Informatics 1 Functional Programming Lectures 13 and 14

Type Classes

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Part I

Type classes

Element

elem :: Eq a => a -> [a] -> Bool -- comprehension elem x ys = or [x == y | y <- ys] -- recursion elem x [] = False elem x (y:ys) = x == y || elem x ys -- higher-order elem x ys = foldr (||) False (map (x ==) ys)

You've seen types like the one for elem, beginning with Eq $a \Rightarrow ...$ resp. Ord $a \Rightarrow ...$ These express the requirement that a is a type whose values can be tested for equality resp. order (<).

Here are 3 ways of writing elem. No matter how you defined it, you need to use ==. That's where the requirement Eq a comes from.

Using element

+Main> alom 1 [2 3 /]

*MaIII/	
False	elem works for Int
*Main>	elem 'o' "word"
Τηιο	alore works for Char
IIUC	elem works for Char
∗Main> True	elem (1,'o') [(0,'w'),(1,'o'),(2,'r'),(3,'d')] elem works for (Int,Char)

True elem works for String = [Char]

***Main>** elem $(\langle x - \rangle x) [(\langle x - \rangle - x), (\langle x - \rangle - (-x))]$ No instance for (Eq (a -> a)) arising from a use of 'elem' Possible fix: add an instance declaration for (Eq (a -> a))

but elem doesn't work for functions

Testing equality of two functions f,g :: Int -> Int would require testing f x == g x for every possible x :: Int. That would take forever. So Haskell refuses to try.

The same goes for any type INVOLVING functions, for instance (Int->Int,Bool).

*Main> elem "word" ["list", "of", "word"]

The error message invites you to define equality for this type yourself - see below for how to do that.

Equality type class

X == V

class Eq a where by an
 (==) :: a -> a -> Bool

Here's how you could define the TYPE CLASS Eq if it wasn't built in. The definition gives one or more functions that need to be provided by any instance of that class.

instance	Eq Int	where
(==) =	eqInt	Then you can declare that a type is an INSTANCE of the type class by saying what that function / those functions are for that type.
instance	Eq Char	where

= ord x == ord y

instance (Eq a, Eq b) => Eq (a,b) where
 (u,v) == (x,v) = (u == x) && (v == v)

instance Eq a => Eq [a] where
[] == [] = True
[] == y:ys = False
x:xs == [] = False
x:xs == y:ys = (x == y) && (xs == ys)

The definitions of the required functions can be as complicated as you like.

Element, translation

```
data EqDict a = EqD (a -> a -> Bool)
eq :: EqDict a -> a -> a -> Bool
eq (EqDict f) = f
elem :: EqD a \rightarrow a \rightarrow [a] \rightarrow Bool
-- comprehension
elem d x ys = or [eq d x y | y < -ys]
-- recursion
elem d x [] = False
elem d x (y:ys) = eq d x y || elem x ys
-- higher-order
elem d x ys = foldr (||) False (map (eq d x) ys)
```

You can define Haskell with type classes by giving a translation into Haskell without type classes. EqDict a is an equality DICTIONARY - an equality function packaged up into a new type. (In general, a dictionary will package up several functions.) eq extracts the equality function from an equality dictionary. We can then define elem with an extra argument d, which tells it how to compute equality on a. Instead of x==y, we write eq d x y

Type classes, translation

dInt	::	EqDict Int			
dInt	=	EqD eqInt			
dChar	::	EqDict Char			
dChar	=	EqD f			
where					
f x y	=	eq dInt (ord x) (ord y)			
dPair	::	(EqDict a, EqDict b) -> EqDict (a,b)			
dPair (da,db)	=	EqD f			
where					
f (u,v) (x,y)	=	eq da u x && eq db v y			
dList	::	EqDict a -> EqDict [a]			
dList d	=	EqD f			
where					
f [] []	=	True			
f [] (y:ys)	=	False			
f (x:xs) []	=	False			
f (x:xs) (y:ys)	=	eq d x y && eq (dList d) xs ys			

We build up dictionaries, sometimes using other dictionaries. Each INSTANCE declaration creates a dictionary.

Using element, translation

*Main> elem dInt 1 [2,3,4] False

*Main> elem dChar 'o' "word"
True

***Main>** elem (dPair dInt dChar) (1,'o') [(0,'w'),(1,'o')] True

*Main> elem (dList dChar) "word" ["list","of","word"]
True

Haskell uses types to write code for you!

