an S.

Algorithm Sketch: Shift-Reduce Parsing

Until the words in the sentences are substituted with S:

- Scan through the input until we recognise something that corresponds to the RHS of one of the production rules (shift)

- Apply a production rule in reverse; i.e., replace the RHS of the rule which appears in the sentential form with the LHS of the rule (reduce)

A shift-reduce parser implemented using a stack:

1. start with an empty stack
2. a shift action pushes the current input symbol onto the stack
3. a reduce action replaces *n* items with a single item

Shift-Reduce Parsing

Stack	Remaining
	my dog saw a man in the park
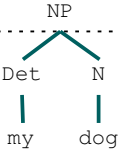
Shift-Reduce Parsing

Stack	Remaining
Det	dog saw a man in the park
my	

Shift-Reduce Parsing
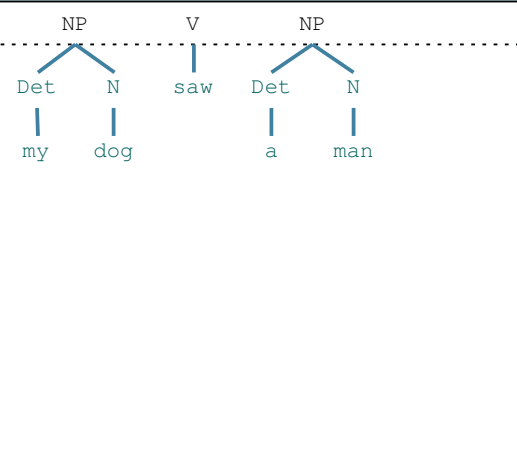
Stack	Remaining
Det N	saw a man in the park
my dog	

Shift-Reduce Parsing

Stack	Remaining
NP	saw a man in the park

```
         NP
        /  \
      Det    N
       |     |
      my    dog
```

Stack	Remaining
NP V NP	in the park

```
    NP        V       NP
   /  \       |      /  \
 Det   N     saw   Det   N
  |    |            |    |
  my  dog           a   man
```

Stack	Remaining

Shift-Reduce Parsing

Stack	


```
      NP           V              NP
     /  \          |           /      \
  Det    N       saw        NP          PP
   |     |                 /  \        /    \
  my    dog             Det    N      P      NP
                         |     |      |     /  \
                         a    man    in   Det    N
                                           |     |
                                          the   park
```
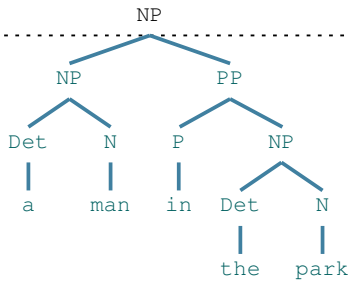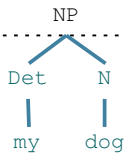
Stack	

Stack	
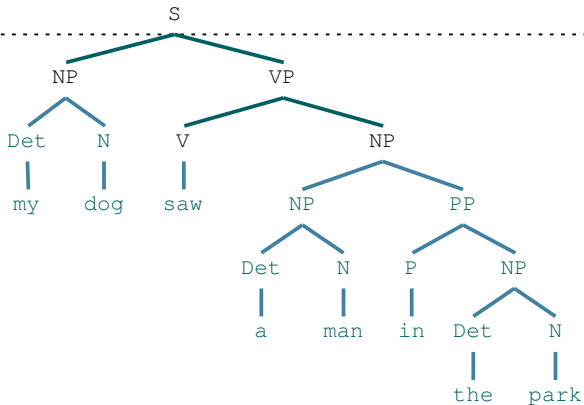
Shift-reduce parsers and pushdown automata

Shift-reduce parsing is equivalent to a pushdown automaton constructed from the CFG (with one state q_0):

- start with empty stack
- shift: a transition in the PDA from q_0 (to q_0) putting a terminal symbol on the stack
- reduce: whenever the righthand side of a rule appears on top of the stack, pop the RHS and push the lefthand side (still staying in q_0). Don't consume anything from the input.
- accept the string if the start symbol is in the stack and the end of string has been reached

If there is some derivation for a given sentence under the CFG, there will be a sequence of actions for which this NPDA accepts the string

If there is some derivation for a given sentence under the CFG, there will be a sequence of actions for which this NPDA accepts the string

But how do we find this derivation?

One way to do this is using so-called generalised LR parsing, which explores all possible paths of the above NPDA

Modern parsers do it differently, because GLR can be expontential in the worst-case

Shift-reduce parsers are highly efficient, they are linear in the length of the string, if they explore only one path

How to do that? **Learn from data** what actions to take at each point, and try to make the optimal decisions so that the correct parser will be found

This keeps the parser linear in the length of the string, but one small error can propagate through the whole parse, and lead to the wrong parse tree

Try it out Yourselves!

Recursive Descent Parser

```
>>> from nltk.app import rdparser
>>> rdparser()
```

Shift-Reduce Parser

```
>>> from nltk.app import srparser
>>> srparser()
```

Summary

- We use CFGs to represent NL grammars
- Grammars need recursion to produce infinite sentences
- Most NL grammars have structural ambiguity
- A parser computes structure for an input automatically
- Recursive descent and shift-reduce parsing
- We'll examine more parsers in Lectures 17–22

Reading: J&M (2nd edition) Chapter 12 (intro – section 12.3), Chapter 13 (intro – section 13.3)

Next lecture: The CYK algorithm