Distributed Systems

Mobile & Sensor Computing

Rik Sarkar

University of Edinburgh Fall 2020

Mobile and Ubiquitous computing

- Devices (computers) are carried by people (mobile)
 - Laptops, phones, watches ...
- They are everywhere
 - Carried by people (mobile)
 - Embedded in the environment
 - Coffee machines, cameras, sensors for light control, elevators...
 - Produce large amounts of data
 - Usage, sensing...

Ubiquitous

Advantages:

- There are computers everywhere
- Everything is "smart"
- Potentially use collaborative distributed computations on these to make them even smarter

Challenges:

- There are more things to go wrong
- Not easy to make things work well coherently
- Consistent platforms for managing ubiquitous devices do not exist (yet)
- Devices do not interoperate easily

Mobile

Advantages:

- The same device is carried by the person easy to give consistent service
- Information whenever, wherever they need
- Devices have sensors potential for sensing the environment and adapting

Disadvantages:

- Connectivity is challenge: data is costly; network does not work the same way; mobility interferes with comunication
- Limited battery: can't do too much communication
- How to make use of sensors, not so well understood

Context aware computing

- Adapt computations to the circumstances
 - Time of day
 - Is the user present?
 - Is the phone in hand or in pocket
 - Scan for wifi only when indoors
 - Turn off ring when in cinema, meeting...
 - Recognize activity and bring up relevant information

— ...

Context aware computing

- Adapt computations to the circumstances
- Basic contexts are easy to identify, but it is not always clear how to adapt
 - Turn down volume at night... but what if it is an important call?
- Many contexts are very hard to detect reliably

Context detection examples

- Detect if user is indoor or outdoor
- Use sound to detect user in a meeting
- Detect transport mode (walking, car, bus, tram..)
 - Using accelerometer
- Detect presence of other users nearby from wifi activity

Context detection

- Generally hard
- Concerns about privacy: you do not want to send context information to a server
- Perhaps distributed computation can help
 - Use data from many phones to detect context
 - But again, do not want to send all data to server
 - Do as much of it as possible on device –
 filter/process data at source

Networking in mobile systems

- Difficulty:
 - The network graph changes
 - A node is not always connected to the same router

- Example system: Mobile ad-hoc networks
 - Ad-hoc: Unplanned
 - Devices simply connect to nearby devices and route packets
 - Also applies to sensor networks

Routing in ad hoc wireless networks

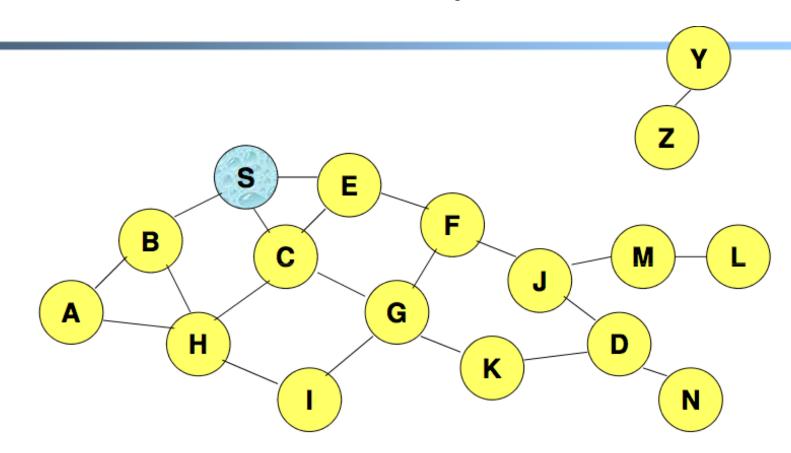
- Find route between pairs of nodes wishing to communicate.
- Proactive protocols: maintain routing tables at each node that is updated as changes in the network topology are detected.
 - Heavy overhead with high network dynamics (caused by link/node failures or node movement).
 - Not practical for networks that change frequently

