Discrete Mathematics & Mathematical Reasoning Chapter 7 (section 7.4): Random Variables, Expectation, and Variance

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Expected Value (Expectation) of a Random Variable **Recall:** A random variable (r.v.), is a function  $X : \Omega \to \mathbb{R}$ , that assigns a real value to each outcome in a sample space  $\Omega$ .

The **expected value**, or **expectation**, or mean, of a random variable  $X : \Omega \to \mathbb{R}$ , denoted by E(X), is defined by:

$$E(X) = \sum_{s \in \Omega} P(s)X(s)$$

Here  $P: \Omega \rightarrow [0, 1]$  is the underlying probability distribution on  $\Omega$ .

**Question:** Let X be the r.v. outputing the number that comes up when a fair die is rolled. What is the expected value, E(X), of X?

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Answer:  $E(X) = \sum_{i=1}^{6} \frac{1}{6} \cdot i = \frac{21}{6} = \frac{7}{2}$ .

### A bad way to calculate expectation

The definition of expectation,  $E(X) = \sum_{s \in \Omega} P(s)X(s)$ , can be used directly to calculate E(X). But sometimes this is horribly inefficient.

**Example:** Suppose that a biased coin, which comes up heads with probability *p* each time, is flipped 11 times consecutively. **Question:** What is the expected # of heads?

### A bad way to calculate expectation

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**Example:** Suppose that a biased coin, which comes up heads with probability p each time, is flipped 11 times consecutively. **Question:** What is the expected # of heads?

**Bad way to answer this:** Let's try to use the definition of E(X) directly, with  $\Omega = \{H, T\}^{11}$ . Note that  $|\Omega| = 2^{11} = 2048$ . So, the sum  $\sum_{s \in \Omega} P(s)X(s)$  has 2048 terms!

This is clearly not a practical way to compute E(X). Is there a better way? Yes.

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### Better expression for the expectation

Recall P(X = r) denotes the probability  $P(\{s \in \Omega \mid X(s) = r\})$ . Recall that for a function  $X : \Omega \to \mathbb{R}$ ,

 $range(X) = \{r \in \mathbb{R} \mid \exists s \in \Omega \text{ such that } X(s) = r\}$ 

**Theorem:** For a random variable  $X : \Omega \to \mathbb{R}$ ,

$$E(X) = \sum_{r \in range(X)} P(X = r) \cdot r$$

**Proof:**  $E(X) = \sum_{s \in \Omega} P(s)X(s)$ , but for each  $r \in range(X)$ , if we sum all terms P(s)X(s) such that X(s) = r, we get  $P(X = r) \cdot r$  as their sum. So, summing over all  $r \in range(X)$  we get  $E(X) = \sum_{r \in range(X)} P(X = r) \cdot r$ .

So, if |range(X)| is small, and if we can compute P(X = r), then we need to sum a lot fewer terms to calculate E(X).

# Expected # of successes in *n* Bernoulli trials

**Theorem:** The expected # of successes in *n* (independent) Bernoulli trials, with probability *p* of success in each, is *np*.

Note: We'll see later that we do not need independence for this. **First, a proof which uses mutual independence:** For  $\Omega = \{H, T\}^n$ , let  $X : \Omega \to \mathbb{N}$  count the number of successes in *n* Bernoulli trials. Let q = (1 - p). Then...

$$E(X) = \sum_{k=0}^{n} P(X = k) \cdot k$$
$$= \sum_{k=1}^{n} {n \choose k} p^{k} q^{n-k} \cdot k$$

The second equality holds because, assuming mutual independence, P(X = k) is the binomial distribution b(k; n, p).

# first proof continued

$$E(X) = \sum_{k=0}^{n} P(X = k) \cdot k = \sum_{k=1}^{n} {n \choose k} p^{k} q^{n-k} \cdot k =$$

$$= \sum_{k=1}^{n} \frac{n!}{k!(n-k)!} p^{k} q^{n-k} \cdot k = \sum_{k=1}^{n} \frac{n!}{(k-1)!(n-k)!} p^{k} q^{n-k}$$

$$= \sum_{k=1}^{n} n \cdot \frac{(n-1)!}{(k-1)!(n-k)!} p^{k} q^{n-k} = n \sum_{k=1}^{n} {n-1 \choose k-1} p^{k} q^{n-k}$$

$$= np \sum_{k=1}^{n} {n-1 \choose k-1} p^{k-1} q^{n-k} = np \sum_{j=0}^{n-1} {n-1 \choose j} p^{j} q^{n-1-j}$$

$$= np(p+q)^{n-1}$$

= np.  $\Box$ 

We will soon see this was an unnecessarily complicated proof.

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# Expectation of a geometrically distributed r.v.

**Question:** A coin comes up heads with probability p > 0 each time it is flipped. The coin is flipped repeatedly until it comes up heads. What is the expected number of times it is flipped?

**Note:** This simply asks: "What is the expected value E(X) of a geometrically distributed random variable with parameter p?"

