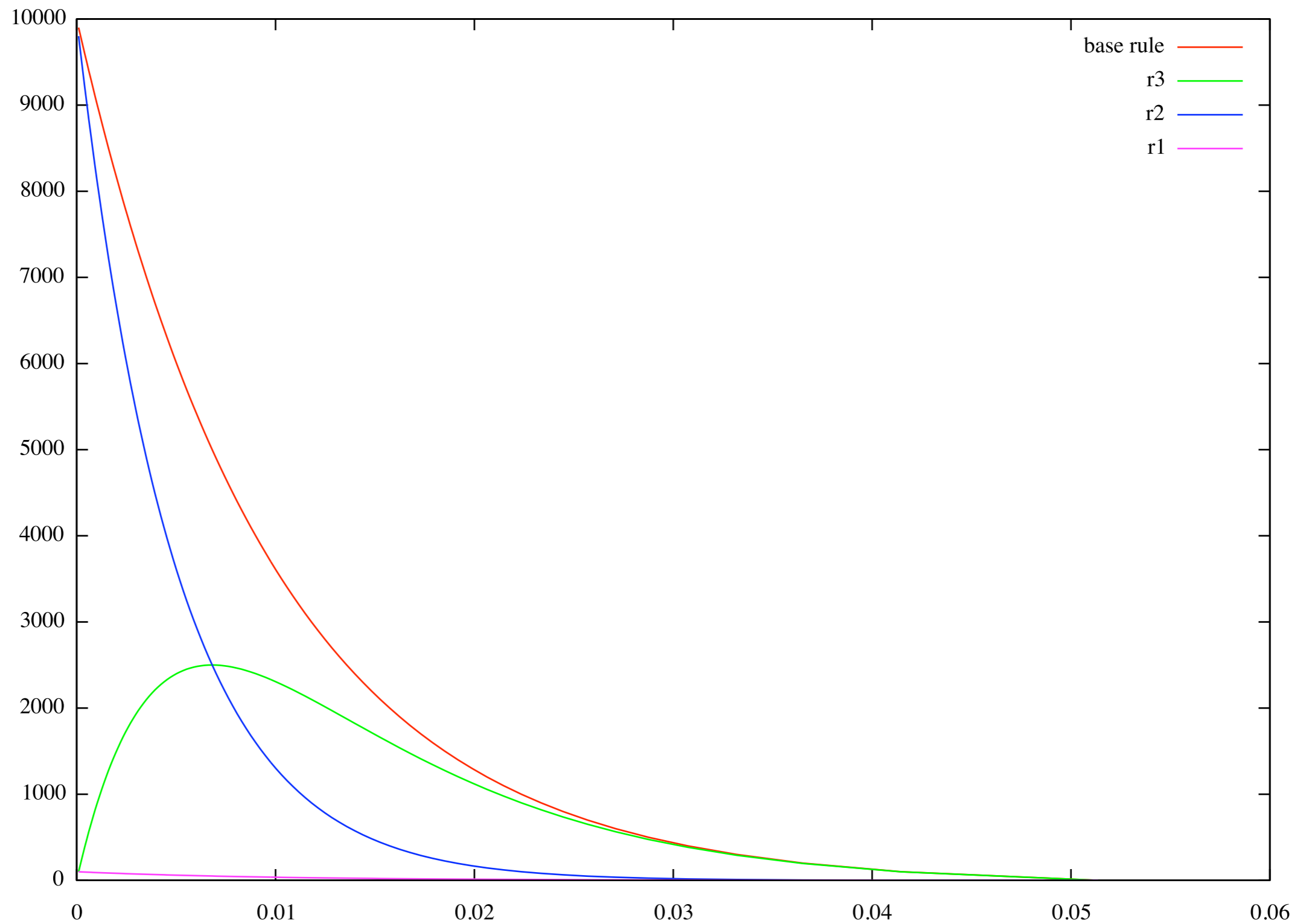


Refinement of rules

- *Don't care, don't write* means that a rule only describes *part* of its actual context
 - many actual contexts are possible (these are all the possible *instances* of the rule)
- We refine a rule by making explicit some (or all) of that missing context
 - allows us to examine the relative importance of different instances
 - allows us to fill in “missing context”

Relative importance



Relative importance (2)

- For this to make sense, we need to ensure that the set of refined rules has the same behaviour as the original rule...
- A refinement is *neutral* if this is the case
 - the appropriate rate constants can be calculated statically
 - the neutral refinement gives a “base line” from which kinetic perturbations can be made...

“Missing context”

- Suppose we have the rule (from before)

$$A(b), B(a) \rightarrow A(b!1), B(a!1) @ k$$

- Then somebody tells us that this binding is far more likely if B is already bound to C...

$$A(b), B(a, c) \rightarrow A(b!1), B(a!1, c) @ k_1$$

$$A(b), B(a, c!_) \rightarrow A(b!1), B(a!1, c!_) @ k_2$$

- Neutral refinement: $k = k_1 = k_2$
- Our desired *non-neutral* refinement: $k_1 \ll k_2$

Formally...

- Describe a “soup” as a *graph-with-sites (gws)*
- An instance of a rule is a *gws-homomorphism* from the rule’s LHS to the current soup
- For any choice of “additional context”, we can find a set (actually multiset) of refinements
 - divides the rule into *disjoint* subcases (but cases may be repeated...)
 - n repeated cases scales the rate constant by n

