

Informatics 1

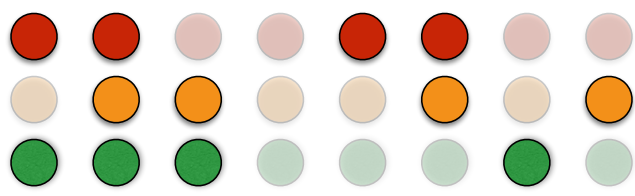
Computation and Logic

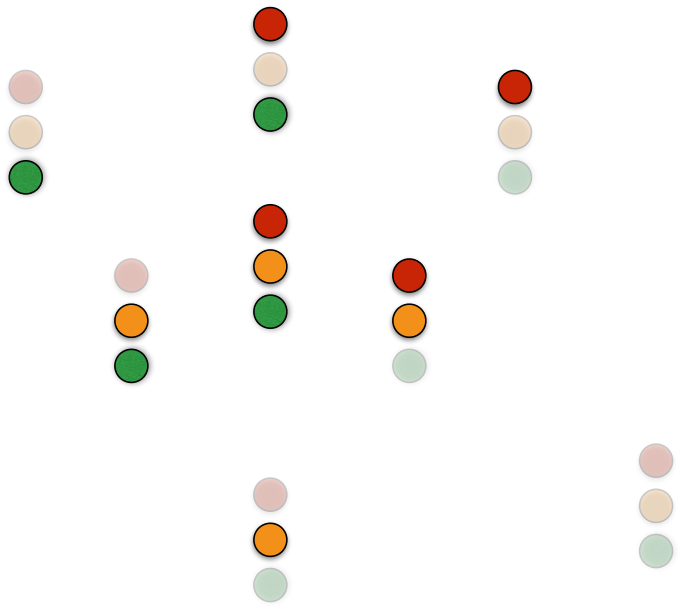


Sets of States: Venn Diagrams and Truth Tables

Michael Fourman
@mp4man

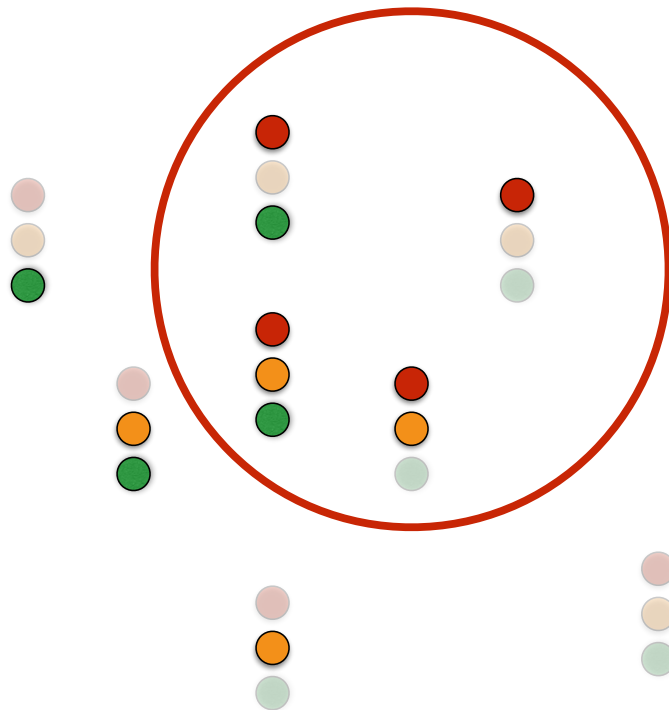
properties & sets
boolean circuits
boolean formulæ
boolean functions