**Answer:**  $\Omega = \{H, TH, TTH, ...\}$ , and  $P(T^{k-1}H) = (1-p)^{k-1}p$ . And clearly  $X(T^{k-1}H) = k$ . Thus  $E(X) = \sum_{s \in \Omega} P(s)X(s) =$ 

$$E(X) = \sum_{k=1}^{\infty} (1-p)^{k-1} p \cdot k = p \sum_{k=1}^{\infty} k (1-p)^{k-1} = p \cdot \frac{1}{p^2} = \frac{1}{p}.$$

This is because:  $\sum_{k=1}^{\infty} k \cdot x^{k-1} = \frac{1}{(1-x)^2}$ , for |x| < 1. **Example:** If p = 1/4, then the expected number of coin tosses before we see Heads for the first time is 4.

# Linearity of Expectation (VERY IMPORTANT)

**Theorem (Linearity of Expectation):** For any random variables  $X, X_1, \ldots, X_n$  on  $\Omega$ ,  $E(X_1 + X_2 + \ldots + X_n) = E(X_1) + \ldots + E(X_n)$ .

Furthermore, for any  $a, b \in \mathbb{R}$ ,

$$E(aX+b)=aE(X)+b.$$

(In other words, the expectation function is a linear function.)

#### Proof:

$$E(\sum_{i=1}^{n} X_i) = \sum_{s \in \Omega} P(s) \sum_{i=1}^{n} X_i(s) = \sum_{i=1}^{n} \sum_{s \in \Omega} P(s) X_i(s) = \sum_{i=1}^{n} E(X_i).$$
$$E(aX+b) = \sum_{s \in \Omega} P(s)(aX(s)+b) = (a \sum_{s \in \Omega} P(s)X(s)) + b \sum_{s \in \Omega} P(s)$$
$$= aE(X) + b. \quad \Box$$

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# Using linearity of expectation

**Theorem:** The expected # of successes in n (not necessarily independent) Bernoulli trials, with probability p of success in each trial, is np.

**Easy proof, via linearity of expectation:** For  $\Omega = \{H, T\}^n$ , let *X* be the r.v. counting the number of successes, and for each *i*, let  $X_i : \Omega \to \mathbb{R}$  be the binary r.v. defined by:

$$X_i((s_1,\ldots,s_n)) = \left\{ egin{array}{cc} 1 & ext{if } s_i = H \ 0 & ext{if } s_i = T \end{array} 
ight.$$

Note that  $E(X_i) = p \cdot 1 + (1 - p) \cdot 0 = p$ , for all  $i \in \{1, ..., n\}$ . Also, clearly,  $X = X_1 + X_2 + ... + X_n$ , so:

$$E(X) = E(X_1 + \ldots + X_n) = \sum_{i=1}^n E(X_i) = np.$$

Note: this holds even if the *n* coin tosses are totally correlated.

# Using linearity of expectation, continued

**Hatcheck problem:** At a restaurant, the hat-check person forgets to put claim numbers on hats.

*n* customers check their hats in, and they each get a random hat back when they leave the restuarant.

What is the expected number, E(X), of people who get their correct hat back?

**Answer:** Let  $X_i$  be the r.v. that is 1 if the *i*'th customer gets their hat back, and 0 otherwise.

Clearly,  $E(X) = E(\sum_i X_i)$ . Furthermore,  $E(X_i) = P(i$ 'th person gets its hat back) = 1/n. Thus,  $E(X) = n \cdot (1/n) = 1$ .

This would be much harder to prove without using the linearity of expectation.

Note: E(X) doesn't even depend on *n* in this case.

## Independence of Random Variables

**Definition:** Two random variables, *X* and *Y*, are called **independent** if for all  $r_1, r_2 \in \mathbb{R}$ :

$$P(X = r_1 \text{ and } Y = r_2) = P(X = r_1) \cdot P(Y = r_2)$$

**Example:** Two die are rolled. Let  $X_1$  be the number that comes up on die 1, and let  $X_2$  be the number that comes up on die 2. Then  $X_1$  and  $X_2$  are independent r.v.'s.

**Theorem:** If *X* and *Y* are independent random variables on the same space  $\Omega$ . Then

$$E(XY) = E(X)E(Y)$$

We will not prove this in class. (The proof is a simple re-arrangement of the sums in the definition of expectation. See Rosen's book for a proof.)

### Variance

The "variance" and "standard deviation" of a r.v., *X*, give us ways to measure (roughly) *"on average, how far off the value of the r.v. is from its expectation"*.

### Variance and Standard Deviation

**Definition:** For a random variable *X* on a sample space  $\Omega$ , the **variance** of *X*, denoted by V(X), is defined by:

$$V(X) = E((X - E(X))^2) = \sum_{s \in \Omega} (X(s) - E(X))^2 P(s)$$

The standard deviation of X, denoted  $\sigma(X)$ , is defined by

$$\sigma(\boldsymbol{X}) = \sqrt{V(\boldsymbol{X})}$$

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### Example, and a useful identity for variance

**Example:** Consider the r.v., X, such that P(X = 0) = 1, and the r.v. Y, such that P(Y = -10) = P(Y = 10) = 1/2. Then E(X) = E(Y) = 0, but  $V(X) = 0 = \sigma(X)$ , whereas V(Y) = 100 and  $\sigma(Y) = 10$ .

**Theorem:** For any random variable *X*,

$$V(X) = E(X^2) - E(X)^2$$

#### Proof:

$$V(X) = E((X - E(X))^{2})$$
  
=  $E(X^{2} - 2XE(X) + E(X)^{2})$   
=  $E(X^{2}) - 2E(X)E(X) + E(X)^{2}$   
=  $E(X^{2}) - E(X)^{2}$ .  $\Box$