Epidemics and gossip algorithms

Social and Technological Networks

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University of Edinburgh, 2019.

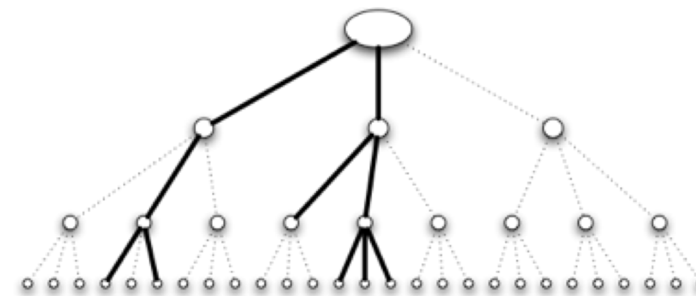
Spread of diseases

- Pattern depends on network structure
- e.g. spread of flu
- Network of people
- Network of airlines

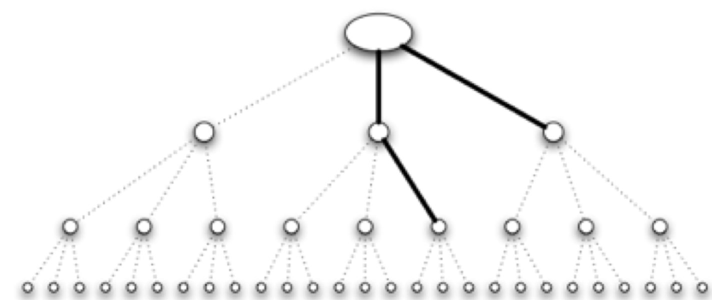
- Different from idea/innovation contagion
  - Does not need a “decision”
  - Does not need multiple support
    - Infectious disease passes easily with some probability

- Suppose everyone meets k new people and infects each with probability p
- That is, they infect $R = kp$ people on average

- If p is high
- The disease will persists through rounds
- If p is low, it will die out after some rounds



(b) With high contagion probability, the infection spreads widely



Property

- When R > 1 number of infected people keeps increasing
  - Outbreak
- When R < 1 Number of infected people decreases
  - Disease dies out
- Phase transition at R = 1
- assuming there are enough “new” people supply to meet
- Generally true in the initial stages

- Around R = 1: small efforts can have large effects on epidemic
  - Awareness causing slight decrease in p
  - Quarantine/fear causing slight decrease in k

SIR Model

- Susceptible (initially)
- Infectious (after being infected)
  - While Infectious, it can pass disease to each neighbor in each step with prob. p
- Removed (after given duration as Infectious)
  - Immune/dead

SIS model

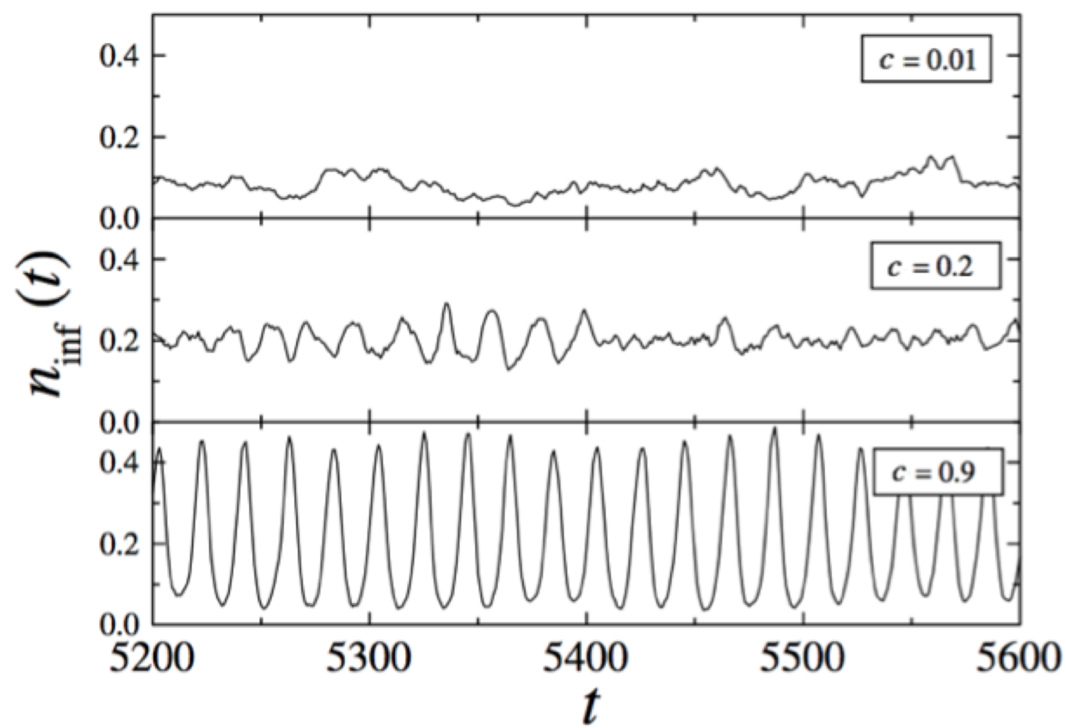
- No “Removed” state. Susceptible follows Infectious