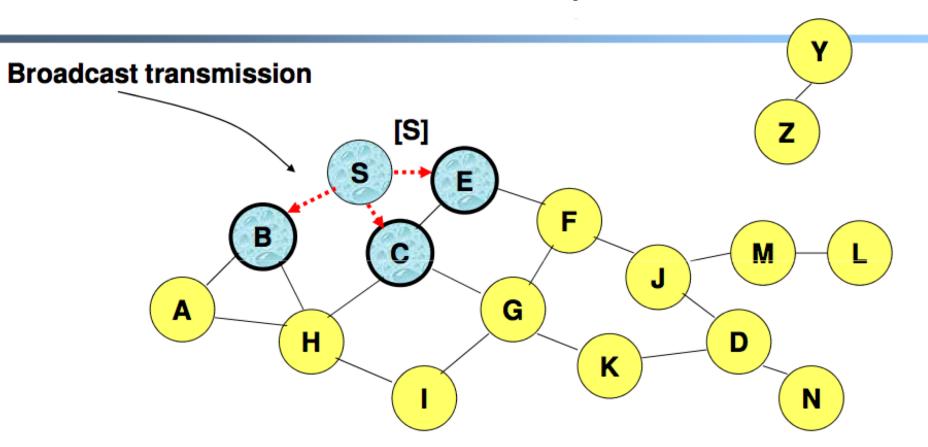
Routing in ad hoc wireless networks

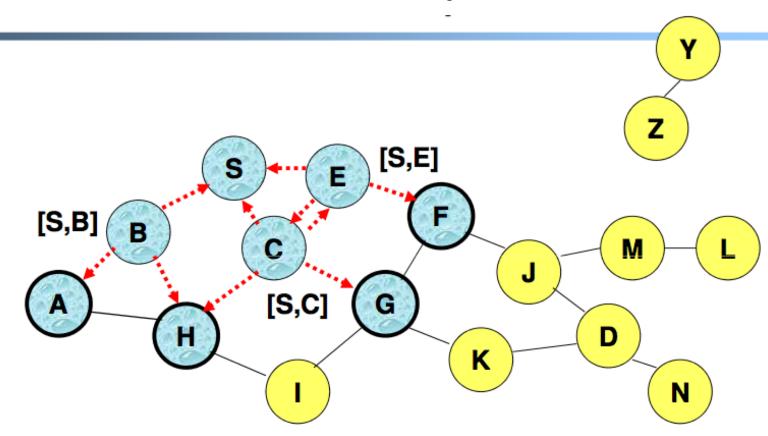
- Reactive protocols: routes are constructed on demand. No global routing table is maintained.
- More appropriate for networks with high rate of changes
 - Ad hoc on demand distance vector routing (AODV)
 - Dynamic source routing (DSR)

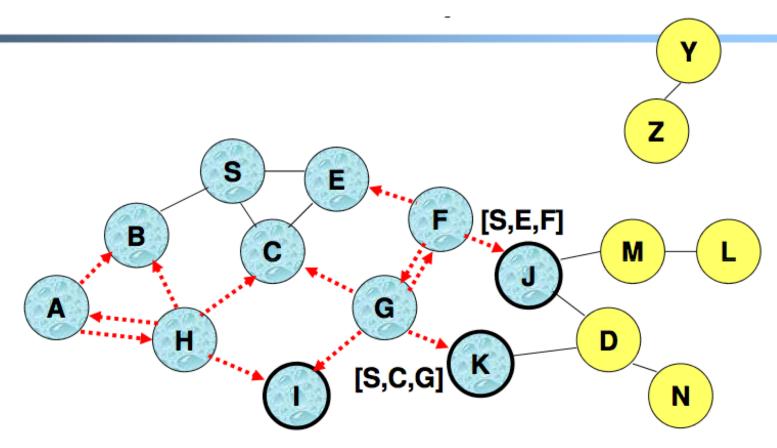
Dynamic Source Routing (DSR)

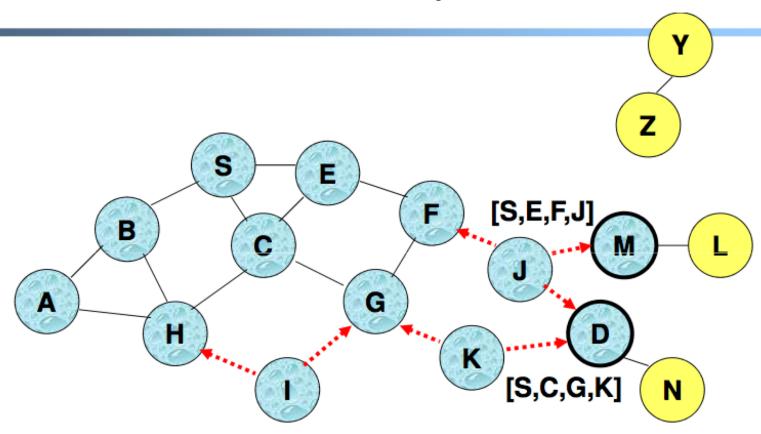
- Node S wants to send a message to node D
- S initiates a a route discovery
- S floods the network with route request (RREQ) message
- Each node appends its own id to the message

