From last time...
A decision function for a relation $R \subseteq U$ is a total function $p$ from $U$ to $\{true, false\}$ such that $p(u) = true$ iff $u \in R$. 
A decision function for a relation $R \subseteq U$ is a total function $p$ from $U$ to $\{true, false\}$ such that $p(u) = true$ iff $u \in R$.

Example:

$$U = \{1, 2, 3\}$$
$$R = \{(1, 2), (2, 3)\}$$

Note that, for now, we are saying absolutely nothing about computability. We’ll come back to this in a moment.
A decision function for a relation $R \subseteq U$ is a total function $p$ from $U$ to \{true, false\} such that $p(u) = true$ iff $u \in R$.

Example:

\[
U = \{1, 2, 3\} \\
R = \{(1, 2), (2, 3)\}
\]

The function $p$ whose graph is

\[
\{(1, 2, true), (2, 3, true), \\
(1, 1, false), (1, 3, false), \\
(2, 1, false), (2, 2, false), \\
(3, 1, false), (3, 2, false), (3, 3, false)\}
\]

is a decision function for $R$.  

A decision function for a relation $R \subseteq U$ is a total function $p$ from $U$ to $\{true, false\}$ such that $p(u) = true$ iff $u \in R$.

Example:

$$U = \{1, 2, 3\}$$
$$R = \{(1, 2), (2, 3)\}$$

The function $p'$ whose graph is

$$\{((1, 2), true), ((2, 3), true)\}$$

is not a decision function for $R$. 
A decision function for a relation $R \subseteq U$ is a total function $p$ from $U$ to \{true, false\} such that $p(u) = true$ iff $u \in R$.

Example:

\[
U = \{1, 2, 3\} \\
R = \{(1, 2), (2, 3)\}
\]

The function $p''$ whose graph is

\[
\{((1, 2), true), ((2, 3), true), ((1, 3), false)\}
\]

is also not a decision function for $R$. 

Decision Procedures (take 1)
Decision Procedures (take 2)

Of course, we want a decision procedure to be a *procedure*.

A *decision procedure* for a relation $R \subseteq U$ is a *computable* total function $p$ from $U$ to $\{true, false\}$ such that $p(u) = true$ iff $u \in R$. 
Example

\[
U = \{1, 2, 3\} \\
R = \{(1, 2), (2, 3)\}
\]
Example

\[ U = \{1, 2, 3\} \]
\[ R = \{(1, 2), (2, 3)\} \]

The function

\[ p(x, y) = \begin{cases} 
  \text{true} & \text{if } x = 2 \text{ and } y = 3 \\
  \text{true} & \text{else if } x = 1 \text{ and } y = 2 \\
  \text{false} & \text{else}
\end{cases} \]

whose graph is

\{ ((1, 2), true), ((2, 3), true), 
  ((1, 1), false), ((1, 3), false), 
  ((2, 1), false), ((2, 2), false), 
  ((3, 1), false), ((3, 2), false), ((3, 3), false) \}

is a decision procedure for \( R \).
Example

\[ U = \{1, 2, 3\} \]
\[ R = \{(1, 2), (2, 3)\} \]

The recursively defined partial function

\[ p(x, y) = \begin{cases} 
  \text{true} & \text{if } x = 2 \text{ and } y = 3 \\
  \text{true} & \text{if } x = 1 \text{ and } y = 2 \\
  \text{false} & \text{if } x = 1 \text{ and } y = 3 \\
  p(x, y) & \text{else}
\end{cases} \]

whose graph is

\[ \{(1, 2), \text{true}, (2, 3), \text{true}, (1, 3), \text{false}\} \]

is not a decision procedure for \( R \).
Subtyping Algorithm

This recursively defined total function is a decision procedure for the subtype relation:

\[
\text{subtype}(S, T) =
\]

- if \( T = \text{Top} \), then \( \text{true} \)
- else if \( S = S_1 \rightarrow S_2 \) and \( T = T_1 \rightarrow T_2 \)
  - then \( \text{subtype}(T_1, S_1) \land \text{subtype}(S_2, T_2) \)
- else if \( S = \{ k_j : S_j \mid j \in 1..m \} \) and \( T = \{ l_i : T_i \mid i \in 1..n \} \)
  - then \( \{ l_i \mid i \in 1..n \} \subseteq \{ k_j \mid j \in 1..m \} \)
    - and for all \( i \in 1..n \) there is some \( j \in 1..m \) with \( k_j = l_i \)
    - and \( \text{subtype}(S_j, T_i) \)
- else \( \text{false} \).