SIRS model

- Susceptible
- Infectious
- Recovered (immune)
- Susceptible

SIRS oscillations in Watts-Strogatz Small worlds

- Nodes connected to few nighbors on a ring
- Fraction c of links modified to connect to random nodes



Epidemic or gossip algorithm

- Emulates the spread of epidemic or a rumor in a network
- A node speaks to a random neighbor to spread the rumor message
- Useful for spreading information in computer networks

Spreading a message via gossip

- Complete graph: Anyone can call anyone
- Problem: One node has a message or rumor to spread. How does it spread it to all nodes in the network?
- Calling everyone will take O(n) rounds.
- After first round, nodes with the rumor can help
  - But how do you avoid collision?

Spreading a message via gossip

- Complete graph: Anyone can call anyone
- Problem: One node has a message or rumor to spread. How does it spread it to all nodes in the network?
- Strategy: In each round, anyone with the rumor calls one random node and passes the rumor

Theorem

- Everyone gets the message in O(log n) rounds
- Idea:
- n/3 nodes get the message in log n rounds
  - Current number of infected nodes m < n/3
  - Probability that a call goes to a new node is at least 2/3
    - Number of calls to new nodes: 2m/3
  - Probability of a collision at the new node is 1/(n-m)
  - $O(m^2)$ possible pairs for collision
  - Max possible collisions: $\frac{m^2}{2(n-m)} \leq \frac{m^2}{2} \cdot \frac{1}{2m} = \frac{m}{4}$
- Number of newly infected nodes at least $\frac{2m}{3} - \frac{m}{4} = \frac{5m}{12}$

- while m < n/3
  - m grows to m(1 + 5/12) = 17m/12 every round
  - m grows to n/3 in O(log n) rounds
- After m> n/3
  - Probability that a node is not called in 1 round is
    $$(1-\frac{1}{n})^{\frac{n}{3}} \leq \mathrm{e}^{\frac{-1}{n}\cdot\frac{n}{3}} = \mathrm{e}^{-1/3}$$
  - Probability that 1 or more nodes are not called after O(log n) rounds
  - Less than $1/n^c$, where c depends on the constant in the O

- See Kempe 18.

Advantages of gossip

- Simple to implement
- Robust algorithm
  - A node failure does not stop the computation
  - Easy to add nodes to the system
  - At the cost of a logarithmic factor of increased costs

Averaging gossip

- Suppose the nodes all have a “value”
- And we wish to compute a linear function of these values
- E.g. the average

- The push-sum protocol
  - In every round
  - Every node takes a fraction of its value and sends to a random neighbor
  - It adds all received values to its current value
- The pairwise averaging protocol
  - In every round, a node talks to one other random neighbor
  - Both nodes set their values to the average of the two

Gossip averaging protocols

- On a complete graph
  - Both protocols converge to the average fast
  - O(log n) rounds
- On other graphs, convergence depends on structure
  - Graphs with small spectral gap have slow convergence