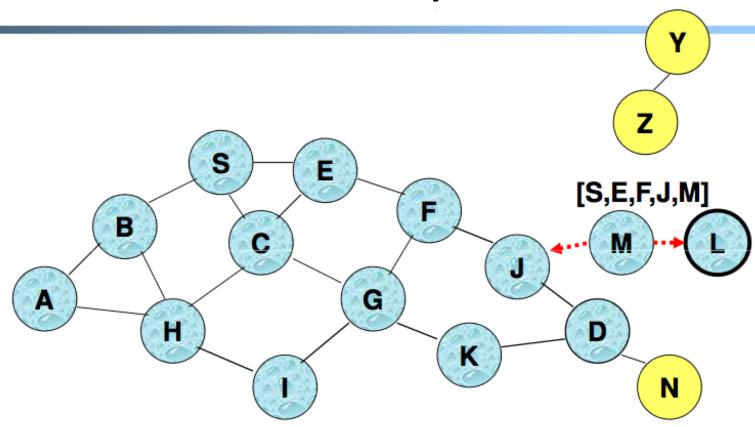






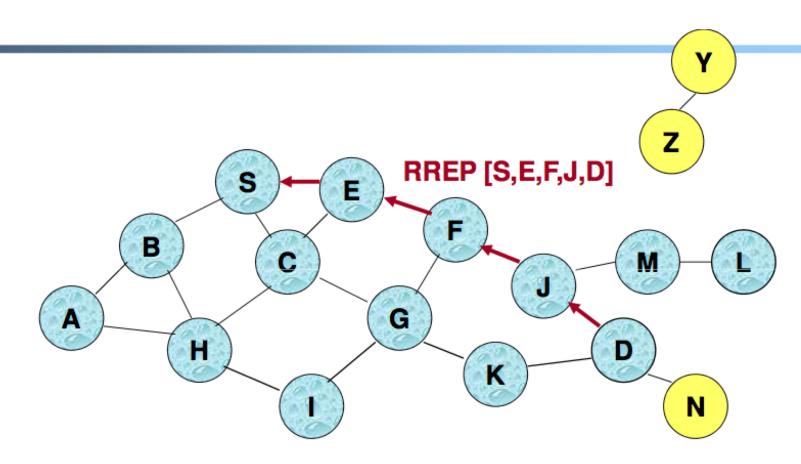


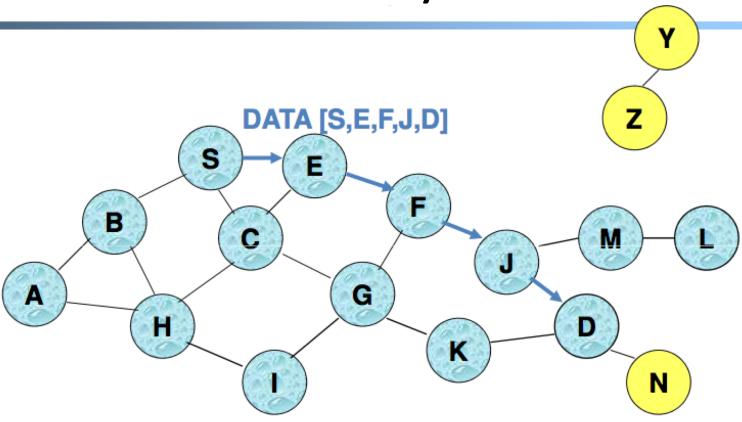




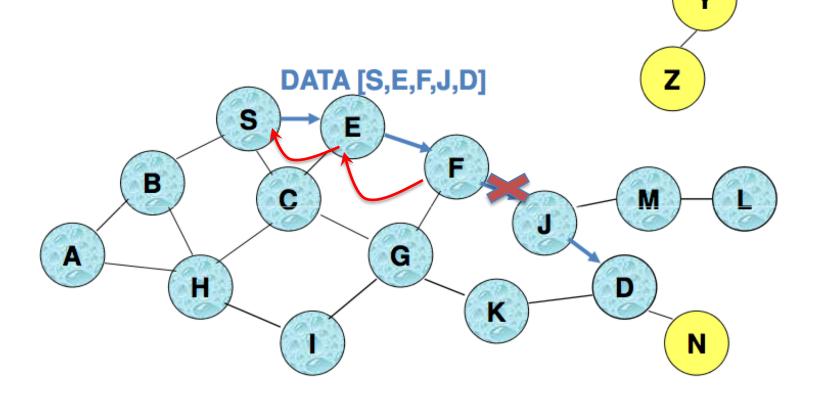
Route Discovery in DSR

- Destination D on receiving the first RREQ sends a route reply (RREP)
- RREP is sent on a route obtained by reversing the route in received RREQ





When node S sends a data packet to D, the entire route is included in the packet header, hence the name source routing



- When a link fails, an error message with the link name is sent back to S.
- S deletes any route using that link and starts discovery.

 Distributed Systems, Edinburgh, 2020

Route caching

 When a node receives or forwards a message, it learns routes to all nodes on the path

