Algorithms and Data Structures 2023/24 Week 5 tutorial sheet

1. In class we mostly worked with DFT for the case where n is a power of 2, and the polynomial being evaluated has degree n - 1. However, we also showed that we can apply the DFT and Inverse DFT when n is not a power of 2, by taking n' to be the closest power of 2 satisfying $n \le n'$, and adding some leading coefficients of value 0.

Make this idea formal: First show how to compute n'. Also show that our DFT or Inverse DFT still takes $\Theta(n \lg(n))$ time in terms of the original value n.

- 2. This exercise asks you to do a few complex number calculations. Evaluate each of these. (See also rules for multiplication and division in the FFT notes.)
 - (a) 2i(3-i).
 - (b) $2i(i+1)^2 + 4(i+1)^3$.
 - (c) 3i/(1+i).
- Compute DFT₄⟨0,1,2,3⟩. (do this directly, rather than by FFT, if you prefer). This is Exercise 30.2-2, p. 838 of [CLRS].
- 4. Use the FFT to efficiently multiply the two polynomials p(x) = x-4 and $q(x) = x^2-1$. Use the following steps:
 - (a) First work out what will be the degree of the product polynomial pq. Take deg(pq) + 1 as our n (and if necessary round up to the nearest power of 2).
 - (b) For this value of n (which we made sure was a power of 2), use trigonometry to write down each of the nth roots-of-unity (so we have them to work with).
 - (c) Calculate the DFT for p(x) for nth roots of unity.
 - (d) Calculate the DFT for q(x) for nth roots of unity.
 - (e) Do pointwise multiplication of the two DFTs to get the DFT of pq(x) for nth roots of unity.
 - (f) Calculate the Inverse DFT of the DFT for pq(x), to obtain the polynomial pq. It's a good idea to do this via DFT (and then swapping), like we saw in class.
 - (g) Check your answer by straight multiplication.