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# Performance Modelling — Lecture 9 Using a GSPN for Performance Evaluation

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### From SPN to Markov Process

Generating the Markov process underlying an SPN model is very straightforward.

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We associate a state, x<sub>i</sub>, in the Markov process with every marking, M<sub>i</sub>, in the reachability graph of the SPN;

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- We associate a state, x<sub>i</sub>, in the Markov process with every marking, M<sub>i</sub>, in the reachability graph of the SPN;
- The transition rate from state x<sub>i</sub> (corresponding to marking M<sub>i</sub>) to state x<sub>j</sub> (M<sub>j</sub>), is obtained as the sum of the firing rates of the transitions that are enabled in M<sub>i</sub> and whose firings generate marking M<sub>j</sub>.

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#### Simple Example



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#### From GSPN to Markov Process

GSPN had two additional features compared to SPN: inhibitor arcs and immediate transitions.

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The effect of the inhibitor arcs is only to eliminate some potential markings and transitions from the reachability graph. Once such a graph is constructed it can be mapped to a Markov process just as in the case for SPN.

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The effect of the inhibitor arcs is only to eliminate some potential markings and transitions from the reachability graph. Once such a graph is constructed it can be mapped to a Markov process just as in the case for SPN.

The effect of the immediate transitions is to create some markings which do not correspond to states in a Markov process, so-called vanishing markings.

### Vanishing and Tangible Markings

If a marking in a GSPN enables an immediate transition, by the firing rules, the immediate transition must fire instantaneously, and so the marking will be changed again without any time elapsing.

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In contrast if a marking enables only timed transitions there will be an exponentially distributed delay before a transition fires, which corresponds to an exponentially distributed sojourn time in the corresponding state in the Markov process.

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In contrast if a marking enables only timed transitions there will be an exponentially distributed delay before a transition fires, which corresponds to an exponentially distributed sojourn time in the corresponding state in the Markov process.

Such markings are called tangible as opposed to vanishing.

# Eliminating Vanishing Markings

Vanishing markings must be eliminated from the reachability graph before the state space of the Markov process can be generated.

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## Eliminating Vanishing Markings

Vanishing markings must be eliminated from the reachability graph before the state space of the Markov process can be generated.

If immediate transitions from a marking can lead to two or more different markings, the transition rates to these markings need to be adjusted (according to the decomposition principle).

# Eliminating Vanishing Markings: a single immediate transition

 Suppose a vanishing marking *M<sub>v</sub>*, enables a single immediate transition, *T<sub>j</sub>*;



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# Eliminating Vanishing Markings: a single immediate transition

- Suppose a vanishing marking M<sub>v</sub>, enables a single immediate transition, T<sub>j</sub>;
- There will be a single successor marking, M<sub>s</sub>, which is the result of firing T<sub>j</sub>.



# Eliminating Vanishing Markings: a single immediate transition

To eliminate  $T_j$  for each predecessor marking  $M_p$ , linked to  $M_v$  by an arc labelled  $T_i$  with rate  $r_i$ 

• delete  $M_v$ ;



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- There will be a successor marking, M<sub>sn</sub>, for the firing of each T<sub>jn</sub>.
- These transitions will be in conflict and so each must have a probability pk such that ∑<sup>n</sup><sub>k=1</sub> pn = 1



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- In the end, all arcs in the modified reachability graph will have a rate originating from a single timed transition associated with the arc.
- This rate may have been adjusted during the elimination of vanishing markings to reflect the relative probability of immediate transitions enabled after the timed transition.
- It is this modified reachability graph that is used to generate the Markov process underlying a GSPN model.

#### Reader-Writer Example

 Consider again the system in which there is a set of processes who share access to a common database.

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- On any particular access a process may wish to perform either a read or a write.
- Any number of readers may access the database concurrently;
- a writer requires exclusive access.
- Between accesses processes undertake independent work (concurrently).

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#### GSPN model of the reader-writer system



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# Reachability graph



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#### Reachability set of the reader-writer model

$M_0$	(2, 0, 0, 0, 1, 0, 0)	tangible
$M_1$	(1, 1, 0, 0, 1, 0, 0)	vanishing
$M_2$	(1,0,1,0,1,0,0)	vanishing
<i>M</i> <sub>3</sub>	(1,0,0,1,1,0,0)	vanishing
$M_4$	(1,0,0,0,1,1,0)	tangible
$M_5$	(1,0,0,0,0,0,1)	tangible
$M_6$	(0, 1, 0, 0, 1, 1, 0)	vanishing
$M_7$	(0,0,1,0,1,1,0)	vanishing
<i>M</i> <sub>8</sub>	(0,0,0,1,1,1,0)	tangible
$M_9$	(0, 0, 0, 0, 1, 2, 0)	tangible
<i>M</i> <sub>10</sub>	(0, 1, 0, 0, 0, 0, 1)	vanishing
<i>M</i> <sub>11</sub>	(0,0,1,0,0,0,1)	tangible
<i>M</i> <sub>12</sub>	(0,0,0,1,0,0,1)	tangible

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### Reachability graph



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### Reduced reachability graph



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### Performance Measures from GSPN

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In other words the aim is to derive performance characteristics of the system based on the steady state probability of being in certain states, or markings.