Advantage:

- S may not need to send RREQ
- Intermediate node on receiving RREQ, can respond with complete route

Disadvantage:

 Caches may be stale: S tries many cached routes before starting a discovery. Or, intermediate nodes return outdated information.

DSR: Summary

Advantages:

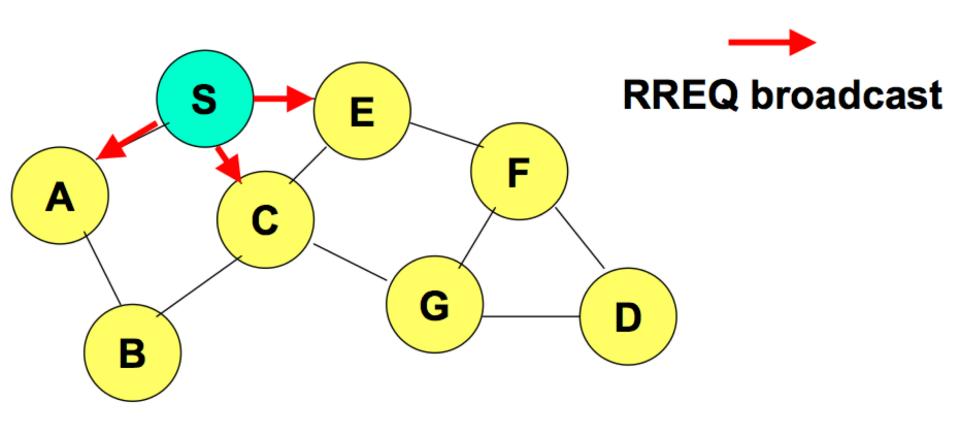
- Routes computed only when needed good for changing networks
- Caching can make things efficient
- Does not create loops

Disadvantages

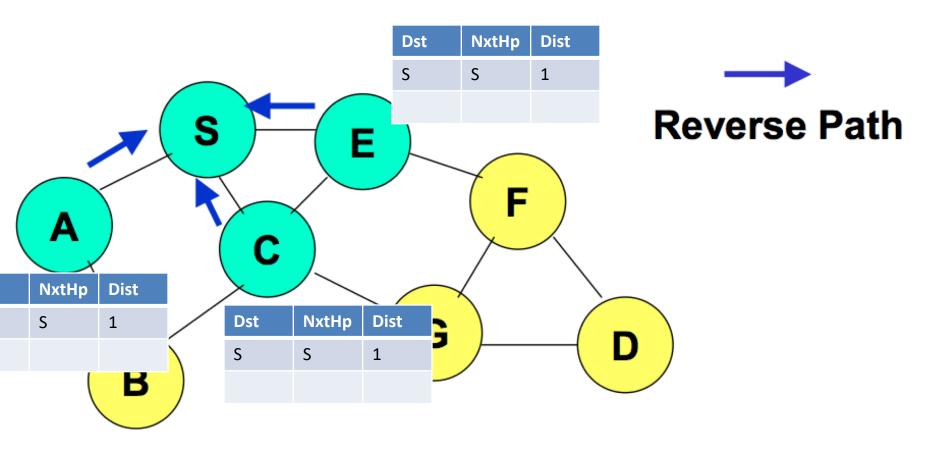
- Entire route must be contained in message: can be long for large networks
- Flooding causes communication to many nodes
- Stale caches can be a problem
- Not suitable for networks where changes are too frequent