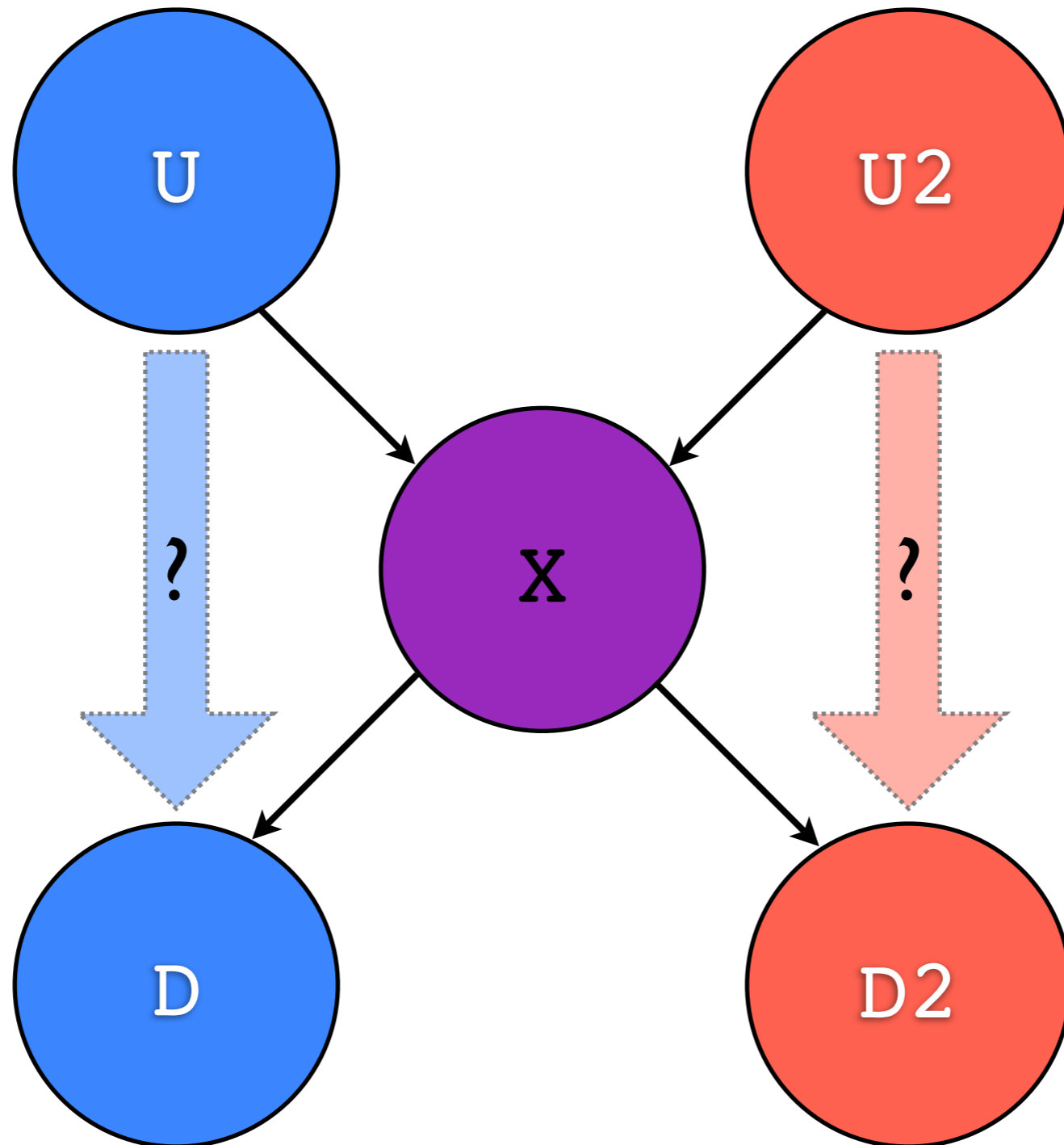
Refinement summary

- Completely formal notion of rule refinement
 - can be used to “split” a rule into several pieces and evaluate their relative contributions
- Neutral refinements assign rates to the rule pieces that leave behaviour unchanged
- Kinetic refinement then modifies this, e.g.
 - an enzyme may have lower affinity for its product than its substrate
 - GPCRs (bond strength between α and $\beta\gamma$)

Refinement (2)

- Another view... refinements reflect small “perturbations” of a rule:
 - expose some previously hidden bit of context
 - use non-neutral kinetics to modify behaviour
- A plausible mechanism by which a signalling network could be subjected to selection
 - e.g. different affinities of splice variants

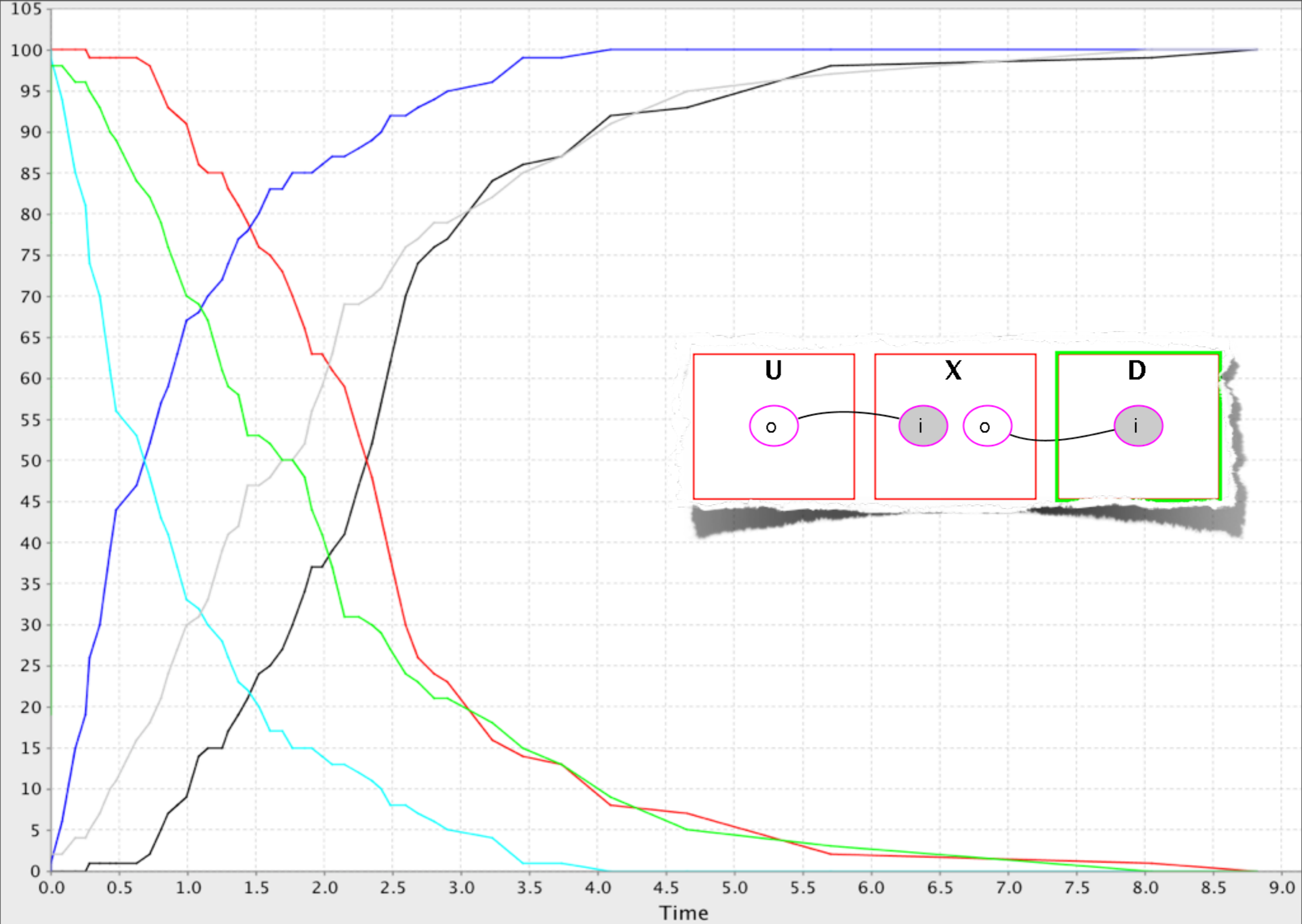
A specificity puzzle



A simple cascade:

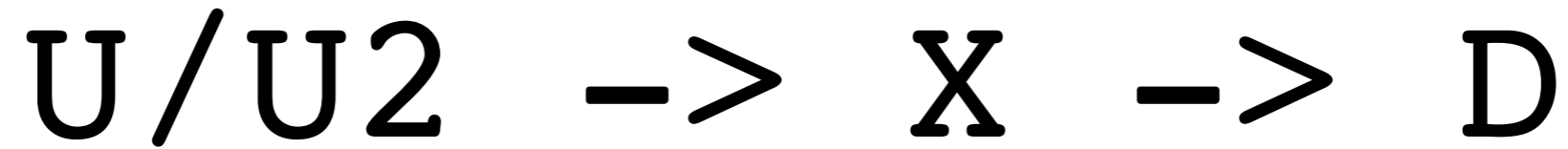
$$U \rightarrow X \rightarrow D$$

- $U(s), X(s \sim u) \rightarrow U(s!1), X(s \sim u!1)$
 $U(s!1), X(s!1) \rightarrow U(s), X(s)$
 $U(s!1), X(s \sim u!1) \rightarrow U(s!1), X(s \sim p!1)$
- $X(s \sim p?, d), D(s \sim u) \rightarrow X(s \sim p?, d!1), D(s \sim u!1)$
 $X(d!1), D(s!1) \rightarrow X(d), D(s)$
 $X(d!1), D(s \sim u!1) \rightarrow X(d!1), D(s \sim p!1)$

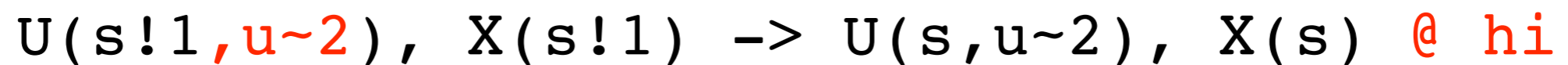
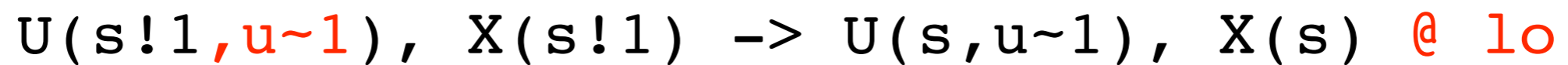


— [D(i~p?)] — [D(i~u?)] — [X(i~p?)] — [X(i~u?)] — [X(i)] — [X(i!_)]

Refining U:



- refine the U/I *off* rule to:



- So U now exists in two distinct forms with different affinities for the common substrate X

Refining D:

$$U/U2 \rightarrow X \rightarrow D/D2$$

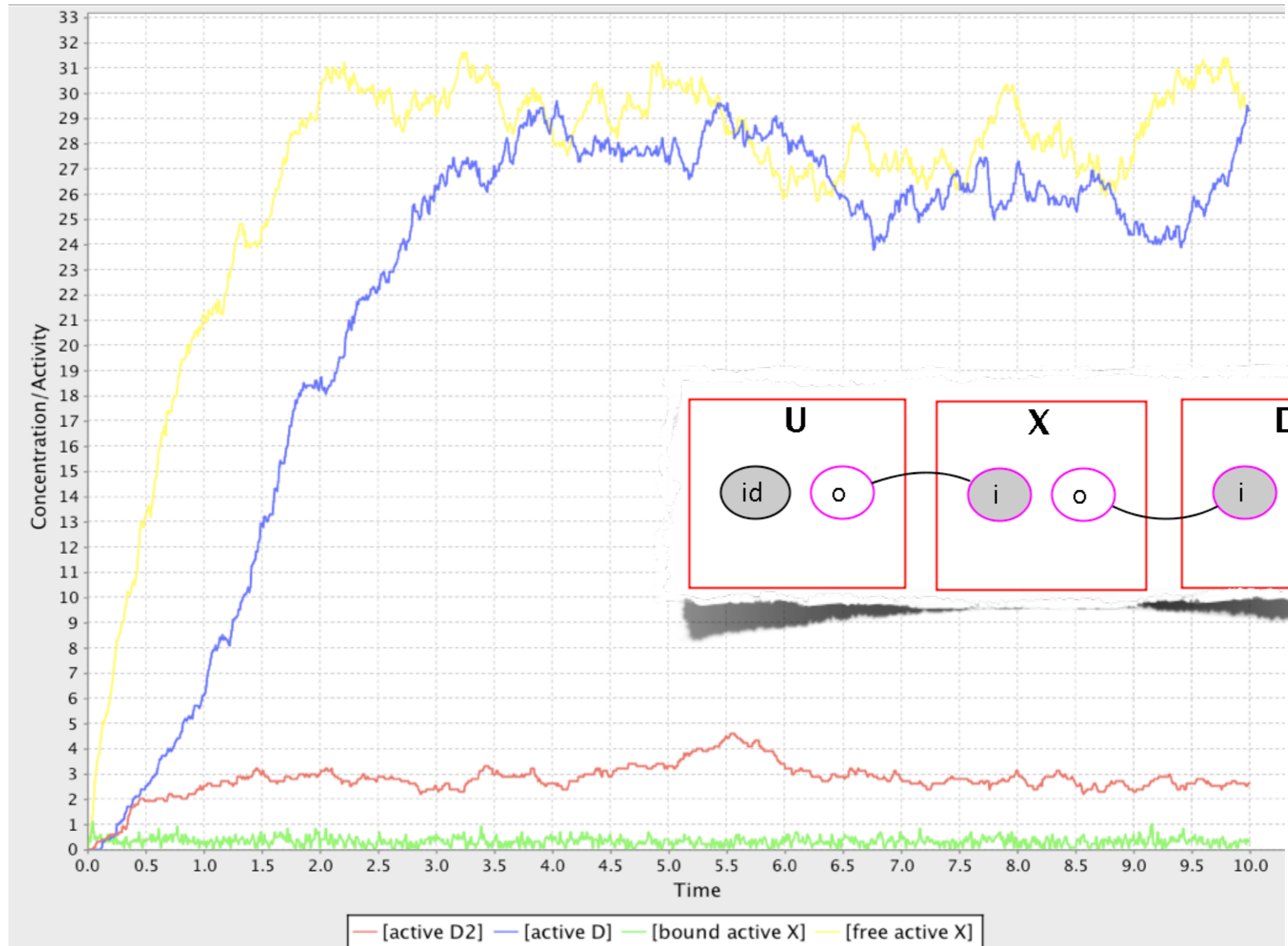
- refine the I/D *on* rule to:

$$X(s \sim p, d), D(s \sim u, d \sim 1) \rightarrow X(s \sim p, d!1), D(s \sim u!1, d \sim 1)$$

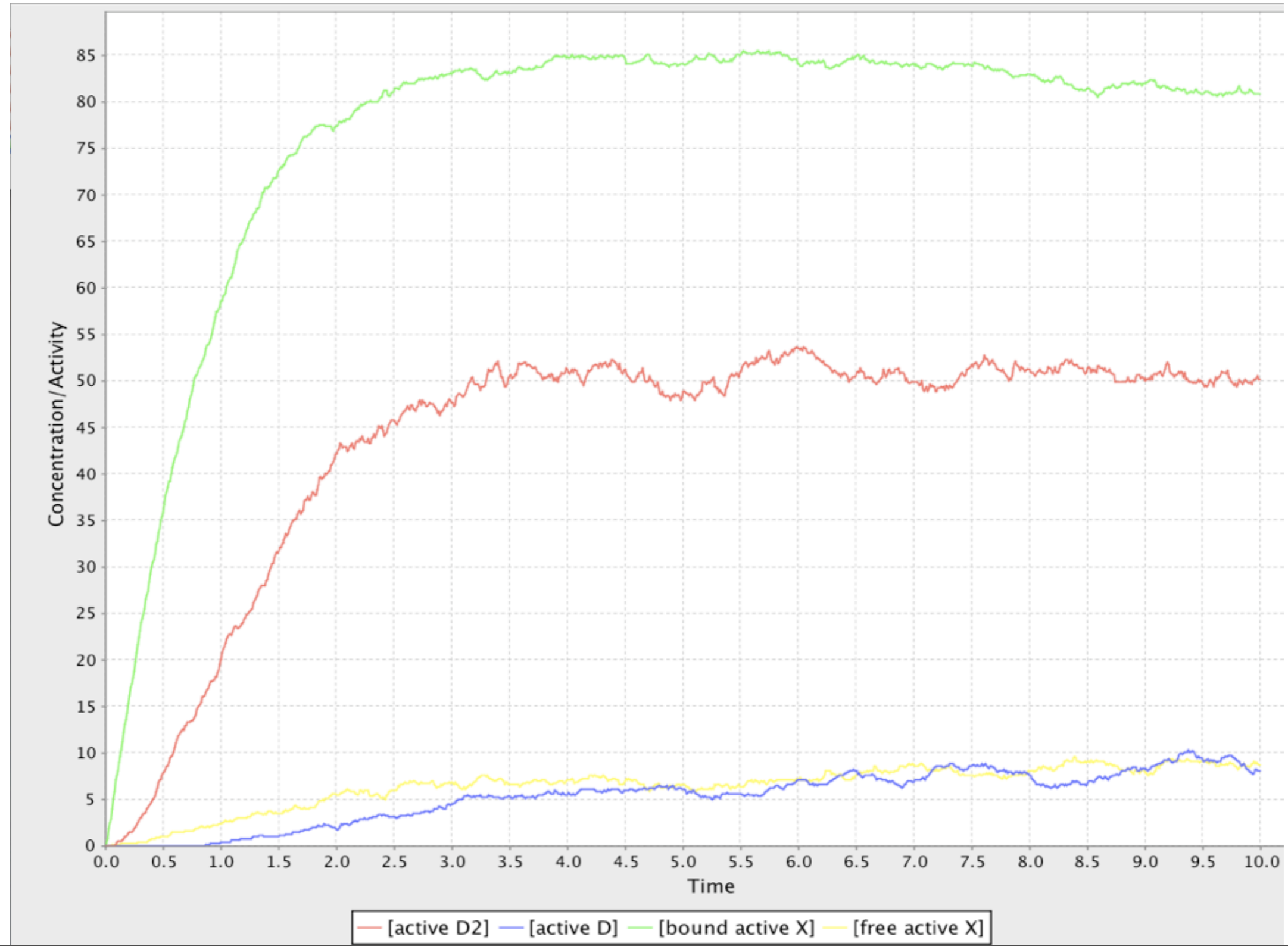
$$X(s \sim p!_, d), D(s \sim u, d \sim 2) \rightarrow X(s \sim p, d!1), D(s \sim u!1, d \sim 2)$$

- So D now exists in two distinct forms
- X has two different “binding configurations” (which depend on whether it is bound or not)