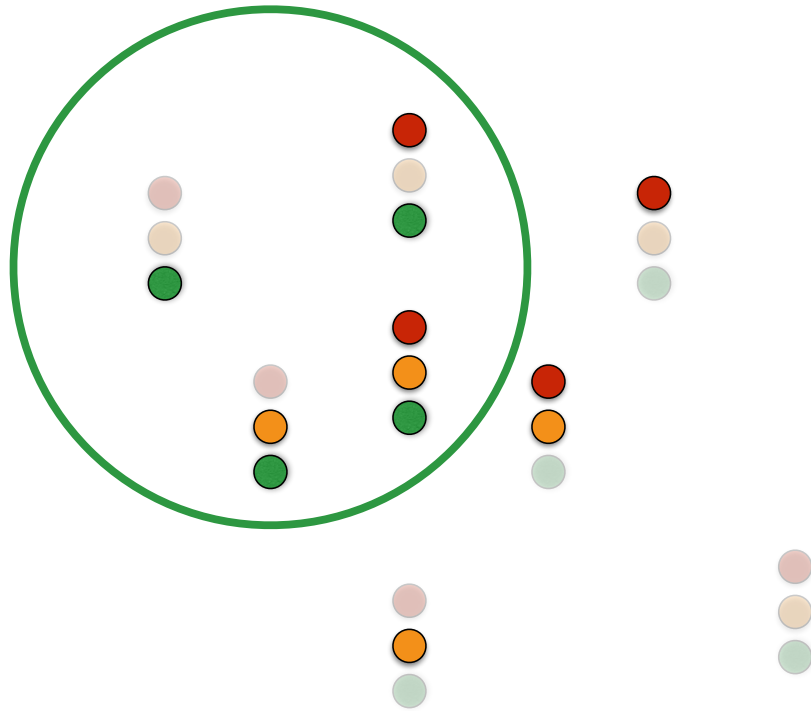
$$\{x \mid R(x)\}$$

R	A	G
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



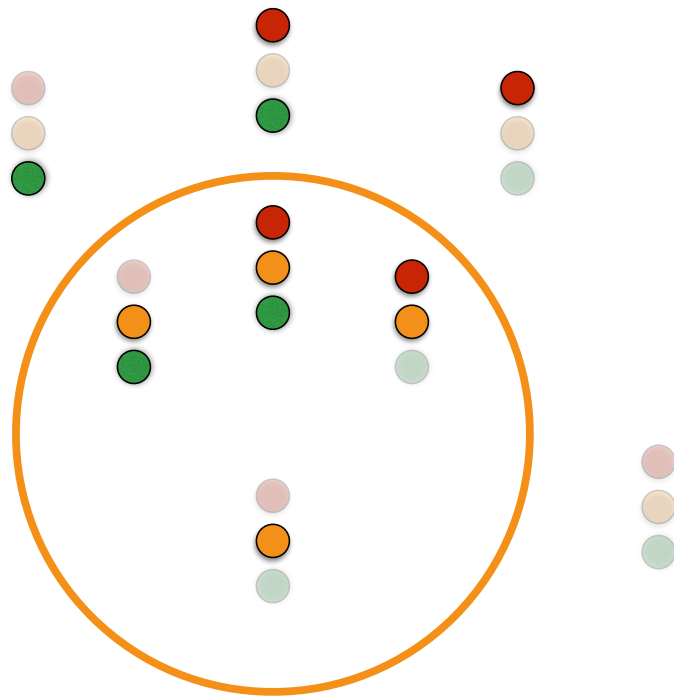
$$\{x \mid G(x)\}$$

R	A	G
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



$$\{x \mid A(x)\}$$

R	A	G
0	0	0
0	0	1
0	1	0
0	1	1
1	0	0
1	0	1
1	1	0
1	1	1



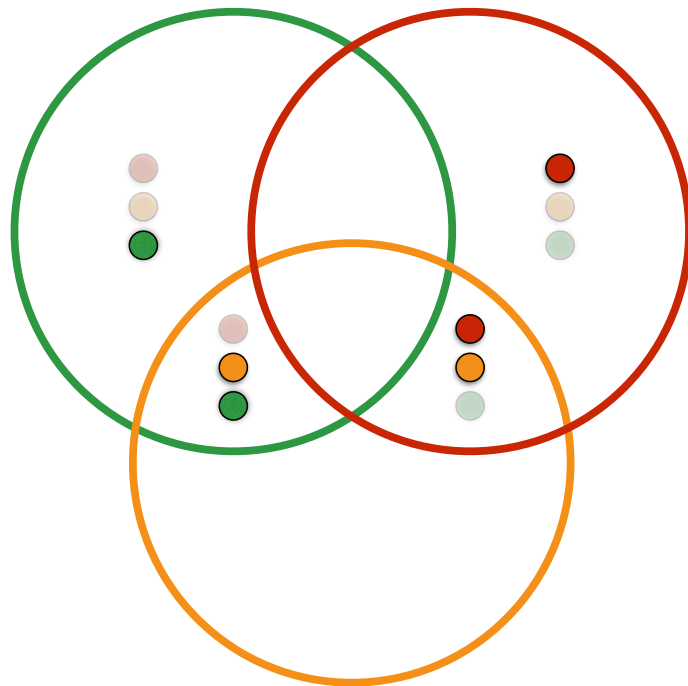
Venn (1834–1923)