Ad hoc On-Demand Distance Vector Routing (AODV)

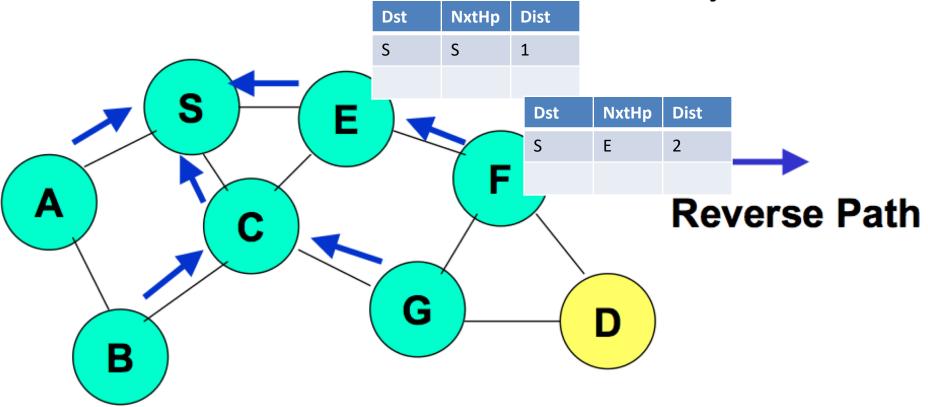
- Maintains routing tables at nodes so that the route need not be stored in the message
- No Caches: Only one route per destination



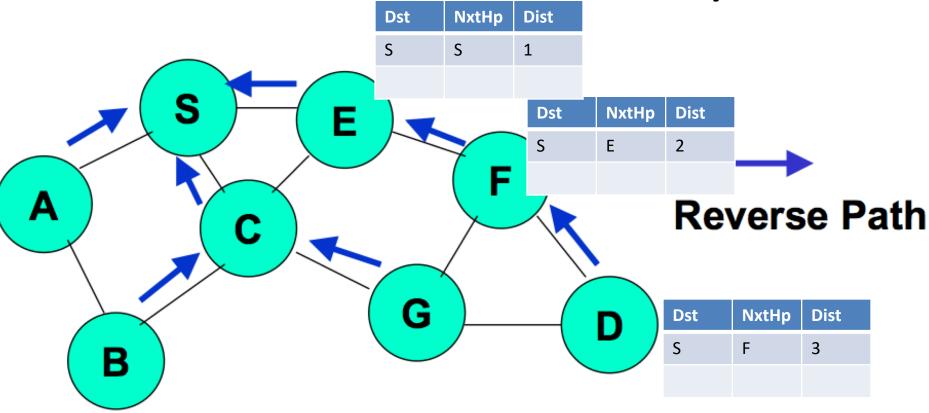
Source floods the network



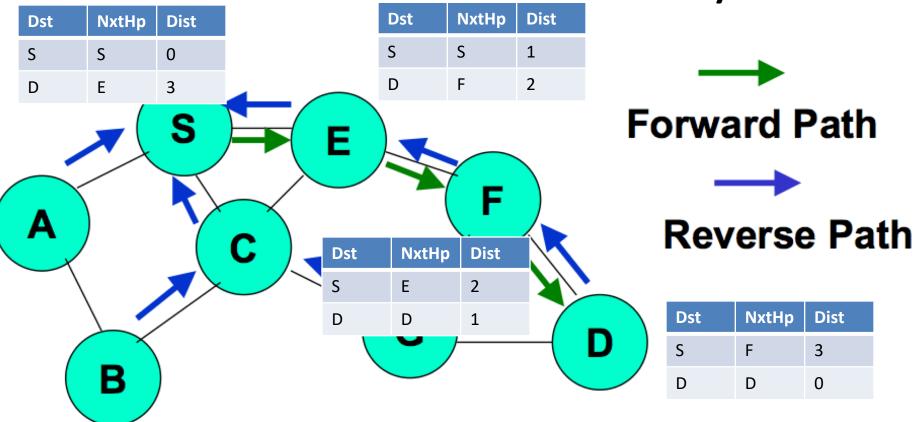
- Other nodes create parent pointer
- A node forwards a RREQ only once