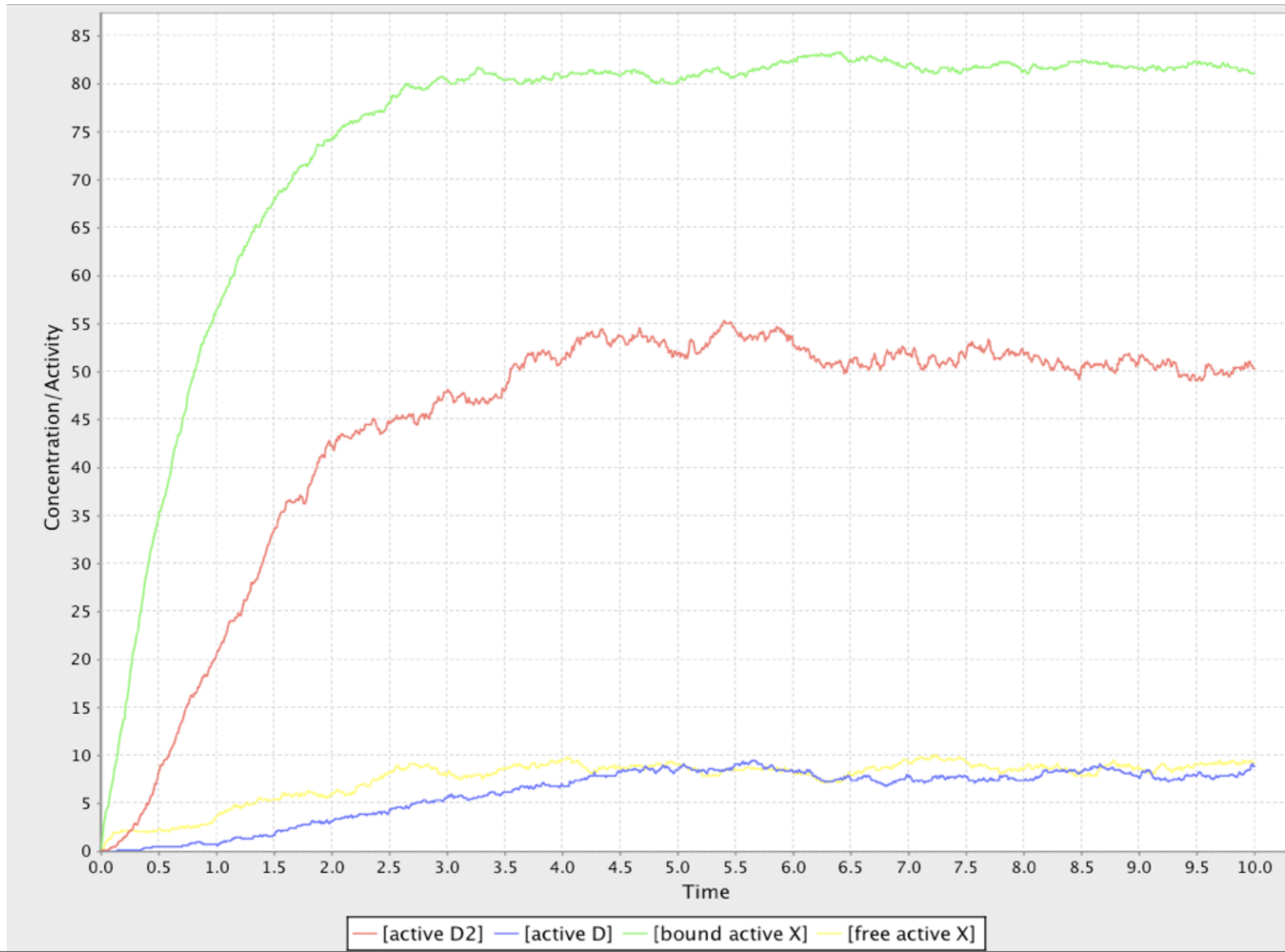
U only



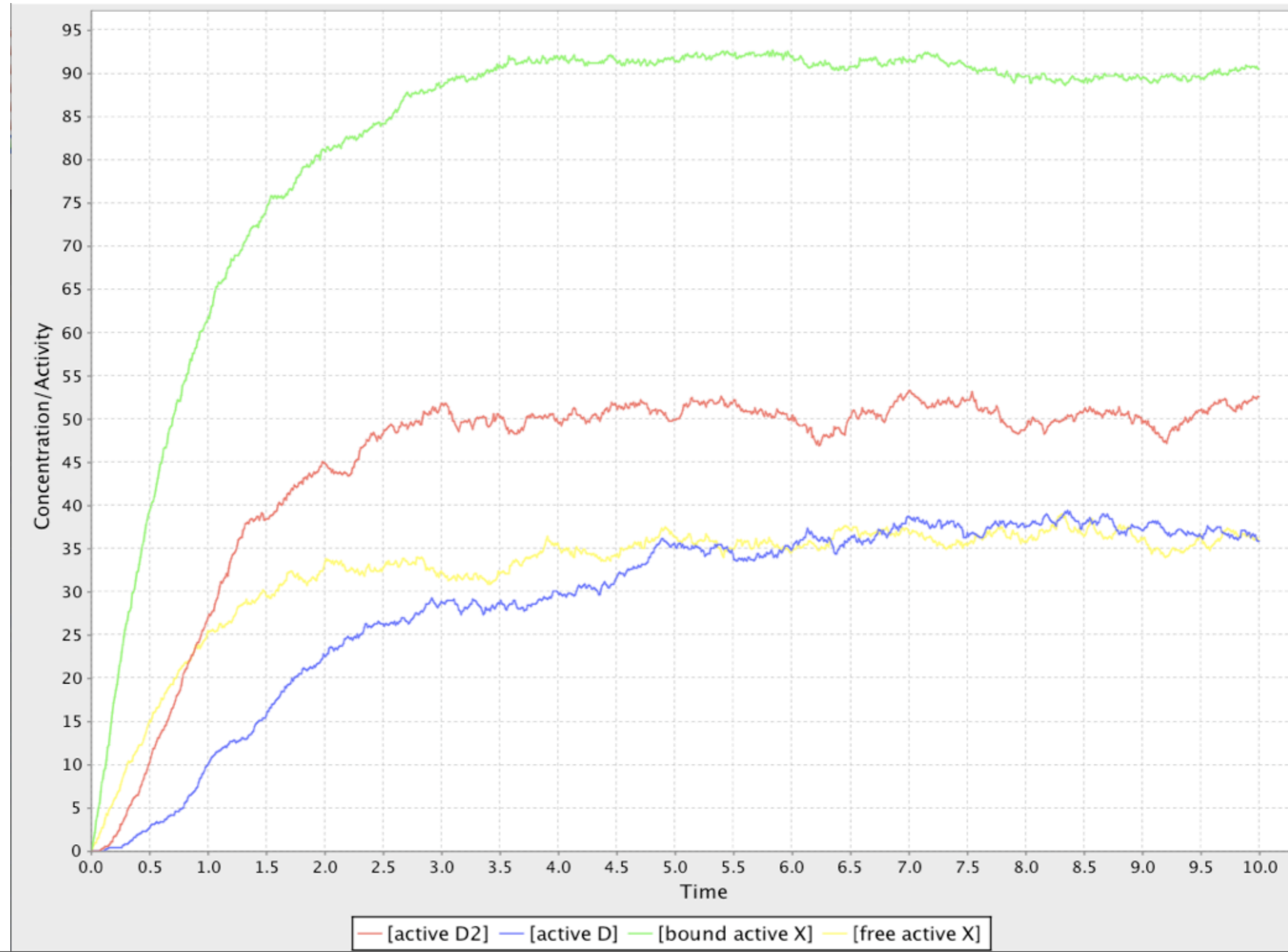
U2 only



U and U2; X limited



U and U2; X in excess



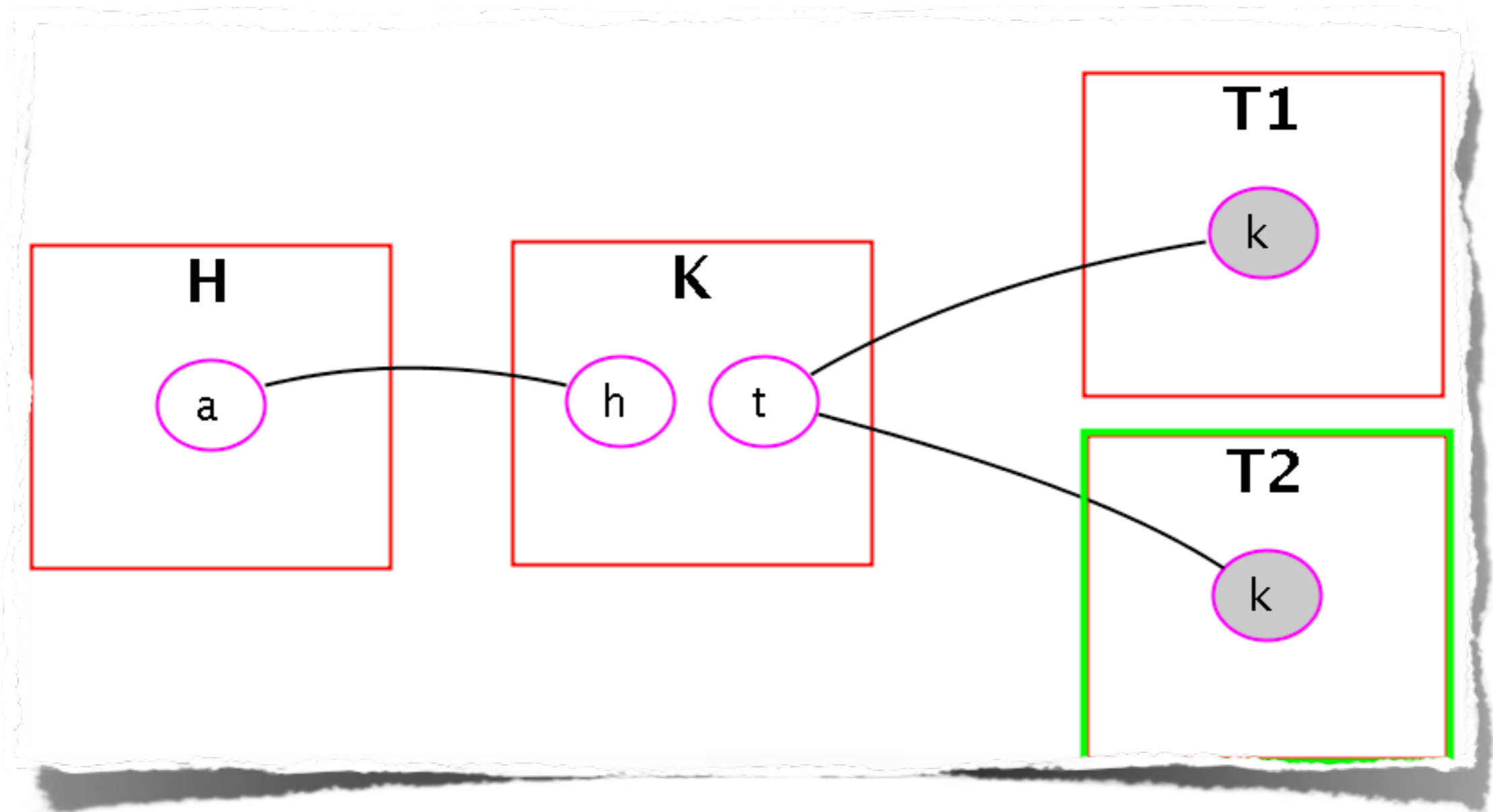
Specificity summary

- Little leakage (aka specificity):
 - U predominantly activates D
 - U2 predominantly activates D2
- If U and U2 are present,
 - limited X implies U2 (high affinity) “wins”
 - excess X allows both U and U2 to signal

Specificity (2)

- We might say that the $U_2 \rightarrow X \rightarrow D_2$ pathway *hijacks* the original $U \rightarrow X \rightarrow D$ pathway (when X is limited)
- Behaves analogously to a *transistor!*
- An example of *network plasticity*: the amount of X determines whether the network behaves as a transistor or simply as two parallel wires
 - a “scalpel” to divide a rule into pieces
 - a possible mechanism of “network evolution”

Another (puzzle)



players and a question

- K -repairs> $T1, T2$ (as target)
- K also binds H (as helper)
- $T2$ needs additional H
- how does H know where to go??
 - Swiss knife approach, saturate K with H
 - stochastic honey pot approach, uses refinements too

model 1: rules

- 'off KTI' $K(t!1), TI(k!1) \rightarrow K(t), TI(k) @ 100$
- 'on KTI' $K(t), TI(k) \rightarrow K(t!1), TI(k!1) @ 1$
- 'mod KTI' $K(t!1), TI(k\sim no!1) \rightarrow K(t!1), TI(k\sim yes!1) @ 50$

