# Spectral analysis of ranking algorithms

Social and Technological Networks

Rik Sarkar

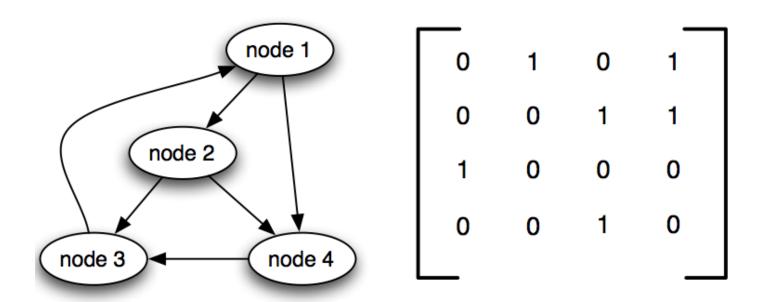
University of Edinburgh, 2019.

## Recap: HITS algorithm

- Evaluate hub and authority scores
- Apply Authority update to all nodes:
  - auth(p) = sum of all hub(q) where q -> p is a link
- Apply Hub update to all nodes:
  - hub(p) = sum of all auth(r) where p->r is a link
- Repeat for k rounds

# Adjacency matrix

Example



# Hubs and authority scores

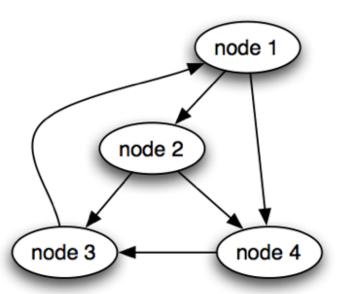
Can be written as vectors h and a

The dimension (number of elements) of the vectors are n

## Update rules

Are matrix multiplications

$$h \leftarrow Ma$$



			_
0	1	0	1
0	0	1	1
1	0	0	0
0	0	1	0

 Hub rule for i : sum of a-values of nodes that i points to:

$$h \leftarrow Ma$$

• Authority rule for i : sum of h-values of *nodes* that point to i:

$$a \leftarrow M^T h$$

## **Iterations**

After one round:

$$a^{\langle 1 \rangle} = M^T h^{\langle 0 \rangle}$$

$$h^{\langle 1 \rangle} = M a^{\langle 1 \rangle} = M M^T h^{\langle 0 \rangle}$$

• Over k rounds:

$$h^{\langle k \rangle} = (MM^T)^k h^{\langle 0 \rangle}$$

## Convergence

- Remember that h keeps increasing
- We want to show that the normalized value

$$\frac{h^{\langle k \rangle}}{c^k}$$

- Converges to a vector of finite real numbers as k goes to infinity
- If convergence happens, then there is a c:

$$(MM^T)h^{\langle * \rangle} = ch^{\langle * \rangle}$$

## Eigen values and vectors

$$(MM^T)h^{\langle * \rangle} = ch^{\langle * \rangle}$$

- Implies that for matrix  $(MM^T)$
- c is an eigen value, with
- $h^{(*)}$  as the corresponding eigen vector

## Proof of convergence to eigen vectors

- Useful Theorem: A symmetric matrix has orthogonal eigen vectors.
  - They form a basis of n-D space
  - Any vector can be written as a linear combination
- $(MM^T)$  is symmetric

- For matrix P with all positive values, Perron's theorem says:
  - A unique positive real valued largest eigen value c exists
  - Corresponding eigen vector y is unique and has positive real coordinates
  - If c=1, then  $P^kx$  converges to y

#### Now to prove convergence:

Suppose sorted eigen values are:

$$|c_1| \ge |c_2| \ge \cdots \ge |c_n|$$

Corresponding eigen vectors are:

$$z_1, z_2, \ldots, z_n,$$

We can write any vector x as

$$x = p_1 z_1 + p_2 z_2 + \dots + p_n z_n$$

• SO:  $(MM^T)x = (MM^T)(p_1z_1 + p_2z_2 + \dots + p_nz_n)$ =  $p_1MM^Tz_1 + p_2MM^Tz_2 + \dots + p_nMM^Tz_n$ =  $p_1c_1z_1 + p_2c_2z_2 + \dots + p_nc_nz_n$ ,

$$(MM^{T})x = (MM^{T})(p_{1}z_{1} + p_{2}z_{2} + \dots + p_{n}z_{n})$$

$$= p_{1}MM^{T}z_{1} + p_{2}MM^{T}z_{2} + \dots + p_{n}MM^{T}z_{n}$$

$$= p_{1}c_{1}z_{1} + p_{2}c_{2}z_{2} + \dots + p_{n}c_{n}z_{n},$$

After k iterations:

$$(MM^T)^k x = c_1^k p_1 z_1 + c_2^k p_2 z_2 + \dots + c_n^k p_n z_n$$

• For hubs:  $h^{\langle k \rangle} = (MM^T)^k h^{\langle 0 \rangle} = c_1^k q_1 z_1 + c_2^k q_2 z_2 + \dots + c_n^k q_n z_n$ 

