Chart Parsing

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Review Top-down Parsing

Chart Parsing

Overview Charts as Graphs The Basic Idea Example States

The Earley Algorithm

Parsing Example Left Recursion

Parsing

Parsing with a CFG is the task of assigning a correct tree (or derivation) to a string given some grammar. A correct tree is:

consistent with the grammar, and

► the leaves of the tree cover all and only the words in the input. There may be a very large number of correct trees for any given input . . .

Problems that arise

- Left Recursion
- Ambiguity
- Inefficiencies due to backtracking

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Common substructures

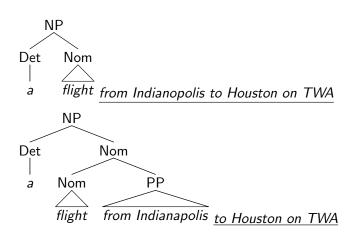
- Despite ambiguity and backtracking there are commons substructures to be taken advantage of.
- Consider parsing the following NP: a flight from Indianapolis to Houston on TWA

with the following rules:

 $NP \rightarrow Det Nom$ $Nom \rightarrow Nom PP$

What happens with a top-down parser?

Backtracking

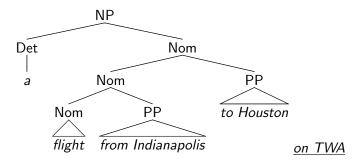


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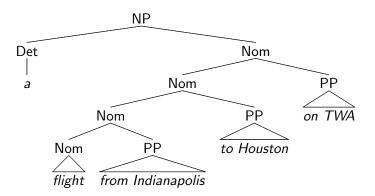
Backtracking, cont



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Backtracking, cont



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Overview Charts as Graphs The Basic Idea Example States

Dynamic Programming

Our current algorithm builds valid trees, discards them during backtracking, then rebuilds them.

▶ the subtree for *a flight* was derived 4 times.

Dynamic programming is one answer to problems that have sub-problems that get solved again and again We'll consider an algorithm that fills a table with solutions to subproblems that:

- does a parallel top-down search with bottom-up filtering
- does not do repeated work
- solves the left-recursion problem

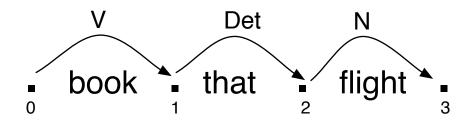
Overview Charts as Graphs The Basic Idea Example States

Dynamic Programming and Parsing

- Systematically fill in tables of solutions to subproblems.
- When complete, the table contains all possible solutions to all of the subproblems needed to solve the whole problem
- For parsing:
 - the table stores subtrees for constituents.
 - Solves reparsing inefficiencies, because subtrees are not reparsed but looked up.
 - Solves ambiguity explosions, because the table *implicitly* stores all parses.
 - Each subtree is represented only once and shared by all parses that need it.

Overview Charts as Graphs The Basic Idea Example States

Lexical Edges

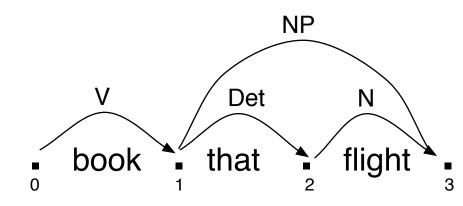


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Charts as Graphs

Nonterminal Edges: NP

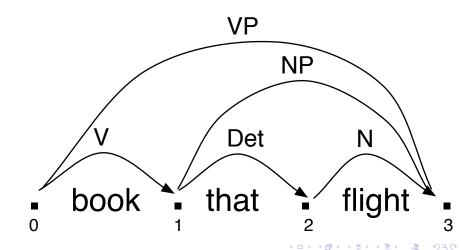


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Nonterminal Edges: VP

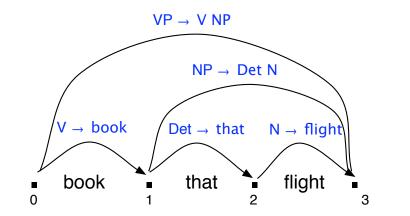


Overview Charts as Graphs The Basic Idea Example States

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Nonterminal Edges: Rules

Useful to label arcs with rules rather than categories:



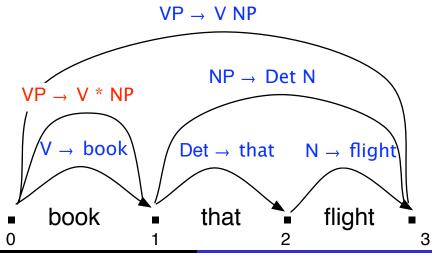
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Incomplete Edges

- Recall that the parser can make predictions on the basis of rules:
 - I'm trying to expand an VP, and I've found a V so I'll start looking for an NP.
- Record incomplete constituents with dotted rules:
 - \blacktriangleright Dot on the RHS of a rule shows what weve found already: $VP \rightarrow V ~\bullet~ NP$
 - We can use this as a label for an incomplete constituent.

Overview Charts as Graphs The Basic Idea Example States

Incomplete Edges, cont.