- 'unmod TI' $TI(k\sim yes?) \rightarrow TI(k\sim no?) @ 1$

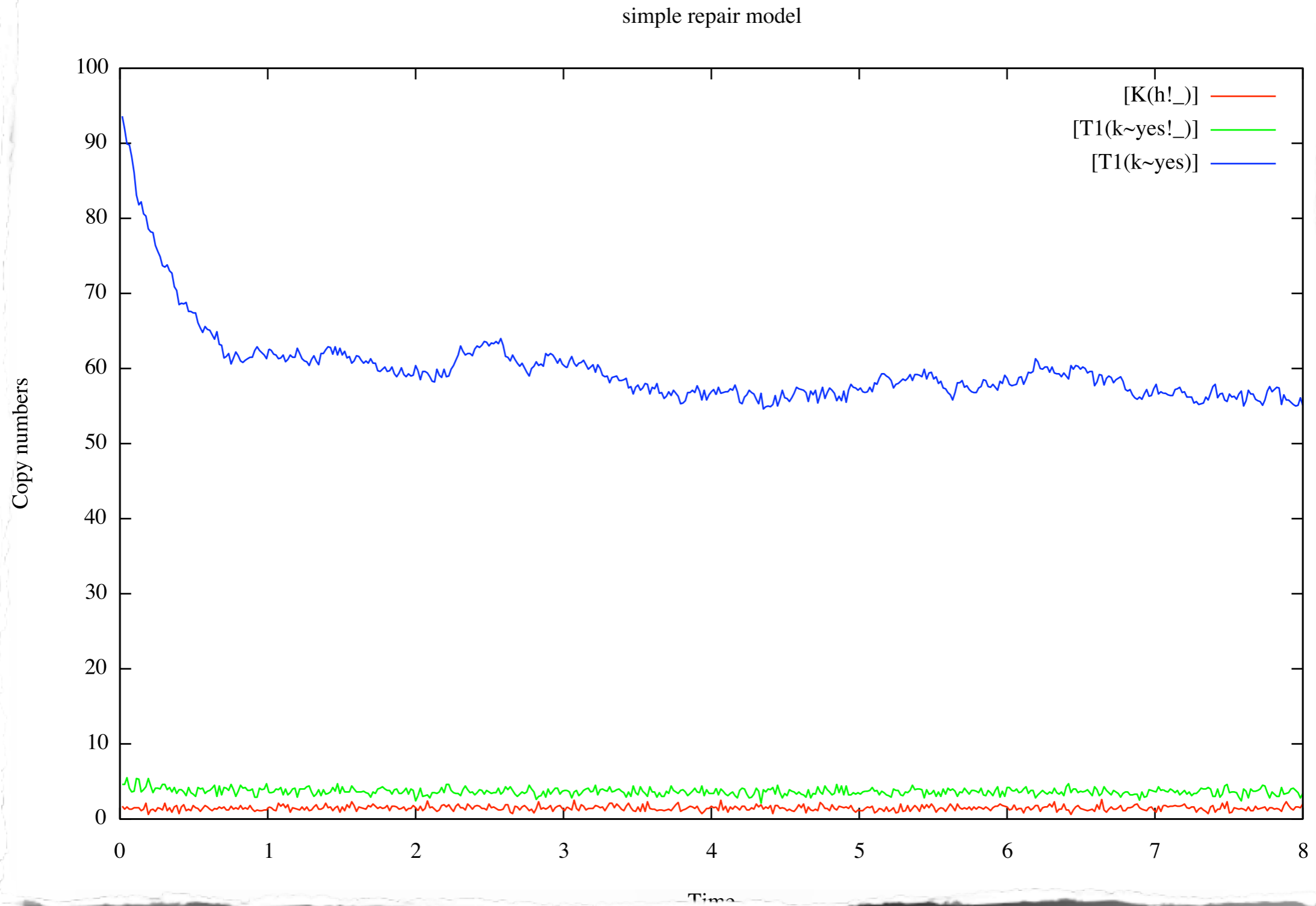
- 'off KH' $K(h!1), H(k!1) \rightarrow K(h), H(k) @ 500$
- 'on KH' $K(h), H(k) \rightarrow K(h!1), H(k!1) @ 10$

- %init: $10 * (K(h,t))$
- %init: $10 * (H(k))$
- %init: $100 * (TI(k\sim yes))$

- %obs: $TI(k\sim yes)$
- %obs: $TI(k\sim yes!_)$
- %obs: $K(h!_)$

model 1: results

model 1: results



model 2: rules

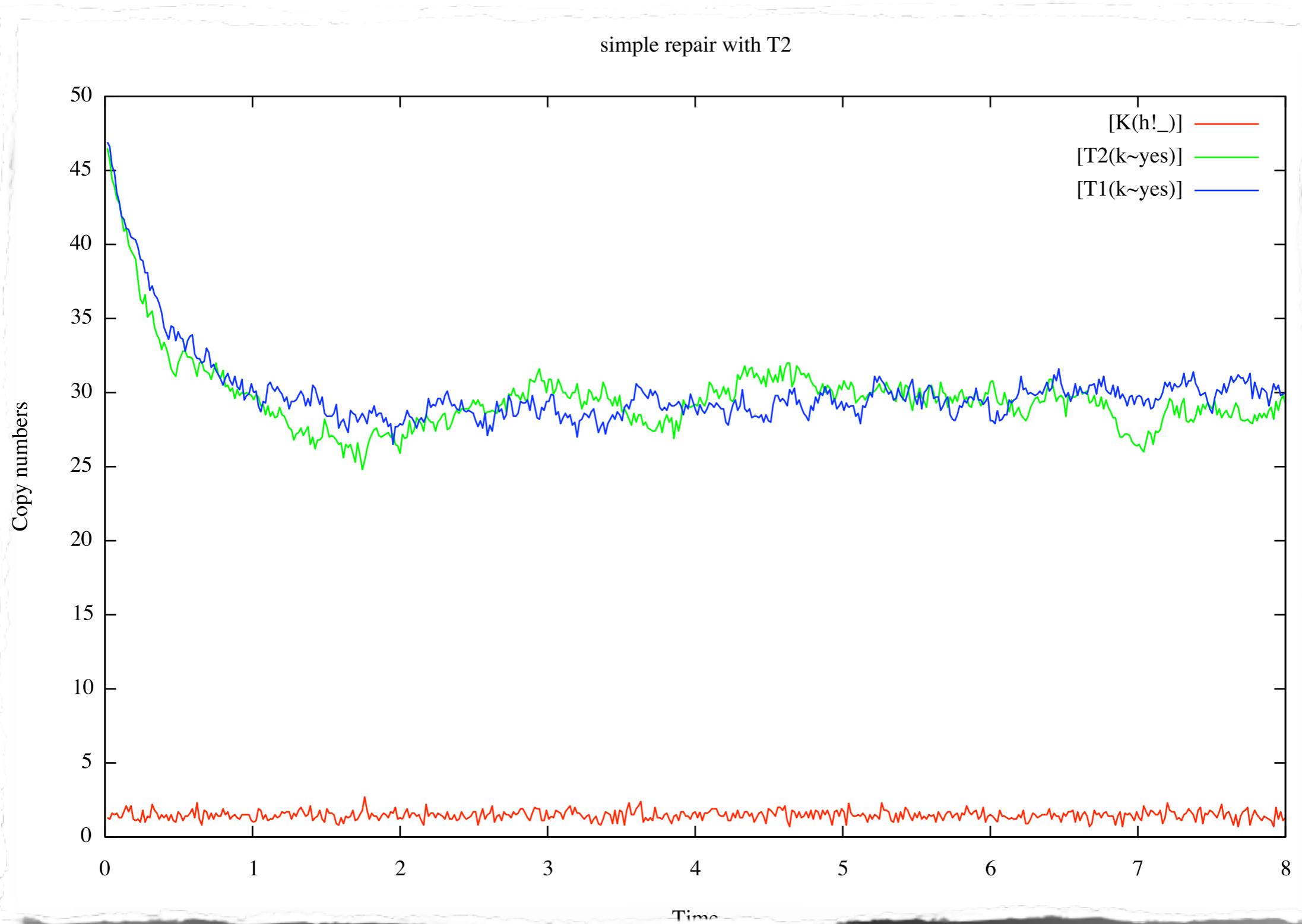
- 'off KT2' $K(t!1), T2(k!1) \rightarrow K(t), T2(k) @ 100$
- 'on KT2' $K(t), T2(k) \rightarrow K(t!1), T2(k!1) @ 1$
- 'mod KT2' $K(t!1), T2(k\sim no!1) \rightarrow K(t!1), T2(k\sim yes!1) @ 50$