xor

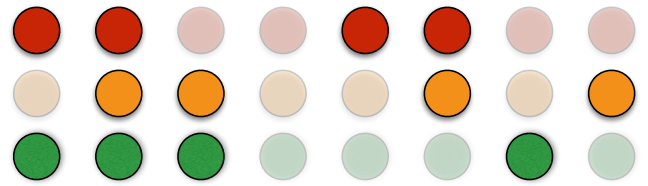
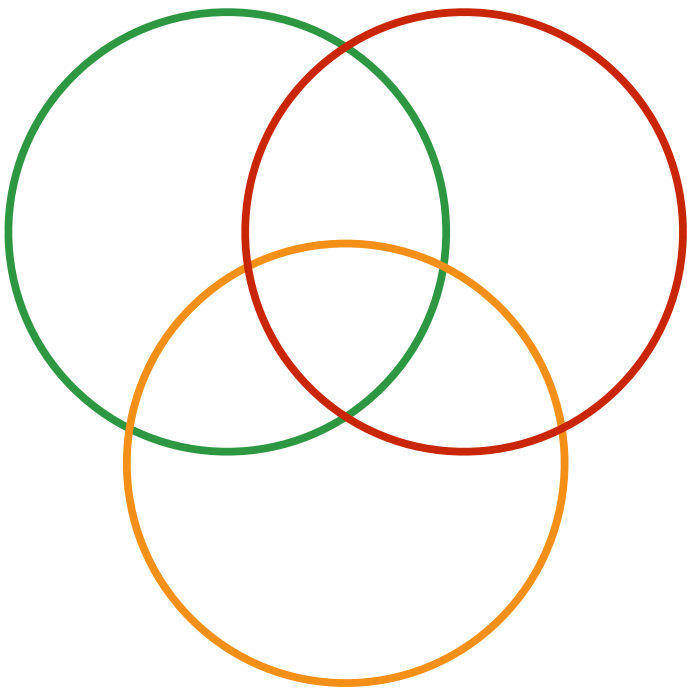
$$\{x \mid G(x) \oplus R(x)\}$$

R	A	G	$G \oplus R$
0	0	0	0
0	0	1	1
0	1	0	0
0	1	1	1
1	0	0	1
1	0	1	0
1	1	0	1
1	1	1	0



xor

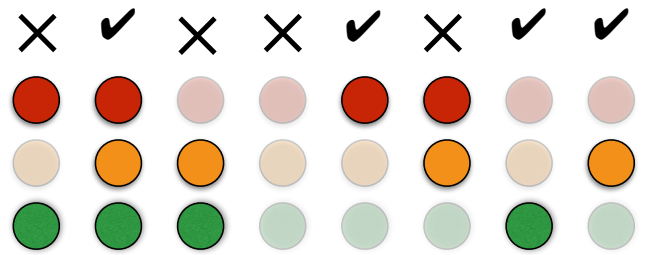
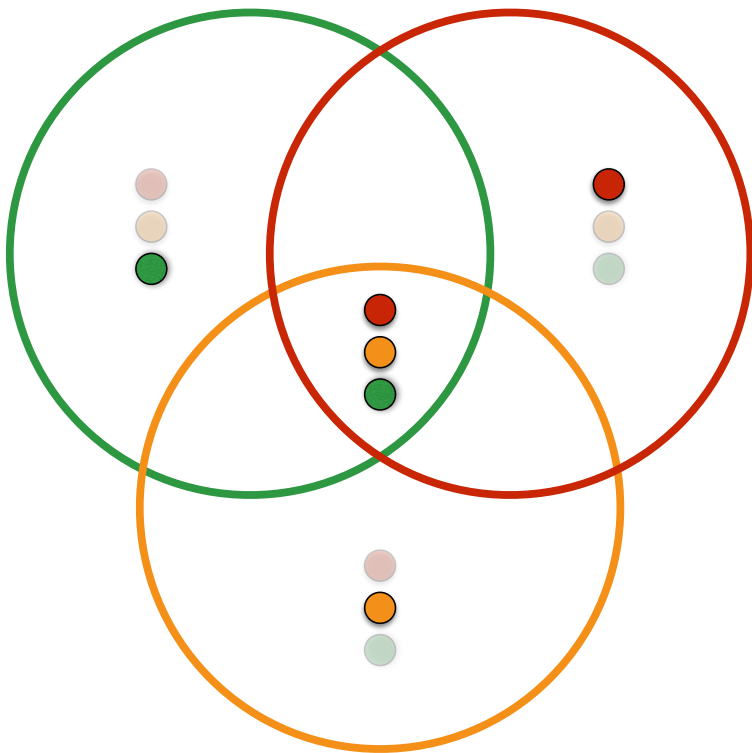
$$\{x \mid G(x) \oplus (R(x) \oplus A(x))\}$$



??