The Basic Idea

Dynamic Programming and Parsing

The Earley algorithm:

- fills a table (the chart) in a single left-to-right pass over the input.
- The chart will be size N + 1, where N is the number of words in the input.
- Chart entries are associated with the gaps between the words — like slice indexing in Python.
- For each word position in the sentence, the chart contains a set of edges representing the partial parse trees generated so far.
- So Chart[0] is the set of edges representing the parse before looking at a word.

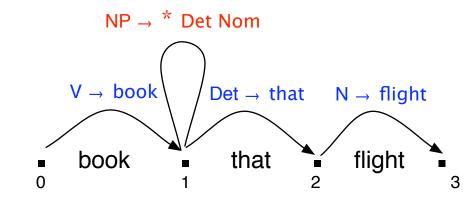
Overview Charts as Graphs **The Basic Idea** Example States

States

- J&M call the chart entries states.
- The chart entries represent three distinct kinds of things:
 - completed constituents;
 - in-progress constituents; and
 - predicted constituents
- The three kinds of states correspond to different dot locations in dotted rules:

Overview Charts as Graphs The Basic Idea Example States

Incomplete NP Edge: Self-loop



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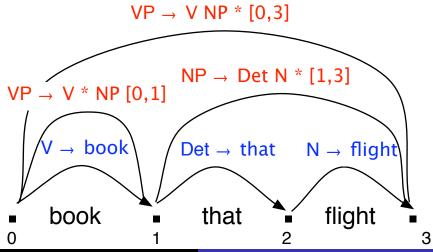
Outline Overview Review Top-down Parsing Charts as Grap Chart Parsing The Basic Idea The Earley Algorithm Example States

States, cont.

- Given dotted rules like those we've just seen, we need to record:
 - where the represented constituent is in the input, and
 - what its parts are.
- ▶ So we add a pair of coordinates [*x*, *y*] to the state:
 - $A \rightarrow \alpha$, [x, y]
 - x indicates the position in input where the state begins
 - y indicates position of dot

Overview Charts as Graphs The Basic Idea Example States

Example with coordinates



Outline Overview Review Top-down Parsing Chart Parsing The Earley Algorithm Earley Algorithm

States, cont.

Example states in parsing Book that flight:

1. S \rightarrow \bullet VP, [0,0]

- First 0 indicates that the constituent begins at the start of the input.
- Second 0 indicates that the dot also begins at start of input, and thus indicates a top-down prediction.
- 2. NP \rightarrow Det Nom, [1,2]
 - the NP begins at position 1
 - the dot is at position 2
 - Det has been successfully completed
 - Nom is predicted next

 Outline
 Overview

 Review Top-down Parsing
 Charts as Graphs

 Chart Parsing
 The Basic Idea

 The Earley Algorithm
 Example States

States, cont.

- 1. VP \rightarrow V NP •, [0,3]
 - VP is completed
 - no further predictions from this rule
 - a successful parse that spans the entire input

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Outline Overview Review Top-down Parsing Charts as Graphs Chart Parsing The Basic Idea The Earley Algorithm Example States

Success

The final answer is found by looking at the last column of the table. In particular, for an input of N words, if we find the following kind of state in the chart then we've succeeded:

 $S \rightarrow \alpha \bullet, [0,N]$

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Parsing Example Left Recursion

Parsing

Parsing is sweeping through the chart creating the three kinds of states as we go. States are never removed, and we never backtrack.

- New predicted states are based on existing table entries (predicted, or in-progress) that predict a certain constituent at that spot.
- New *in-progress* states are created by updating older states to reflect the fact that previously expected completed constituents have been located.
- New completed states are created when the dot in an in-progress state moves to the end.

Parsing Example Left Recursion

More Specifically

- 1. Predict all the states you can.
- 2. Read an input.
 - See what predictions you can match.
 - Extend matched states, add new predictions.
 - Go to next state (goto 2)
- 3. At the end, see if state[N + 1] contains a complete S.

Parsing Example Left Recursion

Earley Algorithm

The Earley algorithm has three main functions that do all the work: Predictor: Adds predictions into the chart

- Completer: Moves the dot to the right when new constituents are found
 - Scanner: Reads the input words and enters states representing those words into the chart

Parsing Example Left Recursion

Predictor

```
procedure PREDICTOR((A \rightarrow \alpha \bullet B \beta, [i, j]))
for each (B \rightarrow \gamma) in GRAMMAR-RULES-FOR(B, grammar)
do
ENQUEUE((B \rightarrow \bullet \gamma, [j, j]), chart[j])
end
```

- Intuition: new states represent top-down expectations.
- Applied when a state has a non-terminal to the right of a dot that is not a part-of-speech.
- Generates one new state for each alternative expansion of the non-terminal in the grammar.
- Adds states to the same chart entry as generating state.

Parsing Example .eft Recursion

Completer

procedure COMPLETER($(B \rightarrow \gamma \bullet, [j, k])$) for each $(A \rightarrow \alpha \bullet B \beta, [i, j])$ in chart[j] do ENQUEUE($(A \rightarrow \alpha B \bullet \beta, [i, k]), chart[k]$) end

- Intuition: parser has discovered a constituent, so must find and advance states that were looking for this grammatical category at this position in input.
- Applied when dot has reached right end of rule.
- New states are generated by copying old state and advancing dot over expected category.