- 'unmod T2' $T2(k\sim yes?) \rightarrow T2(k\sim no?) @ 1$

- %init: $10 * (K(h,t))$
- %init: $10 * (H(k))$
- %init: $50 * (T1(k\sim yes))$
- %init: $50 * (T2(k\sim yes))$

- %obs: $T1(k\sim yes)$
- %obs: $T2(k\sim yes)$
- %obs: $K(h!_)$

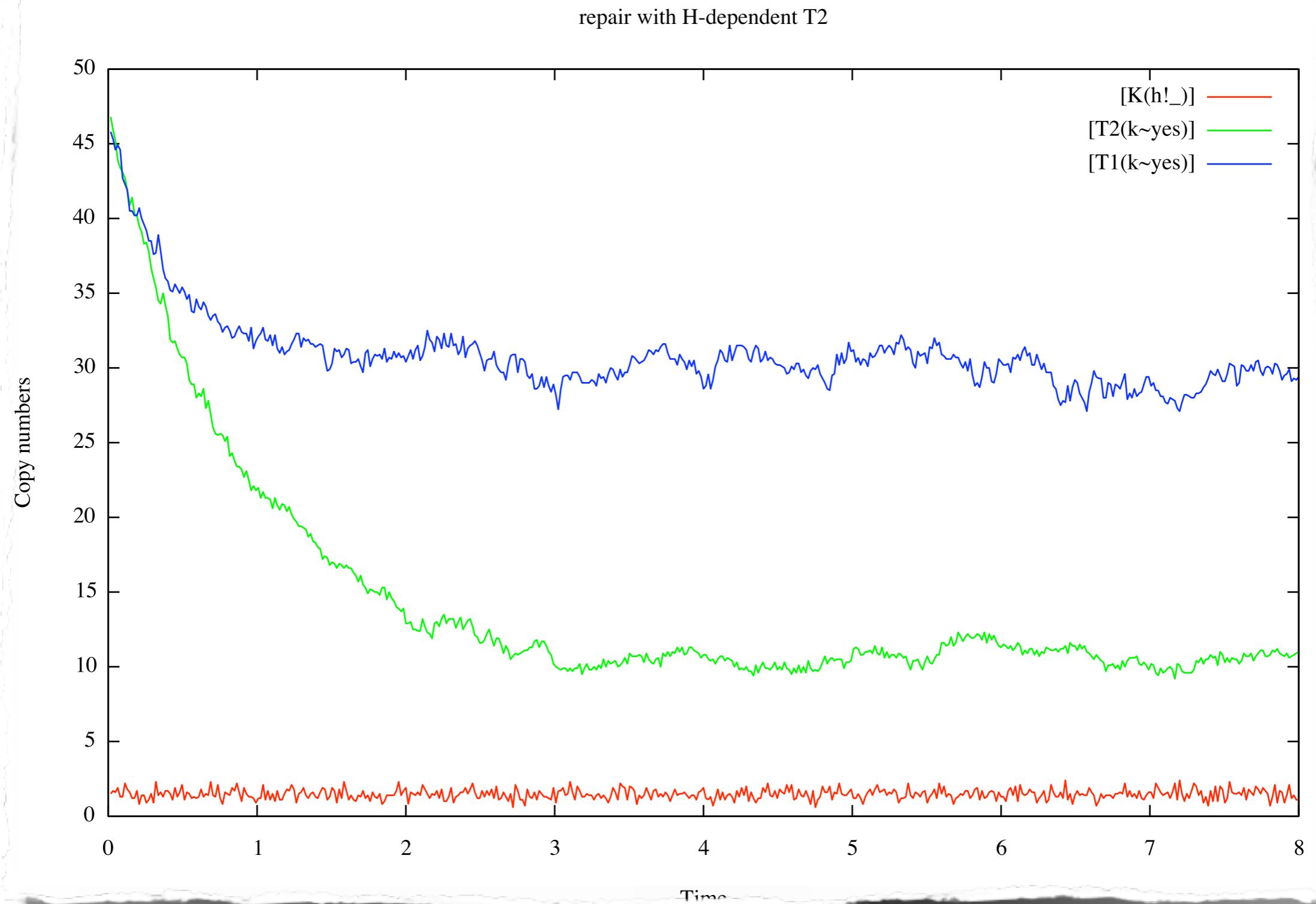
model 2: results



model 3: rules

- `#'mod KT2' K(t! I),T2(k~no! I) -> K(t! I),T2(k~yes! I) @ 50`
- `'mod KT2' K(t! I,h! _),T2(k~no! I) -> K(t! I,h! _),T2(k~yes! I) @ 50`

model 3: results



model 4: rules

- #'off KH' $K(h!1), H(k!1) \rightarrow K(h), H(k) @ 500$
- $K(h!1, t), H(k!1) \rightarrow K(h, t), H(k) @ 1000$
- $K(h!1, t!2), H(k!1), T1(k!2) \rightarrow K(h, t!2), H(k), T1(k!2) @ 1000$
- $K(h!1, t!2), H(k!1), T2(k!2) \rightarrow K(h, t!2), H(k), T2(k!2) @ 0.001$

model 4: results