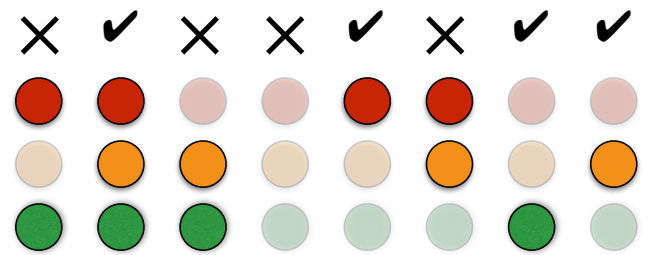
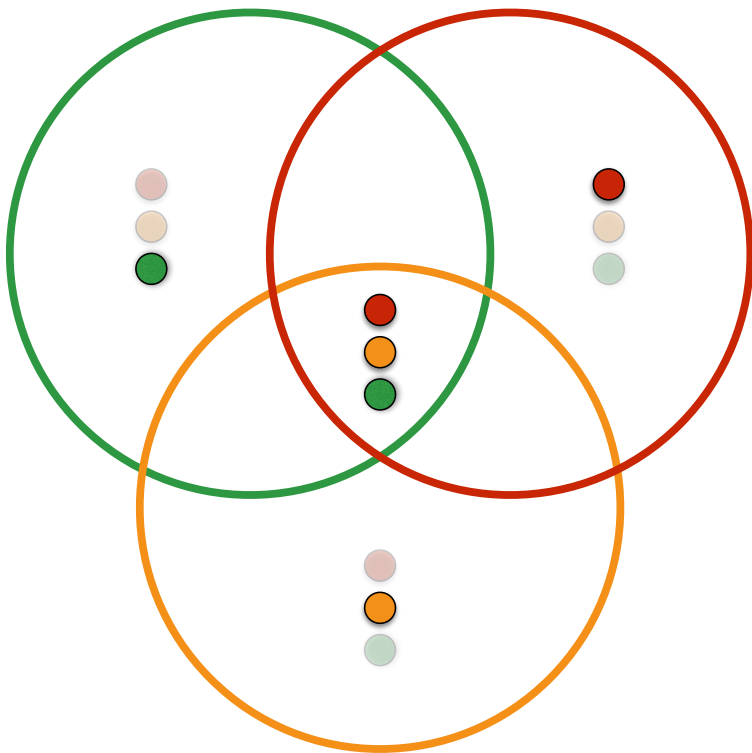
xor

