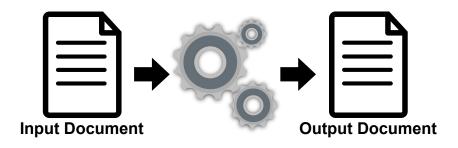
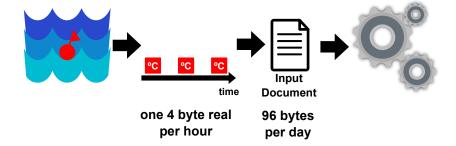
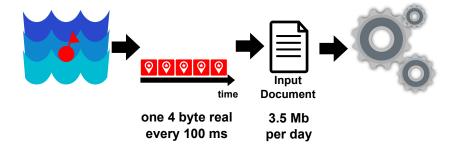
# Data Stream Processing

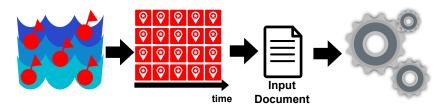
Homework 1 is due this Friday the 20th of October

# Data Processing so far ...









one million 4 byte reals every 100 ms

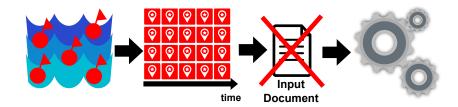
3.5 Tb per day

Stream of large unbounded data

too large for memory

too high latency for disk

We need real time processing!



Process data stream directly

### Data Streams

#### What is a Data Stream?

#### Definition (Golab and Ozsu, 2003

A data stream is a real-time, continuous, ordered (implicitly by arrival time of explicitly by timestamp) sequence of items. It is impossible to control the order in which items arrive, nor it is feasible to locally store a stream in its entirety.

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- continous and sequential input
- typically unpredictable input rate
- can be large amounts of data
- not error free

# Data Stream Applications

- Online, real time processing
- Event detection and reaction
- Aggregation
- Approximation

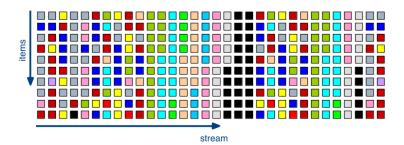
Stock monitoring

Stock monitoring
Website traffic monitoring

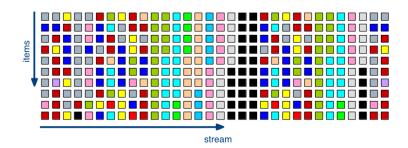
Stock monitoring
Website traffic monitoring
Network management

Stock monitoring
Website traffic monitoring
Network management
Highway traffic

## Data Stream Characteristics

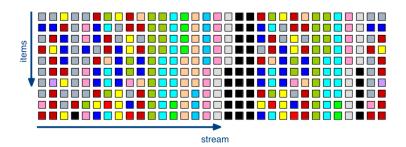


#### Data Stream Characteristics



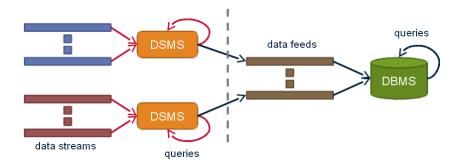
 All items have the same structure. For example a tuple or object: (sender, recipient, text body)

#### Data Stream Characteristics



- All items have the same structure. For example a tuple or object: (sender, recipient, text body)
- timestamps: explicite vs. implicite, physical vs. logical

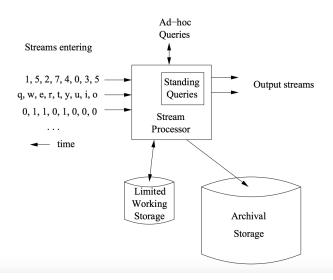
#### Database Management vs. Data Stream Management



### DBMS vs. DSMS

Feature	DBMS	DSMS
Model	persistent relation	transient relation
Relation	tuple set/bag	tuple sequence
Data update	modifications	appends
Query	transient	persistent
Query answer	exact	approximate
Query evaluation	arbitrary	one pass
Query plan	fixed	adaptive