- Other nodes create parent pointer
- A node forwards a RREQ only once



RREP is forwarded via reverse path

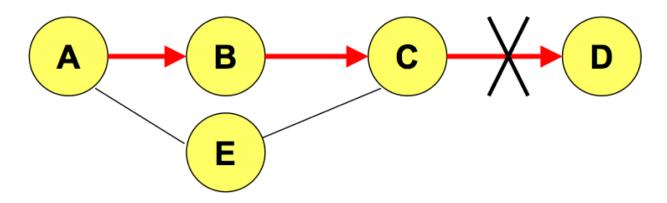


- RREP is forwarded via reverse path
- Creates a forward path, Edinburgh, 2020

Route expiry

- A path expires if not used for a certain time.
- If a node sees that a routing table entry has not been used by this time, it removes this entry
- Even if the path itself is valid
- Good for networks with frequent changes
- Bad for static and stable networks

Can create loops

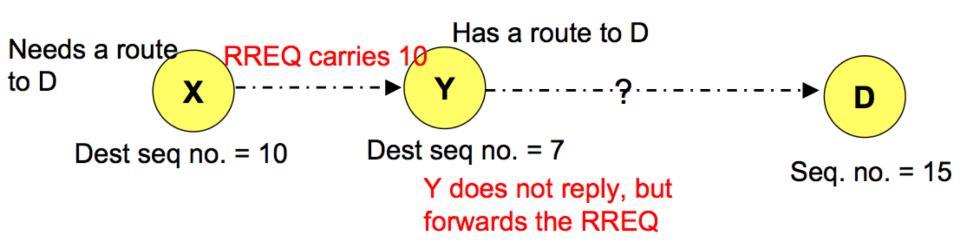


- Assume C->D link has failed, but A does not know because the ERR message was lost
- C is now trying to find path to D
- A responds since A thinks it has a path
- Creates loop: C-E-A-B-C

Sequence numbers in AODV

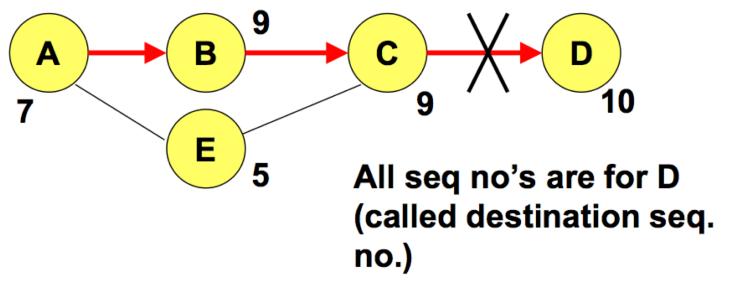
- If A has a route to D, A keeps a sequence number.
- A increments this number periodically: tells how old the information is

Using sequence numbers



Rule: sequence number must increase along any route

Sequence number rule avoids loop



 A does not reply, since its sequence no. is less than that of C

AODV

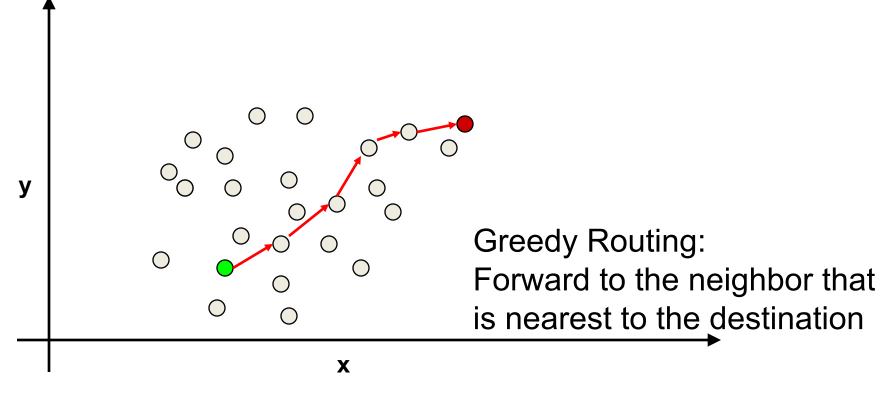
- Routing tables, message does not contain route
- Fresh routes preferred
- Old unused routes expire
- Stale routes less problematic
- Needs sequence numbers to prevent loops
- Better for more dynamic, changing environments

Routing in ad hoc networks

- Reactive protocols: routes are constructed on demand. No global routing table is maintained.
- More appropriate for networks with high rate of changes
 - Ad hoc on demand distance vector routing (AODV)
 - Dynamic source routing (DSR)
- Need flooding
 - Inefficient in large networks

Geographical routing: Using location

 Geographical routing uses a node's location to discover path to that node.