To show this, we need to prove:

1. that it returns \( \text{true} \) whenever \( S <: T \), and
2. that it returns either \( \text{true} \) or \( \text{false} \) on all inputs.
Subtyping Algorithm

But this recursively defined *partial* function is not:

\[
\text{subtype}(S, T) = \\
\text{if } T = \text{Top}, \text{ then } \text{true} \\
\text{else if } S = S_1 \rightarrow S_2 \text{ and } T = T_1 \rightarrow T_2 \\
\quad \text{ then } \text{subtype}(T_1, S_1) \land \text{subtype}(S_2, T_2) \\
\text{else if } S = \{k_j : S_j \mid j \in 1..m\} \text{ and } T = \{l_i : T_i \mid i \in 1..n\} \\
\quad \text{ then } \{l_i \mid i \in 1..n\} \subseteq \{k_j \mid j \in 1..m\} \\
\quad \text{ and for all } i \in 1..n \text{ there is some } j \in 1..m \text{ with } k_j = l_i \\
\quad \text{ and subtype}(S_j, T_i) \\
\text{else subtype}(T, S)
\]
Algorithmic Typing
Algorithmic typing

- How do we implement a type checker for the lambda-calculus with subtyping?
- Given a context $\Gamma$ and a term $t$, how do we determine its type $T$, such that $\Gamma \vdash t : T$?
**Issue**

For the typing relation, we have just one problematic rule to deal with: subsumption.

\[
\frac{\Gamma \vdash t : S \quad S <: T}{\Gamma \vdash t : T} \quad (T\text{-}SUB)
\]

We observed above that this rule is sometimes *required* when typechecking applications:

E.g., the term

\[
(\lambda r:\{x:\text{Nat}\}. \ r.x) \ {x=0,y=1}
\]

is not typable without using subsumption.

But we *conjectured* that applications were the only critical uses of subsumption.
Plan

1. Investigate how subsumption is used in typing derivations by looking at examples of how it can be “pushed through” other rules

2. Use the intuitions gained from this exercise to design a new, algorithmic typing relation that
   - omits subsumption
   - compensates for its absence by enriching the application rule

3. Show that the algorithmic typing relation is essentially equivalent to the original, declarative one
Example (T-\text{SUB} with T-\text{ABS})

\[
\begin{array}{c}
\vdots \\
\hline
\Gamma, x:S_1 \vdash s_2 : S_2 \\
\hline
\end{array}
\]

\[
\begin{array}{c}
\vdots \\
\hline
S_2 <: T_2 \\
\hline
\Gamma, x:S_1 \vdash s_2 : T_2 \\
\hline
\end{array}
\]

\[
\begin{array}{c}
\Gamma \vdash \lambda x:S_1. s_2 : S_1 \rightarrow T_2 \\
\hline
\end{array}
\]

(T-\text{SUB})

(T-\text{ABS})
Example \((T\text{-}Sub \text{ with } T\text{-}Abs)\)

\[
\begin{align*}
\Gamma, x : S_1 &\vdash s_2 : S_2 \\
S_2 &<: T_2 \\
\hline
\Gamma, x : S_1 &\vdash s_2 : T_2 \\
\hline
\Gamma &\vdash \lambda x : S_1 . s_2 : S_1 \to T_2
\end{align*}
\]

becomes

\[
\begin{align*}
\Gamma, x : S_1 &\vdash s_2 : S_2 \\
\hline
\Gamma &\vdash \lambda x : S_1 . s_2 : S_1 \to T_2
\end{align*}
\]

\[
\Gamma &\vdash \lambda x : S_1 . s_2 : S_1 \to S_2 \\
\hline
\Gamma &\vdash \lambda x : S_1 . s_2 : S_1 \to T_2
\]

\[
\begin{align*}
\Gamma, x : S_1 &\vdash s_2 : S_2 \\
S_1 &<: S_1 \\
\hline
S_1 \to S_2 &<: S_1 \to T_2 \\
\hline
\Gamma &\vdash \lambda x : S_1 . s_2 : S_1 \to T_2
\end{align*}
\]
Example \((T\text{-}SUB \text{ with } T\text{-}RCD)\)

\[
\begin{align*}
\vdash t_i : S_i & \quad S_i <: T_i \\
\text{for each } i & \quad \vdash t_i : T_i \\
\vdash \{ l_i = t_i \}_{i \in 1..n} & : \{ l_i : T_i \}_{i \in 1..n}
\end{align*}
\]
Intuitions