#### DSMS Architecture



# Data Stream Mining

# Data Stream Mining

- event detection and reaction
- counting frequency of specific items
- pattern detection
- aggregation
- approximation
- sampling

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# Resevoir Sampling

# Problem: Sampling

Lines from a large text file

Stream: Sample search engine queries, updated live

# The Simple Way

- Scan the text file, counting lines
- Generate random line numbers [0, |lines|)
- Sort the line numbers
- Scan the text file, outputting selected lines

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- Scan the text file, outputting selected lines

Cost: two scans
Impossible / Impractical for stream

# The Simple Way for a Stream

## Problem: Sample top 1000 queries

- assign each query a random number
- 2 keep the queries with the top 1000 highest random numbers
- discard the rest

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So far not reservoir sampling!

## Sample One Line

Probability of keeping a line and dropping all others?

• keep 1st line:

## Sample One Line

# Probability of keeping a line and dropping all others?

• keep 1st line: 1

## Sample One Line

# Probability of keeping a line and dropping all others?

- keep 1st line: 1
- keep 2nd line:

# Probability of keeping a line and dropping all others?

- keep 1st line: 1
- keep 2nd line:





# Probability of keeping a line and dropping all others?

- keep 1st line: 1
- keep 2nd line:  $\frac{1}{2}$





# Probability of keeping a line and dropping all others?

- keep 1st line: 1
- keep 2nd line:  $\frac{1}{2}$
- keep 3rd line:  $\frac{1}{2}$



# Probability of keeping a line and dropping all others?

- keep 1st line: 1
- keep 2nd line:  $\frac{1}{2}$
- keep 3rd line:  $\frac{1}{2}$
- keep nth line:



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- keep 1st line: 1
- keep 2nd line:  $\frac{1}{2}$
- keep 3rd line:  $\frac{1}{2}$
- keep nth line:  $\frac{1}{2}$



Flip a coin at each line. If it's heads, record the line (and forget the others).

```
#!/usr/bin/env python
import sys
import random
resevoir = sys.stdin.readline().strip()
for line in sys.stdin:
   if random.randint(0,1) == 0:
     resevoir = line.strip()
print(resevoir)
```

Flip a coin at each line. If it's heads, record the line (and forget the others).

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This is biased. The last line has probability  $\frac{1}{2}$ .

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```

This is biased. The last line has probability  $\frac{1}{2}$ . It should be the same probability for each line!

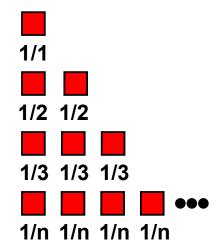
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- keep nth line:

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- keep nth line:  $\frac{1}{n}$

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- keep 3rd line:  $\frac{1}{3}$
- keep nth line:  $\frac{1}{n}$



```
#!/usr/bin/env python
import sys
import random
line_number = 0
for line in sys.stdin:
  if random.randint(0, line_number) == 0:
    resevoir = line.strip()
  line_number += 1
print(resevoir)
```

Line *n* overwrites the resevoir with probability  $\frac{1}{n}$   $\implies$  Uniform sampling

#### Proof Sketch: Induction

Base One line with probability 1.

Inductive Assume n lines were sampled with probability  $\frac{1}{n}$  each. When the n+1th line is added, the resevoir is kept with probability  $\frac{n}{n+1}$ . Thus the first n lines each have probability

$$\frac{1}{n} \cdot \frac{n}{n+1} = \frac{1}{n+1}$$

And the n+1th line also has probability  $\frac{1}{n+1}$  by construction.

2/5 2/5 2/5 2/5 2/5

53

Reservoir Size r = 2

Reservoir Size 
$$r = 1$$
 1/5 1/5 1/5 1/5 1/5 1/5

Reservoir Size  $r = 2$  2/5 2/5 2/5 2/5 2/5 2/5

with reservoir size r and sample count n Substitute an entry with probability:

with reservoir size r and sample count n Substitute an entry with probability:  $\frac{r}{n}$ 

# Sample Multiple Lines Without Replacement

First few lines: Fill the resevoir

Afterwards: Substitute an entry with probability

|samples|

### Summary

Efficiently sample streaming data

Small memory