$$\{x \mid G(x) \oplus (R(x) \oplus A(x))\}$$



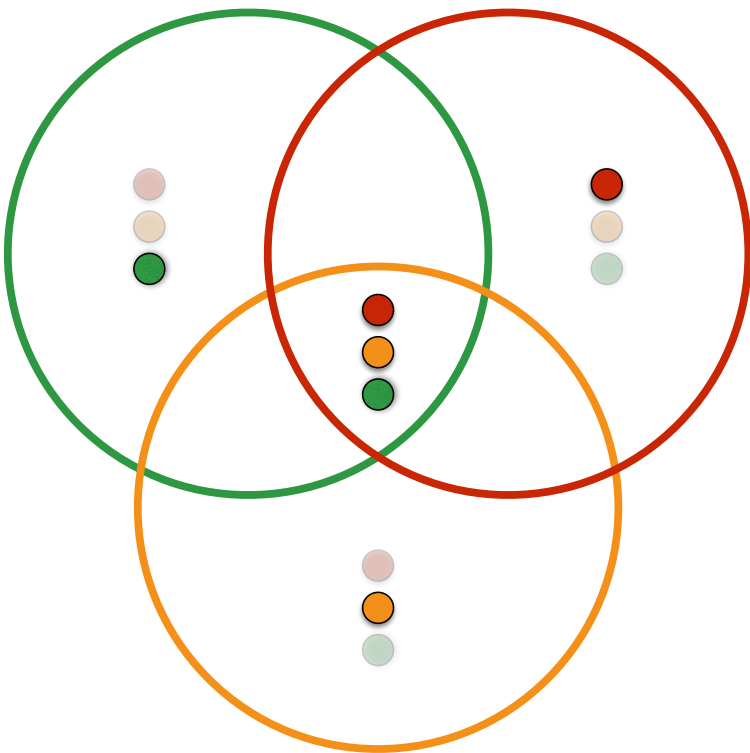
xor

$$\{x \mid G(x) \oplus R(x) \oplus A(x)\}$$



xor

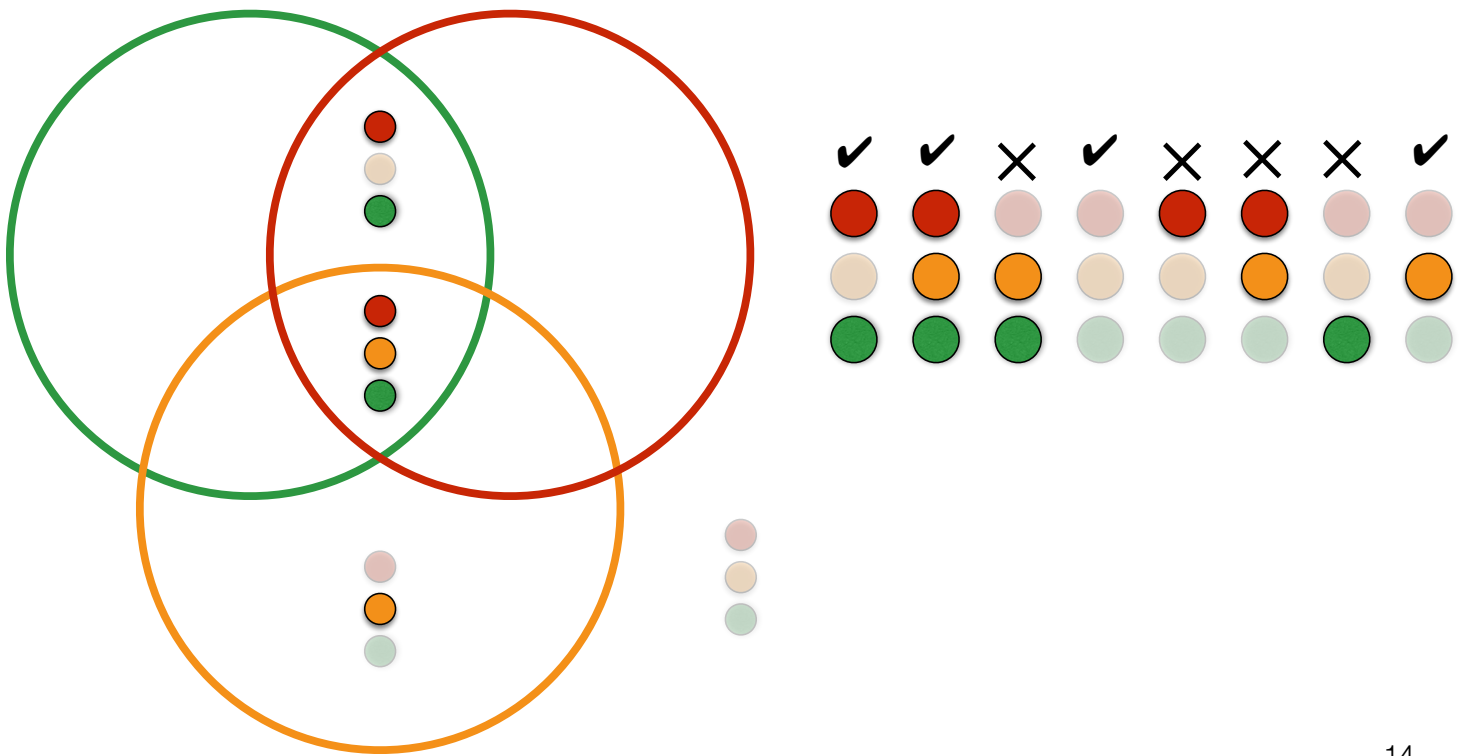
$$\{x \mid G(x) \oplus R(x) \oplus A(x)\}$$



R	A	G	$R \oplus A$	$R \oplus A \oplus G$
0	0	0	0	0
0	0	1	0	1
0	1	0	1	1
0	1	1	1	0
1	0	0	1	1
1	0	1	1	0
1	1	0	0	0
1	1	1	0	1

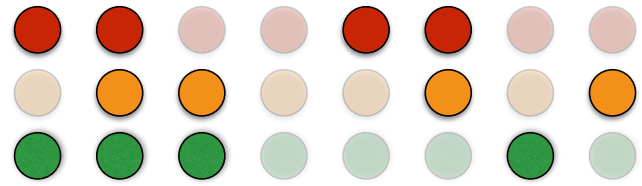
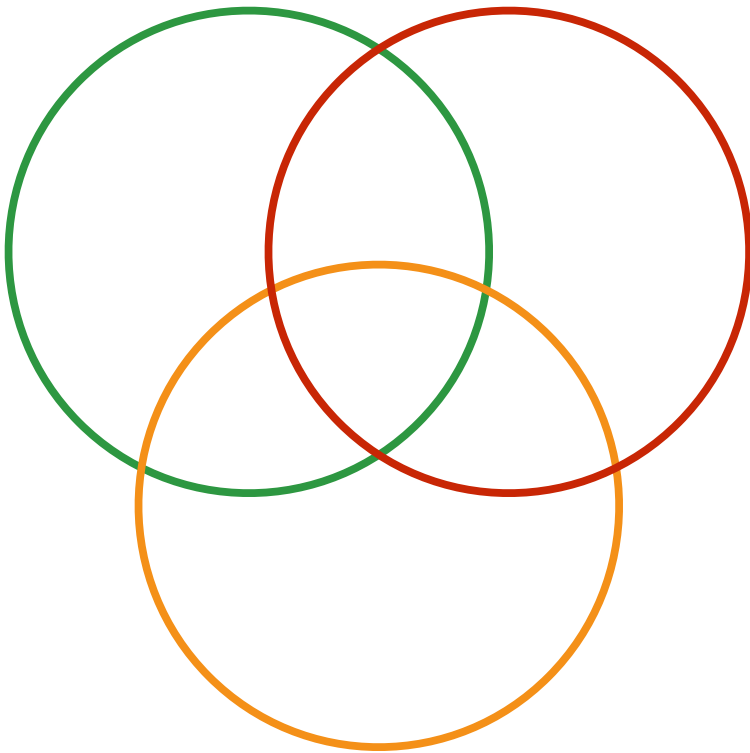
iff