These examples show that we do not need $T\text{-}Sub$ to “enable” $T\text{-}Abs$ or $T\text{-}Rcd$: given any typing derivation, we can construct a derivation with the same conclusion in which $T\text{-}Sub$ is never used immediately before $T\text{-}Abs$ or $T\text{-}Rcd$.

What about $T\text{-}App$?
We’ve already observed that $T\text{-}Sub$ is required for typechecking some applications. So we expect to find that we cannot play the same game with $T\text{-}App$ as we’ve done with $T\text{-}Abs$ and $T\text{-}Rcd$. Let’s see why.
Example \((T\text{-}SUB \text{ with } T\text{-}APP \text{ on the left})\)

\[
\begin{array}{ll}
\Gamma \vdash s_1 : S_{11} \rightarrow S_{12} & \quad \vdash \quad \vdash \\
\Gamma \vdash s_2 : S_{11} \rightarrow S_{12} & \quad \vdash \quad \vdash \\
\Gamma \vdash s_1 : T_{11} \rightarrow T_{12} & \quad \vdash \quad \vdash \\
\Gamma \vdash s_2 : T_{11} & \quad \vdash \quad \vdash \\
\end{array}
\]

\((S\text{-ARROW})\)

\[
\begin{array}{ll}
T_{11} \ll S_{11} & \quad S_{12} \ll T_{12} \\
S_{11} \rightarrow S_{12} \ll T_{11} \rightarrow T_{12} & \quad \vdash \quad \vdash \\
T_{11} \ll \quad \quad \quad \quad \vdash \quad \vdash \\
T_{11} \rightarrow T_{12} & \quad \vdash \quad \vdash \\
\end{array}
\]

\((T\text{-SUB})\)

\[
\begin{array}{ll}
\Gamma \vdash s_1 \quad s_2 : T_{12} & \quad \vdash \quad \vdash \\
\end{array}
\]

\((T\text{-APP})\)
Example \((T\text{-}\text{SUB} \text{ with } T\text{-}\text{APP} \text{ on the left})\)

\[
\begin{align*}
\Gamma \vdash s_1 : S_{11} \rightarrow S_{12} & \quad & \vdash s_1 : T_{11} \rightarrow T_{12} \\
S_{11} \rightarrow S_{12} \ll : T_{11} \rightarrow T_{12} & \quad & \vdash s_2 : T_{11} \\
\Gamma \vdash s_1, s_2 : T_{12}
\end{align*}
\]

becomes

\[
\begin{align*}
\Gamma \vdash s_1 : S_{11} \rightarrow S_{12} & \quad & \vdash s_1 : T_{11} \rightarrow T_{12} \\
\Gamma \vdash s_2 : T_{11} & \quad & \vdash s_2 : S_{11} \\
\Gamma \vdash s_1, s_2 : S_{12} & \quad & \Gamma \vdash s_2 : S_{11} \\
\Gamma \vdash s_1, s_2 : T_{12}
\end{align*}
\]
Example (*T-Sub* with *T-App* on the right)

\[
\begin{array}{c}
\vdash \sigma_1 : T_{11} \rightarrow T_{12} \\
\vdash \sigma_2 : T_2 \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_1 : T_{11} \\
\vdash \sigma_2 : T_{11} \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
T_2 <: T_{11} \\
\vdash \sigma_2 : T_{11}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_2 : T_2 \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_1 : T_{11} \rightarrow T_{12} \\
\vdash \sigma_2 : T_{11} \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_2 : T_{11} \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_1 : T_{11} \rightarrow T_{12} \\
\vdash \sigma_2 : T_{11} \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]

\[
\begin{array}{c}
\vdash \sigma_2 : T_{11} \\
\vdash \sigma_1 \sigma_2 : T_{12}
\end{array}
\]
Example (\(T\text{-}SUB\) with \(T\text{-}APP\) on the right)

\[
\Gamma \vdash s_1 : T_{11} \rightarrow T_{12} \\
\Gamma \vdash s_2 : T_{11} \\
\Gamma \vdash s_1 \ s_2 : T_{12}
\]

becomes

\[
\Gamma \vdash s_1 : T_{11} \rightarrow T_{12} \\
\Gamma \vdash s_2 : T_{11} \\
\Gamma \vdash s_1 \ s_2 : T_{12}
\]