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In other words the aim is to derive performance characteristics of the system based on the steady state probability of being in certain states, or markings.

In GSPN we can identify the states we are interested in by their characteristics at the GSPN level.

### Throughput and Average Marking

Most SPN/GSPN tools will automatically report the performance measures associated with each element of the Petri net:

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- The probability enabled may also be recorded for each transition.
- Sometimes the steady state distribution of tokens in a place will be given.

We can often derive the measures we are interested in directly from these default measures.

#### Throughput and Average Marking: example



In the reader-writer model the throughput of transition  $T_6$  plus the throughput of transition  $T_7$  will give us the throughput of the database in terms of accesses/unit time.

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In the reader-writer model the throughput of transition  $T_6$  plus the throughput of transition  $T_7$  will give us the throughput of the database in terms of accesses/unit time.

Similarly the average number of readers in the system at steady state will be exactly the average marking of place  $P_6$ .

#### Other measures

At the GSPN level the states we are interested in can be identified either by whether a particular transition is enabled, or by whether a particular place is marked.

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To derive other performance measures we associate a value or reward with each of the markings we are interested in, just as we did when working directly at the Markov process level.

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At the GSPN level the states we are interested in can be identified either by whether a particular transition is enabled, or by whether a particular place is marked.

To derive other performance measures we associate a value or reward with each of the markings we are interested in, just as we did when working directly at the Markov process level.

For example, to derive the utilisation of the database in the reader-writer system, we associate a reward 1 with any marking in which transitions  $T_6$  or  $T_7$  are enabled, and a reward 0 with all other markings.

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### PIPE output for Reader-Writer example

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Source net										
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Results										
									Desults	
	GSPN Ste	au	уэ	tate	e A	nar	ysi	SR	tesuits	
Set of Tangible States										
		P0	P1	P2	Р3	Р4	P5	P6		
	MO	2	0	0	0	0	0	1		
	M1	1	0	0	0	0	1	0		
	M2	1	0	0	0	1	0	1		
	M3	0	0	1	0	0	1	0		
	M4	0	0	0	1	0	1	0		-
	M5	0	0	0	1	1	0	1		
	M6	0	0	0	0	2	0	1		
	Steady Stat	te D	istri	buti	on	of T	ang	liple	> States	
		'	Mark	ing	Va	lue				
		1	<u>0N</u>		0.1	9745	5			
		1	<u>M1</u>		0.1	5287	7			
		1	<u> 12</u>		0.0	3185	5			
		1	<u>N3</u>		0.3	0573	3			
		1	<u>14</u>		0.3	0573	3			
		1	<u>15</u>		0.0	0318	3			
		1	<u> 16</u>		0.0	0318	3			
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## PIPE output for Reader-Writer example

#### Average Number of Tokens on a Place

Place	Number of Tokens
P0	0.57962
P1	0
P2	0.30573
P3	0.30892
P4	0.0414
P5	0.76433
P6	0.23567

#### Token Probability Density

	μ=0	µ=1	µ=2
P0	0.61783	0.18471	0.19745
P1	1	0	0
P2	0.69427	0.30573	0
P3	0.69108	0.30892	0
P4	0.96178	0.03503	0.00318
P5	0.23567	0.76433	0
P6	0.76433	0.23567	0

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#### PIPE output for Reader-Writer example

#### Throughput of Timed Transitions

Transition	Throughput
т0	7.64331
Т5	3.82166
Т6	3.82166

#### Sojourn times for tangible states

Marking	Value
<u>M0</u>	0.05
<u>M1</u>	0.04
<u>M2</u>	0.00833
<u>M3</u>	0.2
<u>M4</u>	0.2
<u>M5</u>	0.01
<u>M6</u>	0.01

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#### Assumptions

Previous assumptions about Markov processes are still required but can now be interpreted at the GSPN level.

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Finite implies that the number of markings in the reachability set of a model (both tangible and vanishing markings) is finite. It can be shown that a GSPN is finite if it is bounded.

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Time homogeneity implies that the firing characteristics/system dynamics do not change over time. These characteristics are not necessarily static—marking dependent rates do vary—but the firing rate can depend only on the state of the model, not on how long it has been running.

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Time homogeneity implies that the firing characteristics/system dynamics do not change over time. These characteristics are not necessarily static—marking dependent rates do vary—but the firing rate can depend only on the state of the model, not on how long it has been running.

Irreducibility implies that it is possible to reach an arbitrary state from every other state.