$$\{x \mid G(x) \leftrightarrow R(x)\}$$



iff

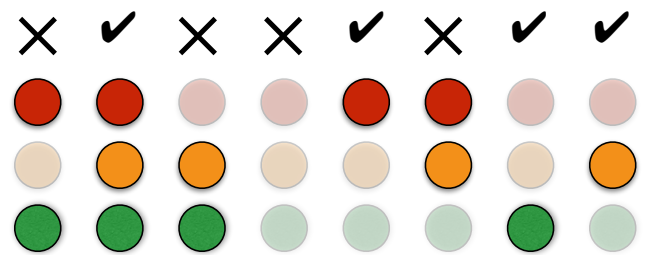
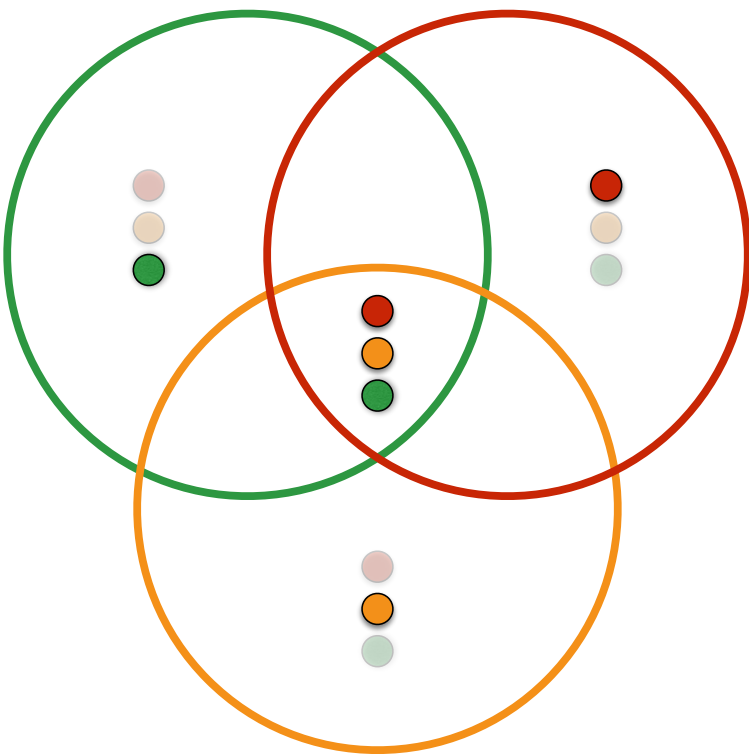
$$\{x \mid G(x) \leftrightarrow (R(x) \leftrightarrow A(x))\}$$



??

iff

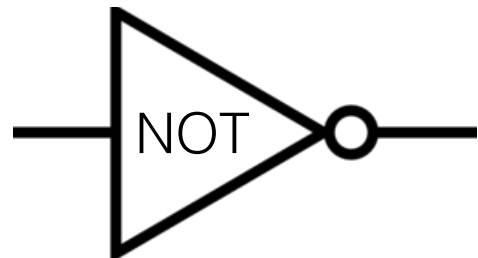
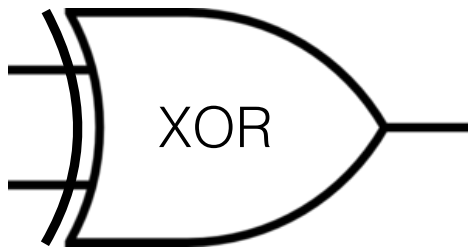
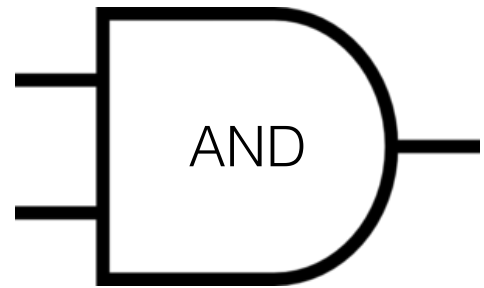
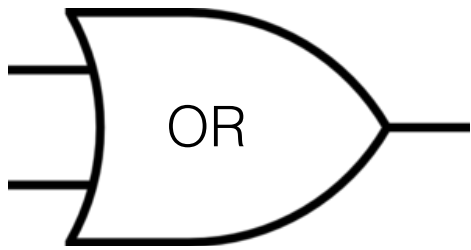
$$\{x \mid G(x) \leftrightarrow R(x) \leftrightarrow A(x)\}$$