Uses of elem then require the appropriate dictionary as an explicit argument. But Haskell does all of this automatically, using the types that it can infer. You don't need to do it yourself and you don't have an opportunity to get it wrong. Part II

Eq, Ord, Show

Eq, Ord, Show

Eq, Ord and Show are built-in type classes.

class Eq a where (==) :: a -> a -> Bool (/=) :: a -> a -> Bool
Eq actually has two functions, == and /= -- minimum definition: (==) x /= y = not (x == y) You can define a default for some functions in terms of others but instances can override the default. class (Eq a) => Ord a where (<) :: a -> a -> Bool (<=) :: a -> a -> Bool (<=) :: a -> a -> Bool (>) :: a -> a -> Bool
Ord EXTENDS Eq Notice that the default definition of < requires equality.

-- minimum definition: (<=) x < y = x <= y && x /= y x > y = y < x x >= y = y <= x

(>=) :: a -> a -> Bool

class Show a where
 show :: a -> String

Show: need a way of converting a value to a String.

Part III

Booleans, Tuples, Lists

Instances for booleans

instance Eq Bool where

False	==	False	=	True
False	==	True	=	False
True	==	False	=	False
True	==	True	=	True

instance Ord Bool where False <= False = True False <= True = True True <= False = False True <= True = True</pre>

instance Show Bool where show False = "False" show True = "True"

Here's how instances of Eq, Ord and Show can be defined for Bool.

Instances for pairs

instance (Eq a, Eq b) => Eq (a,b) where
 (x,y) == (x',y') = x == x' && y == y'

instance (Ord a, Ord b) => Ord (a,b) where
 (x,y) <= (x',y') = x < x' || (x == x' && y <= y')</pre>

instance (Show a, Show b) => Show (a,b) where
 show (x,y) = "(" ++ show x ++ "," ++ show y ++ ")"

Here's how instances of Eq, Ord and Show can be defined for pairs, using Eq, Ord and Show for each component type.

Instances for lists

instance Eq a => Eq [a] where [] == [] = True [] == y:ys = False x:xs == [] = False x:xs == y:ys = x == y && xs == ysinstance Ord a => Ord [a] where [] <= vs = True x:xs <= [] = False $x:xs \leq y:ys = x \leq y \mid | (x = y \& xs \leq ys)$ instance Show a => Show [a] where show [] = "[]" show (x:xs) = "[" ++ showSep x xs ++ "]" where

showSep x [] = show x showSep x (y:ys) = show x ++ "," ++ showSep y ys

List is similar. We've seen equality already.

Order is an extension of the order on pairs: called dictionary ordering or LEXICOGRAPHIC ORDERING.

Deriving clauses

```
data Bool = False | True
   deriving (Eq, Ord, Show)

data Pair a b = MkPair a b
   deriving (Eq, Ord, Show)

data List a = Nil | Cons a (List a)
```

```
deriving (Eq, Ord, Show)
```

Haskell uses types to write code for you!

You can get definitions of instances of Eq, Ord and Show for free for algebraic types.

Part IV

Sets, revisited

Sets, revisited

instance Ord a => Eq (Set a) where
s == t = s 'equal' t

Note that this differs from the derived instance!

Here's how we can make Set a an instance of Eq.

This refers to the equality function that we defined on the underlying representation of sets.

The one that Haskell would give you for free is different (except for sets represented as ordered lists).

Part V

Numbers

Numerical classes

```
class (Eq a, Show a) => Num a where
  (+), (-), (*) :: a -> a -> a
 negate :: a -> a
 fromInteger :: Integer -> a
 -- minimum definition: (+), (-), (*), fromInteger
 negate x = fromInteger 0 - x
class (Num a) => Fractional a where
 (/) :: a -> a -> a
 recip :: a -> a
 fromRational :: Rational -> a
 -- minimum definition: (/), fromRational
 recip x = 1/x
class (Num a, Ord a) => Real a where
 toRational :: a -> Rational
class (Real a, Enum a) => Integral a where
 div, mod :: a \rightarrow a \rightarrow a
 toInteger :: a -> Integer
```

\There are several type classes for different kinds of numbers. Here's a simplified version of some of them.

A built-in numerical type

instance Num Float where

(+)	=	builtInAddFloat
(-)	=	builtInSubtractFloat
(*)	=	builtInMultiplyFloat
negate	=	builtInNegateFloat
fromInteger	=	builtInFromIntegerFloat

instance Fractional Float where

(/)	=	builtInDivideFloat
fromRational	=	builtInFromRationalFloat

Natural.hs (1)

module Natural(Nat) where
import Test.QuickCheck

data Nat = MkNat Integer

We can also define our own numerical types. Natural numbers are integers that are >= 0.