\(\vdash T_{12} \prec : T_{11}\)

\(\vdash T_{12} \prec : T_{12}\)

\(\vdash T_{11} \prec : T_{12}\)

\(\vdash s_1 : T_{12}\)
Intuitions

So we’ve seen that uses of subsumption can be “pushed” from one of immediately before $T$-App’s premises to the other, but cannot be completely eliminated.
Example (nested uses of $\text{T-SUB}$)

\[
\Gamma \vdash s : S \\
\Gamma \vdash s : S <: U \\
\Gamma \vdash s : U \\
\Gamma \vdash s : T
\]
Example (nested uses of $T\text{-}SUB$)

\[
\Gamma \vdash s : S \quad S <: U
\]

\[
\Gamma \vdash s : U \quad U <: T
\]

\[
\Gamma \vdash s : T
\]

becomes

\[
\Gamma \vdash s : S
\]

\[
\Gamma \vdash s : U
\]

\[
\Gamma \vdash s : T
\]
Summary

What we’ve learned:

- Uses of the $T_{\text{SUB}}$ rule can be “pushed down” through typing derivations until they encounter either
  1. a use of $T_{\text{APP}}$ or
  2. the root of the derivation tree.

- In both cases, multiple uses of $T_{\text{SUB}}$ can be collapsed into a single one.
Summary

What we’ve learned:

- Uses of the $T_{\text{SUB}}$ rule can be “pushed down” through typing derivations until they encounter either
  1. a use of $T_{\text{APP}}$ or
  2. the root of the derivation tree.

- In both cases, multiple uses of $T_{\text{SUB}}$ can be collapsed into a single one.

This suggests a notion of “normal form” for typing derivations, in which there is

- exactly one use of $T_{\text{SUB}}$ before each use of $T_{\text{APP}}$
- one use of $T_{\text{SUB}}$ at the very end of the derivation
- no uses of $T_{\text{SUB}}$ anywhere else.
Algorithmic Typing

The next step is to “build in” the use of subsumption in application rules, by changing the $T\text{-App}$ rule to incorporate a subtyping premise.

\[
\Gamma \vdash t_1 : T_{11} \rightarrow T_{12} \quad \Gamma \vdash t_2 : T_2 \quad \vdash T_2 <: T_{11}
\]

\[
\Gamma \vdash t_1 \ t_2 : T_{12}
\]

Given any typing derivation, we can now

1. normalize it, to move all uses of subsumption to either just before applications (in the right-hand premise) or at the very end

2. replace uses of $T\text{-App}$ with $T\text{-Sub}$ in the right-hand premise by uses of the extended rule rule above

This yields a derivation in which there is just one use of subsumption, at the very end!
Minimal Types

But... if subsumption is only used at the very end of derivations, then it is actually *not needed* in order to show that any term is typable! It is just used to give *more* types to terms that have already been shown to have a type.

In other words, if we dropped subsumption completely (after refining the application rule), we would still be able to give types to exactly the same set of terms — we just would not be able to give as many types to some of them.

If we drop subsumption, then the remaining rules will assign a *unique, minimal* type to each typable term.