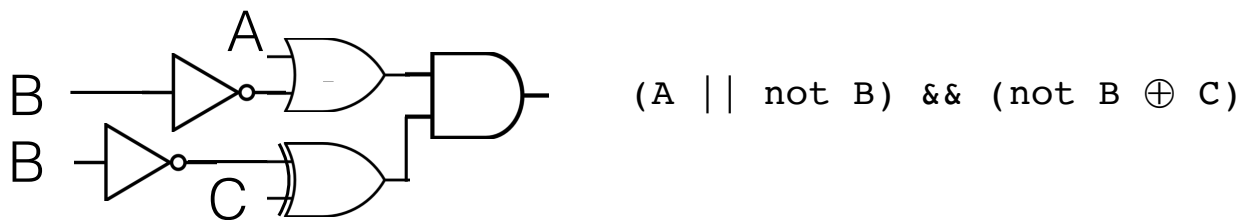
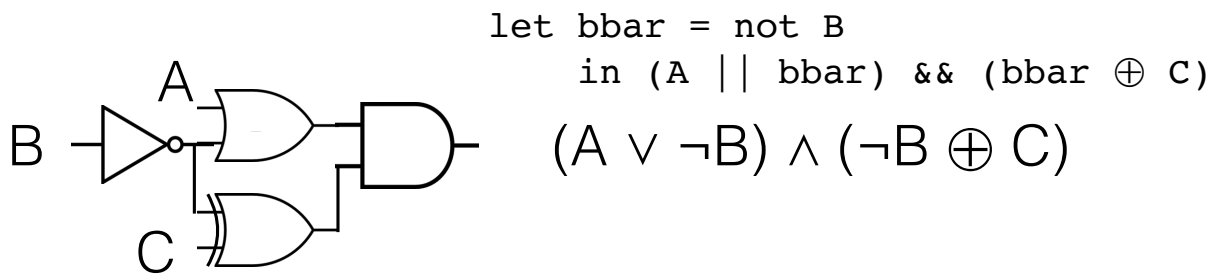


$$G(x) \leftrightarrow R(x) \leftrightarrow A(x)$$

≡

$$G(x) \oplus R(x) \oplus A(x)$$

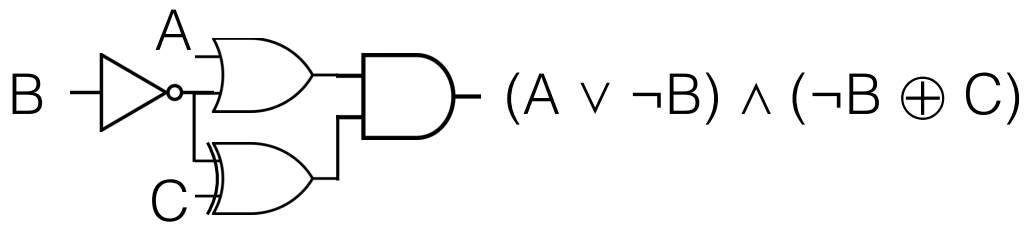




Exercise: define \oplus in Haskell

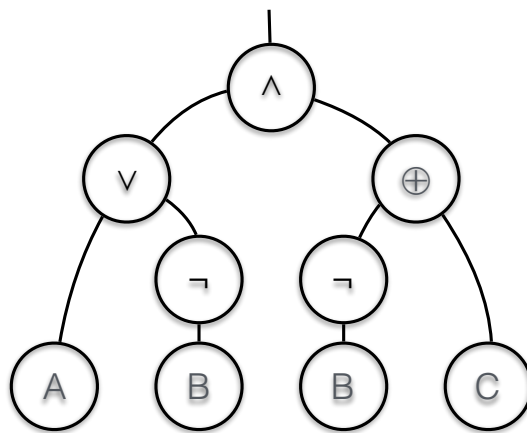
Circuit

Formula

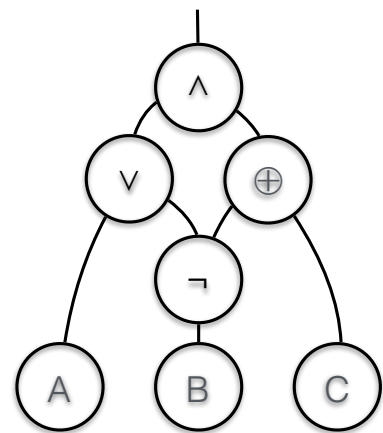


ABC	$A \vee \neg B$	$\neg B \oplus C$	out
000	1	1	1
001	1	0	0
010	0	0	0
011	0	1	0
100	1	1	1
101	1	0	0
110	1	0	0
111	1	1	1

