Advances in Programming Languages APL2: Types and type systems

Ian Stark

School of Informatics The University of Edinburgh

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http://www.inf.ed.ac.uk/teaching/courses/apl/

Some types

A selection of types from some languages.

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C/C++
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int, long, float, unsigned int, char int [], char*, char&, int(*)(float,char) OCaml

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int, int64, bool, char, string, unit
string*string, int list, bool array
int->int, int->string->char, 'a list -> 'a list
```

Java

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Object, byte[], boolean
StringBuffer, LinkedList, TreeSet, ArrayList<String>
IllegalPathStateException, BeanContextServiceRevokedListener
```

- Type checking
- Static type checking
- Dynamic type checking
- Type annotation
- Type inference

- Subtyping
- Structural typing
- Nominative typing
- Duck typing
- Effect types

A *type system* is a syntactically defined subset T of programs such that:

 $\mathsf{P} \in \mathsf{T} \quad \Longrightarrow \quad \mathsf{Compile}(\mathsf{P}) \models \varphi$

(read: "if P is in T then Compile(P) satisfies ϕ ")

where Compile(P) is the object code corresponding to P and ϕ is some desired property of its execution.

For example,

 $\mathsf{T}=\text{``well-typed Java programs''}$

 $\phi =$ "methods are always correctly invoked"

Slogan: Well-typed programs cannot go wrong. [Robin Milner, 1978]

Java is serious about abstraction

Java works almost entirely through class-based object-oriented programming; it encourages the use of abstract classes through inheritance and interfaces; and it does not expose the private workings of classes and packages.

Java is serious about typing

Java has strong static typing: all programs are checked for type-correctness at compile-time. Bytecode is checked again when classes are loaded, by the *bytecode verifier*, before execution. The recent introduction of *generics* extends the power of the type system.

Even so, things do not always go as well as one might hope...

Java has subtyping: a value of one type may be used at any more general type. So String \leq Object, and every String is an Object.

Not all is well with Java types

String[] $a = \{ "Hello" \};$	// A small string array
Object[] b = a;	// Now a and b are the same array
b[0] = Boolean.FALSE;	// Drop in a Boolean object
String $s = a[0];$	// Oh, dear
System.out.println(s.toUpperCase());	// This isn't going to be pretty

This compiles without error or warning: in Java, if $S \leq T$ then $S[] \leq T[]$. Except that it isn't. So every array assignment gets a runtime check.

Subtype variance

The issue here is with *parameterized types* like String[] and List<Object>; or in OCaml ('a list -> 'a list) and ('a * 'b).

Suppose some type A<X> depends on type X, and types S \leqslant T. Then the dependency is:

For example, in the Scala language, type parameters can be annotated with variance information: List[+T], Function[-S,+T].

In Java, arrays are typed as if they were covariant. But they aren't. We shall revisit this later. . .

By the next lecture, on Monday:

- Test out the Java array subtyping example, and confirm that (a) it compiles, and (b) there is a type error when run.
- Read the Java fable Execution in the Kingdom of Nouns.

If you are uncertain about OCaml programming, try these online guides:

- Chapter 1 of OCaml for Scientists
- The Objective Caml Tutorial
- Developing Applications with Objective Caml
- For those who already know Standard ML, Andreas Rossberg has written a handy conversion guide.

- Languages use types and type systems for several reasons.
- A *type system* is a syntactically defined subset of programs which are certain to have some desired property.
- Java has covariance subtyping of arrays, which can cause runtime type errors.