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Adds new states to same chart entry as generating state.

Parsing Example .eft Recursion

Scanner

procedure SCANNER($(A \rightarrow \alpha \bullet B \beta, [i, j])$) if $B \in PARTS-OF-SPEECH(word[j])$ then ENQUEUE($(B \rightarrow word[j] \bullet, [j, j + 1])$, chart[j+1])

- New states for predicted part-of-speech.
- Applicable when part-of-speech is to the right of a dot.
- Adds states to next chart entry.

Note: Earley parser uses top-down predictions to help disambiguate part-of-speech ambiguities. Only those parts-of-speech of a word that are predicted by some state will find their way into the chart.

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Parsing Example Left Recursion

Mini grammar and lexicon

$S \rightarrow NP VP \mid Aux NP VP \mid VP$	$Det \to \mathit{that} \mid \mathit{this} \mid \mathit{a}$
$NP \to Det Nom \mid PropN$	$N \to \mathit{book} \mid \mathit{flight} \mid \mathit{meal}$
$Nom \to Nom PP \mid N Nom$	$V \rightarrow book \mid include \mid prefer$
$PP \to P NP$	$Aux \rightarrow does$
$VP \to V \mid V NP$	$P \to \mathit{from} \mid \mathit{to} \mid \mathit{on}$
$Nom \to N PP \mid N Nom$	$PropN \to \mathit{Houston} \mid \mathit{TWA}$

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Example: Chart[0] and Chart[1]

Chart[0]			
$\gamma \rightarrow \bullet S$	[0,0]	Dummy start state	
$S \rightarrow \bullet NP VP$	[0,0]	Predictor	
$S \rightarrow$ • Aux NP VP	[0,0]	Predictor	
$S \rightarrow \bullet VP$	[0,0]	Predictor	
$\textit{NP} \rightarrow \textit{\bullet} \textit{Det NOMINAL}$	[0,0]	Predictor	
$\textit{NP} \rightarrow \bullet \textit{Proper-Noun}$	[0,0]	Predictor	
$V\!P \ o \ ullet$ VP \to Verb	[0,0]	Predictor	
$VP \rightarrow \bullet Verb NP$	[0,0]	Predictor	

Chart[1]				
Verb $ ightarrow$ book $ullet$	[0,1]	Scanner		
$V\!P \ o \ V\!erb$ $ullet$	[0,1]	Completer		
$S \ ightarrow \ VP$ $ullet$	[0,1]	Completer		
$VP \rightarrow Verb \bullet NP$	[0,1]	Completer		
$NP \rightarrow Opt NOMINAL$	[1 1]	Predictor	-12	900
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Example: Chart[1] and Chart[2]

Chart[1]		
Verb $ ightarrow$ book $ullet$	[0,1]	Scanner
$V\!P \ o \ V\!erb$ $ullet$	[0,1]	Completer
$S \ ightarrow \ VP$ $ullet$	[0,1]	Completer
$VP \rightarrow Verb \bullet NP$	[0,1]	Completer
$\textit{NP} \rightarrow \bullet \textit{Det NOMINAL}$	[1,1]	Predictor
$NP \rightarrow \bullet Proper - Noun$	[1,1]	Predictor

Chart[2]		
Det \rightarrow that \bullet	[1,2]	Scanner
$NP \rightarrow Det \bullet NOMINAL$	[1,2]	Completer
NOMINAL \rightarrow • Noun	[2,2]	Predictor
NOMINAL \rightarrow • Noun NOMINAL	[2,2]	Predictor

Parsing Example Left Recursion

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Example: Chart[3]

Chart[3]		
Noun \rightarrow flight •	[2,3]	Scanner
NOMINAL \rightarrow Noun •	[2,3]	Completer
$NOMINAL \rightarrow Noun \bullet NOMINAL$	[2,3]	Completer
NP $ ightarrow$ Det NOMINAL •	[1,3]	Completer
$VP \ ightarrow \ Verb \ NP \ ullet$	[0,3]	Completer
$S \rightarrow VP \bullet$	[0,3]	Completer
NOMINAL \rightarrow • Noun	[3,3]	Predictor
NOMINAL \rightarrow • Noun NOMINAL	[3,3]	Predictor

Parsing Example Left Recursion

Examples: Left Recursion

What about parsing the NP*a flight from Denver to Boston* with the following rules:

 $NP \rightarrow NP PP$ $NP \rightarrow Det Nom$ $NP \rightarrow Proper-Noun$

- We construct the state (NP → NP PP, [0,0]) and add it to chart[0]
- The PREDICTOR function then requires us to find a rule which expands the (non-lexical) category immediately to the right of the dot.
- So let's pick the first rule above, and ENQUEUE the state (NP → • NP PP, [0,0]).
- But this is already in the state, so we don't add it again.

Reading

- Read section 10.4 of J&M
- Read the NLTK-Lite Tutorial on Chart Parsing

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