Energy-Aware Computing

Lecture 10: Power-aware caches

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Outline

- Value compression in caches
- Reducing cache parallel activity in associative caches
 - Phased cache
 - Way prediction
 - Way selection (halting)
- · Controlling cache idle capacity
 - Selective cache ways
- Saving snoop power

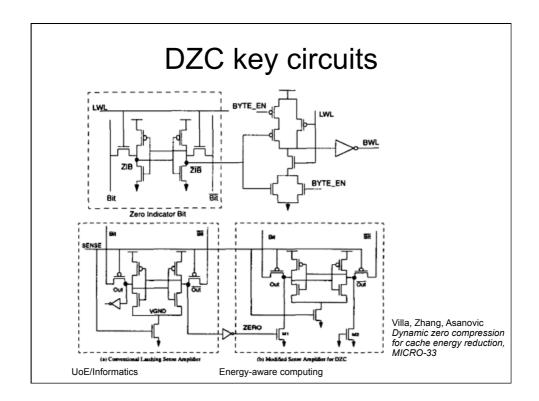
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Value compression

Dynamic zero compression

- Compress zero bytes using a tag bit per byte (zero indicator bit)
- Detect zero when storing and only store the ZIB, the bitlines are not driven
- Byte word line disabled when ZIB = 1
- Special sense amps force output byte to 0 when ZIB=1

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DZC evaluation

- Spice simulation from extracted layout of a 16KB cache in 0.25um
- Array of ZIBs adds 9% area
- Read
 - Latency overhead 2FO4 delays
 - D\$ 26% reduction in power
 - I\$ 10% reduction in power
- Write
 - No delay; zero detection overlaps with tag access
 - Write energy 1/5th for 0x0 compared to non-zero

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Other compression methods

- · Simple compression scheme required
- Frequent value locality
 - Small number of distinct values appear frequently in memory accesses
- · Dictionary based approaches
 - Keep a table of frequent values
 - Store index to dictionary in place of value
- Pack more lines into a cache!
 - Increases the capacity, hit rate

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Activity in associative caches

- · Associativity is desirable, if it comes cheaply
 - It improves the hit rate
- The standard RAM-tagged implementation is not power efficient
 - We know that at most one data bank will contain the address, but all are accessed
- A number of activity reducing methods have been proposed

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Phased caches

Conventional 4 way associative		
➤ Phased		
Full tag access required		
 On a hit, only one of the data sub-blocks is accessed 		
 Significant power savings at the expense of 1 additional cycle for hits 		
 Can be pipelined 		
• Similar to CAM tagged can	ches	

Way prediction

Conventional 4 way	
associative	
Way predicting	

- · Extra prediction unit required
- On correct prediction, only 1 way (tag, data) accessed
- On incorrect prediction, energy is as with standard cache
 - But slower

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Way prediction evaluation

- For I\$, way prediction accuracy is 96% on average, over 90% for all benchmarks
- For D\$, average accuracy is 86%
 - Half benchmarks over 90%
- Improved EDP by 60%(D\$) to 70%(I\$)
- · Compared to phased cache
 - Similar power savings
 - Approx half average cache access time

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Way prediction problems

- · High associativity reduces miss rates
 - Power efficient
- But the success rate of way-prediction drops with increasing associativity
 - More options to choose from
 - Not a serious problem for L1 caches, but increasingly important for other levels
- Cache misses on way-predicting caches are very slow as all options are explored before a miss is declared

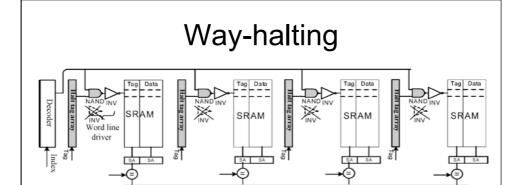
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Way-selection

- · Deterministically select the right way
 - No prediction is made
- Location cache
 - Hold way info for L2 \$ next to L1
 - Update location \$ for L1 replacements
- · Way-halting
- · Decaying Bloom filters

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- Tag storage split btw RAM, CAM mems
- Quick partial look up of tags (using CAM)
- · All mismatches halt tag and data array accesses
- 2% area overhead, no performance penalty, 45-60% energy reduction

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Decaying Bloom Filters

- Use a leakage power saving method to save dynamic power
 - Decay (switch off) inactive cache lines to save leakage power
- · No need to search decayed lines
- A Bloom filter tracks which ways may hold a line
- · Parallel access to all such lines
 - rather than 1 with way-prediction

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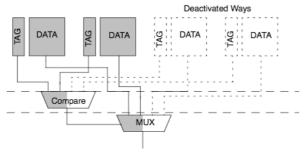
Selective cache ways

- Not all cache needed by all programs all the time
- In sub-banked caches, "ways" are implemented as independent banks
- Easiest (only?) way to adapt cache dynamically: disable one or more ways
 - No accesses to data array, no prechanging,...
 - Tags remain active

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Selective cache ways



- · Dirty data in disabled ways
 - Temporarily enable on hit or coherence access
 - Flushing to L2 would be expensive
- · When to adapt the cache size
 - Software control

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Saving snoop power

- Multiprocessor caches stay coherent by snooping all other accesses
 - Every miss in a cache, searches all other caches
- 54% of all snoops miss in the L2 tags 4processor SMP SPLASH-2 benchmarks
 - Wasted power searching tags

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Jetty - snoop filtering

- Filter out unneeded snooping
- Bloom filter efficient structure for nonmembership test
 - to determine if item not in set
- Three Jettys:
 - Exclude: "I've seen this before and I am sure its not locally cached"
 - Include: captures a superset of cached data
 - Hybrid: check both (eliminates 76%)

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Recent work on caches

- Non-uniform cache access (NUCA)
 - Access to banks that are far away from the core are slower than to banks that are near
 - Either operate at slowest speed or expose the nonuniformity to the processor
 - Caching within the cache: keep frequently used data in nearest bank
- Sharing a number of banks between cores in CMPs
 - Dynamically adjust the size of each core's cache
 - Used in L2 of Intel multicore chips

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Summary

- Value compression
 - Zero value compression
- Cache activity reduction
 - Phased cache
 - Way prediction
 - Way selection (halting)
- · Controlling cache size
 - Selective cache ways
- · Saving snoop power

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