For purposes of building a typechecking algorithm, this is enough.
Final Algorithmic Typing Rules

\[
\frac{x:T \in \Gamma}{\Gamma \vdash x:T} \quad \text{(TA-VAR)}
\]

\[
\frac{\Gamma, x:T_1 \vdash t_2 : T_2}{\Gamma \vdash \lambda x:T_1. t_2 : T_1 \to T_2} \quad \text{(TA-ABS)}
\]

\[
\frac{\Gamma \vdash t_1 : T_1 \quad T_1 = T_{11} \to T_{12} \quad \Gamma \vdash t_2 : T_2}{\Gamma \vdash t_1 \ t_2 : T_{12}} \quad \text{(TA-APP)}
\]

\[
\frac{\text{for each } i}{\Gamma \vdash t_i : T_i} \quad \text{(TA-RCD)}
\]

\[
\frac{\Gamma \vdash t_1 : R_1 \quad R_1 = \{l_1:T_1 \ldots l_n:T_n\}}{\Gamma \vdash t_1.l_i : T_i} \quad \text{(TA-PROJ)}
\]
Soundness of the algorithmic rules

**Theorem:** If $\Gamma \models t : T$, then $\Gamma \vdash t : T$. 
Completeness of the algorithmic rules

**Theorem [Minimal Typing]:** If $\Gamma \vdash t : T$, then $\Gamma \triangleright t : S$ for some $S <: T$.

Proof: Induction on typing derivation. (Details on this week’s homework.) (N.b.: All the messing around with transforming derivations was just to build intuitions and decide what algorithmic rules to write down and what property to prove: the proof itself is a straightforward induction on typing derivations.)
Completeness of the algorithmic rules

**Theorem [Minimal Typing]:** If $\Gamma \vdash t : T$, then $\Gamma \triangleright t : S$ for some $S <: T$.

**Proof:** Induction on typing derivation. (*Details on this week’s homework.*)

(N.b.: All the messing around with transforming derivations was just to build intuitions and decide what algorithmic rules to write down and what property to prove: the proof itself is a straightforward induction on typing derivations.)
Meets and Joins
Adding Booleans

Suppose we want to add booleans and conditionals to the language we have been discussing. For the *declarative* presentation of the system, we just add in the appropriate syntactic forms, evaluation rules, and typing rules.

\[
\begin{align*}
\Gamma \vdash \text{true} : \text{Bool} & \quad \text{(T-TRUE)} \\
\Gamma \vdash \text{false} : \text{Bool} & \quad \text{(T-FALSE)} \\
\Gamma \vdash t_1 : \text{Bool} & \quad \Gamma \vdash t_2 : T & \quad \Gamma \vdash t_3 : T \\
\Gamma \vdash \text{if } t_1 \text{ then } t_2 \text{ else } t_3 : T & \quad \text{(T-IF)}
\end{align*}
\]
A Problem with Conditional Expressions

For the *algorithmic* presentation of the system, however, we encounter a little difficulty.

What is the minimal type of

\[
\text{if true then } \{x=\text{true}, y=\text{false}\} \text{ else } \{x=\text{true}, z=\text{true}\}
\]
The Algorithmic Conditional Rule

More generally, we can use subsumption to give an expression

\[
\text{if } t_1 \text{ then } t_2 \text{ else } t_3
\]

any type that is a possible type of both \( t_2 \) and \( t_3 \).

So the \textit{minimal} type of the conditional is the \textit{least common supertype} (or \textit{join}) of the minimal type of \( t_2 \) and the minimal type of \( t_3 \).

\[
\Gamma \vdash t_1 : \text{Bool} \quad \Gamma \vdash t_2 : T_2 \quad \Gamma \vdash t_3 : T_3
\]

\[
\Gamma \vdash \text{if } t_1 \text{ then } t_2 \text{ else } t_3 : T_2 \lor T_3
\]  

(T-IF)
More generally, we can use subsumption to give an expression

\[
\text{if } t_1 \text{ then } t_2 \text{ else } t_3
\]

any type that is a possible type of both \( t_2 \) and \( t_3 \).

So the \textit{minimal} type of the conditional is the \textit{least common supertype} (or \textit{join}) of the minimal type of \( t_2 \) and the minimal type of \( t_3 \).

\[\Gamma \vdash ▶ t_1 : \text{Bool} \quad \Gamma \vdash ▶ t_2 : T_2 \quad \Gamma \vdash ▶ t_3 : T_3 \]
\[\Gamma \vdash ▶ \text{if } t_1 \text{ then } t_2 \text{ else } t_3 : T_2 \lor T_3 \] \hfill (T-I_F)

Does such a type exist for every \( T_2 \) and \( T_3 \)?
Existence of Joins

**Theorem:** For every pair of types $S$ and $T$, there is a type $J$ such that

1. $S <: J$
2. $T <: J$
3. If $K$ is a type such that $S <: K$ and $T <: K$, then $J <: K$.