Function

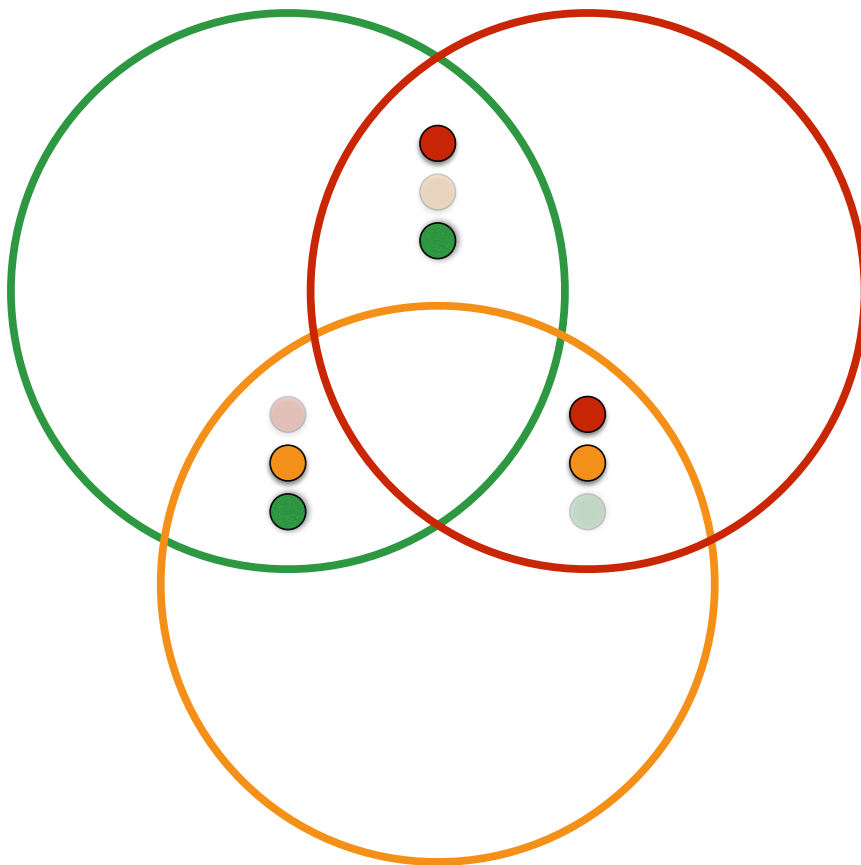


Syntax tree



DAG

Find a proposition



??

Basic Boolean operations

1, \top

\vee

\wedge

\neg

0, \perp



Boole (1815 – 1864)

true, top
disjunction, or
conjunction, and
negation, not
false, bottom

$$\mathbb{Z}_2 = \{0, 1\}$$

+	0	1
0	0	1
1	1	0

$$x \wedge y \equiv xy$$

$$x \vee y \equiv x + y - xy$$

$$\neg x \equiv 1 - x$$

\vee	0	1
0	0	1
1	1	1

\times	0	1
0	0	0
1	0	1

Here, we use arithmetic mod 2

\wedge	0	1
0	0	0
1	0	1

-	
0	0
1	1

The same equations work if we use ordinary arithmetic!

\neg	
0	1
1	0

The algebra of sets

$$\mathcal{P}(S) = \{X \mid X \subseteq S\}$$

$$X \vee Y = X \cup Y$$

union

$$X \wedge Y = X \cap Y$$

intersection

$$\neg X = S \setminus X$$

complement

$$0 = \emptyset$$

empty set

$$1 = S$$

entire set

Derived Operations

Definitions:

$x \rightarrow y \equiv \neg x \vee y$	implication
$x \leftarrow y \equiv x \vee \neg y$	
$x \leftrightarrow y \equiv (\neg x \wedge \neg y) \vee (x \wedge y)$	equality (iff)
$x \oplus y \equiv (\neg x \wedge y) \vee (x \wedge \neg y)$	inequality (xor)

Some equations:

$$\begin{aligned}x \leftrightarrow y &= (x \rightarrow y) \wedge (x \leftarrow y) \\x \oplus y &= \neg(x \leftrightarrow y) \\x \oplus y &= \neg x \oplus \neg y \\x \leftrightarrow y &= \neg(x \oplus y) \\x \leftrightarrow y &= \neg x \leftrightarrow \neg y\end{aligned}$$