• So: 
$$rac{h^{\langle k 
angle}}{c_1^k} = q_1 z_1 + \left(rac{c_2}{c_1}
ight)^k q_2 z_2 + \cdots + \left(rac{c_n}{c_1}
ight)^k q_n z_n$$

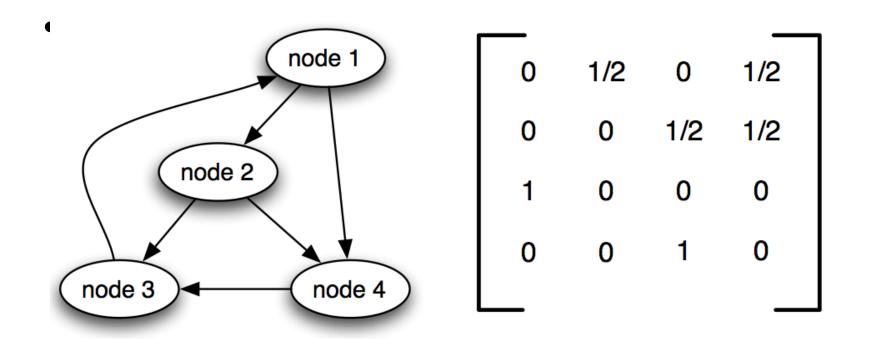
• If  $|c_1| > |c_2|$ , only the first term remains.

• So,  $\frac{h^{\langle k \rangle}}{c_1^k}$  converges to  $q_1 z_1$ 

## **Properties**

- The vector q<sub>1</sub>z<sub>1</sub> is a simple multiple of z<sub>1</sub>
  - A vector essentially similar to the first eigen vector
  - Therefore independent of starting values of h
- q1 can be shown to be non-zero always, so the scores are not zero
- Authority score analysis is analogous

# Pagerank Update rule as a matrix derived from adjacency



$$r \leftarrow N^T r$$

Scaled pagerank:

$$r \leftarrow \tilde{N}^T r$$

Over k iterations:

$$r^{\langle k \rangle} = (\tilde{N}^T)^k r^{\langle 0 \rangle}$$

Pagerank does not need normalization.

$$\tilde{N}^T r^{\langle * \rangle} = r^{\langle * \rangle}$$

 We are looking for an eigen vector with eigen value=1

## Random walk interpretation

- The pagerank "Fluid" from a source s spreads to neighbors and their neighbors
- Take the small quantity that ends at node x
- We can trace back the walk that it took
- This is a random walk, since at each step it was sent to a random neighbor
- Thus the fluid dividing and spreading randomly is equivalent to
  - Several small particles of fluid starting at s and doing a random walk

### Random walks

- A random walker is moving along random directed edges
- Suppose vector b shows the probabilities of walker currently being at different nodes
- Then vector  $N^Tb$  gives the probabilities for the next step

## Random walks

- Thus, pagerank values of nodes after k iterations is equivalent to:
  - The probabilities of the walker being at the nodes after k steps
- The final values given by the eigen vector are the steady state probabilities
  - Note that these depend only on the network and are independent of the starting points

## History of web search

- YAHOO: A directory (hierarchic list) of websites
  - Jerry Yang, David Filo, Stanford 1995
- 1998: Authoritative sources in hyperlinked environment (HITS), symposium on discrete algorithms
  - Jon Kleinberg, Cornell
- 1998: Pagerank citation ranking: Bringing order to the web
  - Larry Page, Sergey Brin, Rajeev Motwani, Terry Winograd, Stanford techreport

## Course

- No class or office hours next week
- No limit on number of self-proposed projects
  - Send an email if doodle is full
- Instructions for project plan submission out today.
- Look at last year's slides on some other topics in STN for approaches in networks to use in projects
  - Community detection
  - Graph kernels (as a method to find similar graphs)
  - Link prediction (under miscellaneous topics)