I.e., $J$ is the smallest type that is a supertype of both $S$ and $T$. 
Examples

What are the joins of the following pairs of types?

1. \{x: Bool, y: Bool\} and \{y: Bool, z: Bool\}? 
2. \{x: Bool\} and \{y: Bool\}? 
3. \{x:\{a: Bool, b: Bool\}\} and
   \{x:\{b: Bool, c: Bool\}, y: Bool\}? 
4. \{\} and Bool? 
5. \{x:{}\} and \{x: Bool\}? 
6. Top→\{x: Bool\} and Top→\{y: Bool\}? 
7. \{x: Bool\}→Top and \{y: Bool\}→Top?
Meets

To calculate joins of arrow types, we also need to be able to calculate *meets* (greatest lower bounds)!

Unlike joins, meets do not necessarily exist. E.g., $\text{Bool} \rightarrow \text{Bool}$ and $\{\}$ have *no* common subtypes, so they certainly don’t have a greatest one!

However...
Existence of Meets

**Theorem:** For every pair of types $S$ and $T$, if there is any type $N$ such that $N <: S$ and $N <: T$, then there is a type $M$ such that

1. $M <: S$
2. $M <: T$
3. If $O$ is a type such that $O <: S$ and $O <: T$, then $O <: M$.

I.e., $M$ (when it exists) is the largest type that is a subtype of both $S$ and $T$.

**Jargon:** In the simply typed lambda calculus with subtyping, records, and booleans...

- The subtype relation *has joins*
- The subtype relation *has bounded meets*
Examples

What are the meets of the following pairs of types?

1. \{x:Bool, y:Bool\} and \{y:Bool, z:Bool\}?
2. \{x:Bool\} and \{y:Bool\}?
3. \{x:{a:Bool, b:Bool}\} and
   \{x:{b:Bool, c:Bool}, y:Bool\}?
4. {} and Bool?
5. \{x:{}\} and \{x:Bool\}?
6. Top → \{x:Bool\} and Top → \{y:Bool\}?
7. \{x:Bool\} → Top and \{y:Bool\} → Top?
Calculating Joins

\[
S \lor T = \begin{cases} 
  \text{Bool} & \text{if } S = T = \text{Bool} \\
  M_1 \rightarrow J_2 & \text{if } S = S_1 \rightarrow S_2 \quad T = T_1 \rightarrow T_2 \\
  \{ j_l : J_l \mid l \in 1..q \} & \text{if } S = \{ k_j : S_j \mid j \in 1..m \} \\
  \{ l_i : T_i \mid i \in 1..n \} & \\
  \{ j_l \mid l \in 1..q \} = \{ k_j \mid j \in 1..m \} \cap \{ l_i \mid i \in 1..n \} \\
  S_j \lor T_i = J_l & \text{for each } j_l = k_j = l_i \\
  \text{Top} & \text{otherwise}
\end{cases}
\]
Calculating Meets

\[ S \land T = \begin{cases} 
S & \text{if } T = \text{Top} \\
T & \text{if } S = \text{Top} \\
\text{Bool} & \text{if } S = T = \text{Bool} \\
J_1 \rightarrow M_2 & \text{if } S = S_1 \rightarrow S_2, \ T = T_1 \rightarrow T_2 \\
& \quad \quad S_1 \lor T_1 = J_1, \ S_2 \land T_2 = M_2 \\
\{m_l : M_l \mid l \in 1..q\} & \text{if } S = \{k_j : S_j \mid j \in 1..m\} \\
& \quad \quad T = \{l_i : T_i \mid i \in 1..n\} \\
& \quad \quad \{m_l \mid l \in 1..q\} = \{k_j \mid j \in 1..m\} \cup \{l_i \mid i \in 1..n\} \\
& \quad \quad S_j \land T_i = M_l \quad \text{for each } m_l = k_j = l_i \\
& \quad \quad M_l = S_j \quad \text{if } m_l = k_j \text{ occurs only in } S \\
& \quad \quad M_l = T_i \quad \text{if } m_l = l_i \text{ occurs only in } T \\
\text{fail} & \text{otherwise}
\end{cases} \]