Remember, we introduce a constructor that is not exported in order to protect the abstraction.

invariant :: Nat \rightarrow Bool invariant (MkNat x) = x >= 0

instance Eq Nat where
 MkNat x == MkNat y = x == y

instance Ord Nat where
 MkNat x <= MkNat y = x <= y</pre>

instance Show Nat where
 show (MkNat x) = show x

Natural.hs (2)

```
instance Num Nat where
MkNat x + MkNat y = MkNat (x + y)
MkNat x - MkNat y
| x >= y = MkNat (x - y)
| otherwise = error (show (x-y) ++ " is negative")
MkNat x * MkNat y = MkNat (x * y)
fromInteger x
| x >= 0 = MkNat x
| otherwise = error (show x ++ " is negative")
negate = undefined
```

Now we can declare Nat as an instance of Num.

We need these operations to PRESERVE THE INVARIANT: if x, y satisfy the invariant, so should x+y etc.

Natural.hs (3)

```
prop_plus :: Integer -> Integer -> Property
prop_plus m n =
   (m >= 0) && (n >= 0) ==> (m+n >= 0)

prop_times :: Integer -> Integer -> Property
prop_times m n =
   (m >= 0) && (n >= 0) ==> (m*n >= 0)

prop_minus :: Integer -> Integer -> Property
prop_minus m n =
   (m >= 0) && (n >= 0) && (m >= n) ==> (m-n >= 0)
```

Here are QuickCheck properties for checking that the invariant is preserved. The invariant isn't preserved if Nat is represented using Int (computer integers) because adding big numbers can give a negative result, but it is preserved if they are represented using Integer (infinite-precision integers).

NaturalTest.hs

module NaturalTest where
import Natural

- m, n :: Nat
- m = fromInteger 2
- n = fromInteger 3

Test run

```
ghci NaturalTest
Ok, modules loaded: NaturalTest, Natural.
*NaturalTest> m
2
*NaturalTest> n
3
*NaturalTest> m+n
5
*NaturalTest> n-m
1
*NaturalTest> m-n
*** Exception: -1 is negative
*NaturalTest> m*n
6
*NaturalTest> fromInteger (-5) :: Nat
*** Exception: -5 is negative
*NaturalTest> MkNat (-5)
Not in scope: data constructor 'MkNat'
```

Hiding—the secret of abstraction

```
module Natural(Nat) where ...
```

```
> ghci NaturalTest
*NaturalTest> let m = fromInteger 2
*NaturalTest> let s = fromInteger (-5)
*** Exception: -5 is negative
*NaturalTest> let s = MkNat (-5)
Not in scope: data constructor `MkNat'
```

VS.

module NaturalUnabs(Nat(MkNat)) where ...

```
> ghci NaturalUnabs
*NaturalUnabs> let p = MkNat (-5) -- breaks invariant
*NaturalUnabs> invariant p
False
```

If I export Nat and not MkNat, I can't break the abstraction.

If you check that all of the functions preserve the invariant, then all values are guaranteed to satisfy it.

Part VI

Seasons

Seasons

data Season = Winter | Spring | Summer | Fall

```
next :: Season -> Season
next Winter = Spring
next Spring = Summer
next Summer = Fall
next Fall = Winter
warm :: Season -> Bool
warm Winter = False
warm Spring = True
warm Summer = True
warm Fall = True
```

Eq, Ord

instance	Εq	Season	whe:	re
Winter	==	Winter	=	True
Spring	==	Spring	=	True
Summer	==	Summer	=	True
Fall	==	Fall	=	True
_	==	_	=	False

instance	Ord	Season	whe	ere
Spring	<= 1	Winter	=	False
Summer	<= 1	Winter	=	False
Summer	<= ;	Spring	=	False
Fall	<= 1	Winter	=	False
Fall	<= ;	Spring	=	False
Fall	<= ;	Summer	=	False
_	<=	_	=	True

instand	ce Show	Seas	son where
show	Winter	=	"Winter"
show	Spring	=	"Spring"
show	Summer	=	"Summer"
show	Fall	=	"Fall"

Class Enum

class Enum a where

toEnum	:: Int -> a	
fromEnum	:: a -> Int	
succ, pred	:: a -> a	
enumFrom	:: a -> [a]	[x]
enumFromTo	:: a -> a -> [a]	[xy]
enumFromThen	:: a -> a -> [a]	[x,y]
enumFromThenTo	:: a -> a -> a -> [a]	[x,yz]

-- minimum definition: toEnum, fromEnum succ x = toEnum (fromEnum x + 1) pred x = toEnum (fromEnum x - 1) enumFrom x = map toEnum [fromEnum x ..] enumFromTo x y = map toEnum [fromEnum x .. fromEnum y] enumFromThen x y = map toEnum [fromEnum x, fromEnum y ..] enumFromThenTo x y z = map toEnum [fromEnum x, fromEnum y .. fromEnum z]

Here's another type class, Enum, used for giving meaning to expressions like [x..y].

