

Transition rules (including axioms)

Communication and Concurrency Lecture 3

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$$\begin{aligned}
 & R(.) \quad a.E \xrightarrow{a} E \\
 & R(in) \quad a(x).E \xrightarrow{a(v)} E\{v/x\} \quad \text{if } v \in D \\
 & R(out) \quad \bar{a}(e).E \xrightarrow{\bar{a}(v)} E \quad \text{if } Val(e) = v \\
 & R(\stackrel{\text{def}}{=}) \quad \frac{P \xrightarrow{a} F}{E \xrightarrow{a} F} \quad P \stackrel{\text{def}}{=} E \\
 & R(+) \quad \frac{E_1 + E_2 \xrightarrow{a} F}{E_1 \xrightarrow{a} F} \quad \frac{E_1 + E_2 \xrightarrow{a} F}{E_2 \xrightarrow{a} F} \\
 & R(| \text{ com}) \quad \frac{E | F \xrightarrow{\tau} E' | F'}{E \xrightarrow{a} E' \quad F \xrightarrow{\bar{a}} F'} \\
 & R(|) \quad \frac{E | F \xrightarrow{a} E' | F}{E \xrightarrow{a} E'} \quad \frac{E | F \xrightarrow{a} E | F'}{F \xrightarrow{a} F'} \\
 & \quad \frac{E \setminus J \xrightarrow{a} F \setminus J}{E \xrightarrow{a} F} \quad a \notin J \cup \bar{J}
 \end{aligned}$$

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Example: protocol that may lose messages

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Sender     $\stackrel{\text{def}}{=} \text{in}(x).\overline{\text{sm}}(x).\text{Send1}(x)$ 
Send1(x)   $\stackrel{\text{def}}{=} \text{ms}.\overline{\text{sm}}(x).\text{Send1}(x) + \text{ok}.\text{Sender}$ 
Medium     $\stackrel{\text{def}}{=} \text{sm}(y).\text{Med1}(y)$ 
Med1(y)    $\stackrel{\text{def}}{=} \overline{\text{mr}}(y).\text{Medium} + \tau.\overline{\text{ms}}.\text{Medium}$ 
Receiver   $\stackrel{\text{def}}{=} \text{mr}(x).\overline{\text{out}}(x).\text{ok}.\text{Receiver}$ 

Protocol   $\equiv (\text{Sender} \mid \text{Medium} \mid \text{Receiver}) \setminus \{\text{sm}, \text{ms}, \text{mr}, \text{ok}\}$ 
    
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Abstracting from silent activity

► Difference between τ and “observable” actions

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Abstracting from silent activity

- ▶ Difference between τ and “observable” actions
- ▶ Suppose E may at some time perform ok.

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Abstracting from silent activity

- ▶ Difference between τ and “observable” actions
- ▶ Suppose E may at some time perform ok.
- ▶ $\text{In } (E \mid \overline{\text{ok}}.\text{Resource}) \setminus \{\text{ok}\}$
access to Resource is triggered by ok by E
- ▶ Observation of ok = release of Resource

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- ▶ $\text{In } (E \mid \overline{\text{ok}}.\text{Resource}) \setminus \{\text{ok}\}$
access to Resource is triggered by ok by E
- ▶ Observation of ok = release of Resource
- ▶ τ cannot be observed in this way

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Observable transitions

$$\begin{aligned} C &\stackrel{\text{def}}{=} \text{in}(x).\overline{\text{out}}(x).\overline{\text{ok}}.C \\ U &\stackrel{\text{def}}{=} \text{write}(x).\overline{\text{in}}(x).\text{ok}.U \end{aligned}$$

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What is difference between?

$$(C \mid U) \setminus \{\text{in}, \text{ok}\}$$

$$U_{\text{cop}} \stackrel{\text{def}}{=} \text{write}(x).\overline{\text{out}}(x).U_{\text{cop}}$$

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Exercise: observable transition graphs

1. Assuming just one datum value, draw the observable graphs for processes $(C \mid U) \setminus \{\text{in}, \text{ok}\}$ and U_{cop}

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$$E \xRightarrow{\varepsilon} F \text{ or } E \xRightarrow{a} F \text{ where } a \neq \tau$$

$$R(\xRightarrow{\varepsilon}) \quad \frac{E \xRightarrow{\varepsilon} E \quad \frac{E \xRightarrow{\varepsilon} F}{E \xrightarrow{\tau} E'} \quad E' \xRightarrow{\varepsilon} F}{E \xRightarrow{\varepsilon} F}$$

$$R(\xRightarrow{a}) \quad \frac{E \xRightarrow{a} F}{\frac{E \xRightarrow{\varepsilon} E' \quad E' \xrightarrow{a} F' \quad F' \xRightarrow{\varepsilon} F}{E \xRightarrow{a} F}}$$

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2. Draw both kinds of transition graph for the following pair of processes, $\tau.0$ and $Div' \stackrel{\text{def}}{=} \tau.Div' + \tau.0$



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3. Assuming just one datum value, draw the observable graph for $(Cop \mid User) \setminus \{in\}$



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2. Draw both kinds of transition graph for the following pair of processes, $\tau.0$ and $Div' \stackrel{\text{def}}{=} \tau.Div' + \tau.0$
3. Assuming just one datum value, draw the observable graph for $(Cop \mid User) \setminus \{in\}$
4. Draw the observable graph for Peterson

