

We can use the rules to show this is universally valid, or, if it is not, to generate a counterexample, a model in which

$$\frac{1}{1} \not= \frac{\Delta}{\Delta}$$
some $\Lambda \Gamma$ is not $\nabla \Delta$

Can we use the rules to show this is somewhere valid?

We say the sequent is **satisfiable** if we can find a model in which

some
$$\bigwedge \Gamma$$
 is $\bigvee \Delta$

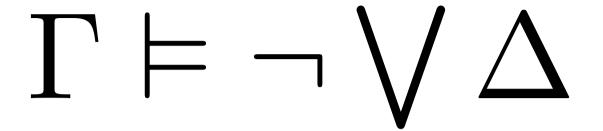
Can we use the rules to show this is somewhere valid?

We say the sequent is satisfiable if we can

find a model where

some
$$\ \ \Gamma$$
 is $\ \ \ \Delta$

$$\Gamma \not\models \neg \ \ \ \Delta$$



We can use the rules to show this is universally valid, or, if it is not, to generate a counterexample, which shows

$$\Gamma \not\models \neg \bigvee \Delta$$

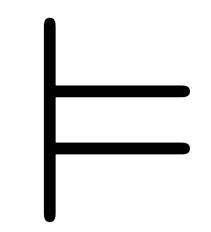
$$\Gamma \vDash \neg \lor \Delta$$
 $\Gamma, \bigvee \Delta \vDash$

We can use the rules to show this is universally valid,

$$\Gamma, \bigvee \Delta$$
 is inconsistent

or, if it is not, to generate a counterexample, a model in which

$$\Gamma \not\models \neg \bigvee \Delta$$
some $\bigwedge \Gamma$ is $\bigvee \Delta$



$$\bigwedge \varnothing \vDash \bigvee \varnothing$$
$$\top \vDash \bot$$

which is only valid in the empty universe

which is only valid in the empty universe

$$\varnothing \vDash \varnothing$$
 $a \vDash b \quad (a = b = \varnothing = \bot)$
 $\bot \vDash \bot$

which is universally true

This is a type error

— but for a mathematician

a set is just a set

there is only one emptyset