Syntactic sugar

- -- [x..] = enumFrom x
- -- [x..y] = enumFromTo x y
- -- [x, y..] = enumFromThen x y
- -- [x, y...z] = enumFromThenTo x y z

Enumerating Int

instance Enum Int where toEnum x = x fromEnum x = x succ x = x+1pred x = x-1enumFrom x = iterate (+1) x enumFromTo x y = takeWhile (<= y) (iterate (+1) x) enumFromThen x y = iterate (+(y-x)) x enumFromThenTo x y z = takeWhile (<= z) (iterate (+(y-x)) x) iterate :: $(a \rightarrow a) \rightarrow a \rightarrow [a]$ iterate f x = x: iterate f (f x)takeWhile :: $(a \rightarrow Bool) \rightarrow [a] \rightarrow [a]$ takeWhile p [] = [] takeWhile p(x:xs) | p x = x : takeWhile p xs| otherwise = []

Now we can declare Int as an instance of Enum.

Enumerating Seasons

instance Enum Season where

fromEnum Winter = 0 fromEnum Spring = 1 fromEnum Summer = 2

fromEnum Fall = 3

toEnum 0 = Winter toEnum 1 = Spring toEnum 2 = Summer toEnum 3 = Fall

Here is Season defined as an instance of Enum.

Deriving Seasons

Haskell uses types to write code for you!

Seasons, revisited

```
next :: Season -> Season
next x = toEnum ((fromEnum x + 1) `mod` 4)
warm :: Season -> Bool
warm x = x `elem` [Spring .. Fall]
-- [Spring .. Fall] = [Spring, Summer, Fall]
```

Having defined Season as an instance of Enum, we can give better definitions of next and warm.

Part VII

Shape

Shape

type	Radius	=	Float
type	Width	=	Float
type	Height	=	Float
data	Shape	=	Circle Radius Rect Width Height
area	:: Shape	->	> Float
area	(Circle	r)	= pi * r^2
area	(Rect w	h)	= w * h

Eq, Ord, Show

instance Eq Shape where Circle r == Circle r' = r == r' Rect w h == Rect w' h' = w == w' && h == h' ______ = False

instance Ord Shape where Circle r <= Circle r' = r < r' Circle r <= Rect w' h' = True Rect w h <= Rect w' h' = w < w' || (w == w' && h <= h') _ <= _ False</pre>

instance Show Shape where show (Circle r) = "Circle " ++ showN r show (Radius w h) = "Radius " ++ showN w ++ " " ++ showN h showN :: (Num a) => a -> String showN x | x >= 0 = show x

| otherwise = "(" ++ show x ++ ")"

Here's Shape as an instance of Eq, Ord and Show.

Deriving Shapes

Haskell uses types to write code for you!

You get all of that for free using deriving.

Part VIII

Expressions

Expression Trees

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```
eval :: Exp -> Int
eval (Lit n) = n
eval (e :+: f) = eval e + eval f
eval (e :*: f) = eval e * eval f
*Main> eval (Lit 2 :+: (Lit 3 :*: Lit 3))
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*Main> eval ((Lit 2 :+: Lit 3) :*: Lit 3)
```

Eq, Ord, Show

instance Ord Exp where Lit n <= Lit n' = n < n' Lit n <= e' :+: f' = True Lit n <= e' :*: f' = True e :+: f <= e' :+: f' = e < e' || (e == e' && f <= f') e :+: f <= e' :*: f' = True e :*: f <= e' :*: f' = e < e' || (e == e' && f <= f') _ = False

instance Show Exp where show (Lit n) = "Lit " ++ showN n show (e :+: f) = "(" ++ show e ++ ":+:" ++ show f ++ ")" show (e :*: f) = "(" ++ show e ++ ":*:" ++ show f ++ ")"

Here's Exp as an instance of Eq, Ord and Show.

Deriving Expressions

```
data Exp = Lit Int
    | Exp :+: Exp
    | Exp :*: Exp
    deriving (Eq, Ord, Show)
```

Haskell uses types to write code for you!

You get all of that for free using deriving.