Geographical routing

Assumptions:

- Nodes know their own geographical location
- Nodes know their 1-hop neighbors
- Routing destinations are specified geographically (a location, or a geographical region)
- Each packet can hold a small amount of routing information.

Sensor network

- Sensors enabled with wireless
 - Can communicate with nearby sensors
 - Communication to server relatively costly
- Low power, but lots of data
 - Not worth sending everything to server

- Try use the data directly inside the network
 - In-network distributed computing

Problem: How to find the relevant data?

- A tourist in a park asks
- "Where is the elephant?"

• Out of all the sensors/cameras which one is close

to an elephant?



Data centric routing

- Traditional networks try to route to an IP address
- Find path to the node with a particular ID
- But what if we try to find data, not specific nodes?
- After all, delivering data is the ultimate goal of routing and networks
- Data centric storage
 - Storage depends on the data (elephant, giraffe, song...)
- Data centric routing (search)
 - Route to the data

Distributed Database

- Information Producer
 - Can be anywhere in the network
 - May be mobile
 - Many producers may generate data of the same type
- User or Information Consumer
 - Can be anywhere
 - May be many

Distributed Database: Challenges

- Consumer does not know where the producer is, and vice versa
- Need to search : Must be fast, efficient

Basic methods:

- Push: Producer disseminates data (flood)
- Pull: Consumer looks for the data (flood the query)
- Push-pull: Both producer, consumer search for each-other

Distributed hash tables

- Use a hash on the data: h(song1.mp3) = node#26
- Anyone that has song1.mp3 informs node#26
- Anyone that needs Song1.mp3 checks with node#26
- Used in peer to peer systems like Chord, pastry etc

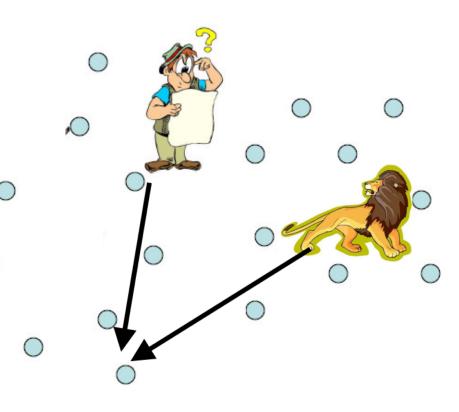
Geographic Hash Tables

Content based hash gives coordinates:

$$- h(lion) = (12, 07)$$

 Producer sends msg to (12, 07) by geographic routing and stores data

 Consumer sends msg to (12, 07) by geographic routing and gets data



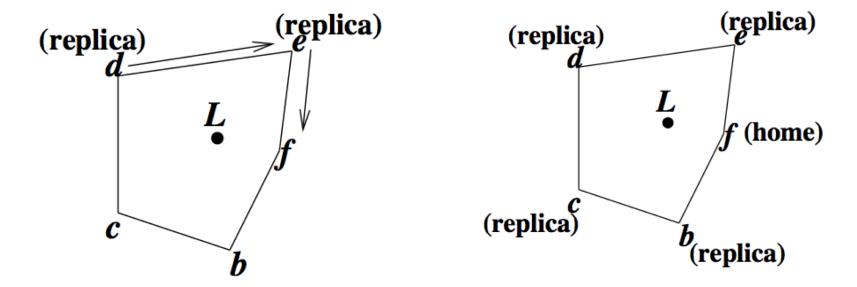
GHT

• What if there is no sensor at (12, 07)?

Use the sensor nearest to it

Fault handling

- What if home node a dies?
- Replicas have a timer that triggers a new check
- A new node becomes home



GHT

- Advantages
 - Simple
 - Handles load balancing and faults
- Disadvantages
 - Not distance sensitive: everyone has to go to hash node even if producer and consumer are close
 - If a data is queried or updated often, that node has a lot of traffic – bottleneck

Mobile, Ad-hoc and Sensor network

- A difficult model least infrastructure, low power nodes, communication/computation expensive
- Not entirely realistic
- However, it makes least number of assumptons
 - useful as a basis for developing distributed protocols/algorithms
 - Which can then be enhanced using available infrastructure in specific cases