

# Performance Modelling — Lecture 3

## Constructing and Solving Markov Processes

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23rd January 2017

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- A stochastic process is a set of random variables  $\{X(t), t \in T\}$ .
- $T$  is called the **index set** usually taken to represent time.
- Since we consider continuous time models  $T = \mathbb{R}^{\geq 0}$ , the set of non-negative real numbers.

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These paths are called **sample paths** or **realisations** of the stochastic process.

# Properties of Stochastic Processes

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$\{X(t)\}$  is a **Markov process**.

This implies that  $\{X(t)\}$  has the **Markov** or **memoryless property**: given the value of  $X(t)$  at some time  $t \in T$ , the future path  $X(s)$  for  $s > t$  does not depend on knowledge of the past history  $X(u)$  for  $u < t$ , i.e. for  $t_1 < \dots < t_n < t_{n+1}$ ,

$$\Pr(X(t_{n+1}) = x_{n+1} \mid X(t_n) = x_n, \dots, X(t_1) = x_1) = \Pr(X(t_{n+1}) = x_{n+1} \mid X(t_n) = x_n)$$

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$\{X(t)\}$  is **irreducible**.

This implies that all states in  $S$  can be reached from all other states, by following the transitions of the process. If we draw a directed graph of the state space with a node for each state and an arc for each event, or transition, then for any pair of nodes there is a path connecting them, i.e. the graph is strongly connected.

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$\{X(t)\}$  **is stationary**:

for any  $t_1, \dots, t_n \in T$  and  $t_1 + \tau, \dots, t_n + \tau \in T$  ( $n \geq 1$ ), then the process's joint distributions are unaffected by the change in the time axis and so,

$$F_{X(t_1+\tau)\dots X(t_n+\tau)} = F_{X(t_1)\dots X(t_n)}$$

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$\{X(t)\}$  **is time homogeneous:**

the behaviour of the system does not depend on **when** it is observed. In particular, the transition rates between states are independent of the time at which the transitions occur. Thus, for all  $t$  and  $s$ , it follows that

$$\Pr(X(t + \tau) = x_k \mid X(t) = x_j) = \Pr(X(s + \tau) = x_k \mid X(s) = x_j).$$

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Note: by the **Markov** property, the sojourn times are **memoryless**.

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Thus, for  $i \neq j, i, j \in S$ ,

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The **transition probability**  $p_{ij}$  is the probability, given that a transition out of state  $i$  occurs, that it is the transition to **state**  $j$ .  
By the definition of conditional probability, this is  $p_{ij} = q_{ij}/q_i$ .



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For a state space of size  $N$ , this is a  $N \times N$  matrix, where entry  $q(i, j)$  or  $q_{i,j}$ , records the transition rate of moving from state  $x_i$  to state  $x_j$ .

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By convention, the diagonal entries  $q_{i,i}$  are the negative row sum for row  $i$ , i.e.

$$q_{i,i} = - \sum_{j=1, j \neq i}^N q_{i,j}$$



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From this probability distribution we will derive performance measures based on subsets of states where some condition holds.

# Existence of a steady state probability distribution

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This distribution is reached when the initial state no longer has any influence.

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Thus, in an instant of time, the probability that a transition will occur from state  $x_i$  to state  $x_j$  is the probability that the model was in state  $x_i$ ,  $\pi_i$ , multiplied by the transition rate  $q_{ij}$ .

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This is called the **probability flux** from state  $x_i$  to state  $x_j$ .

## Global balance equations

In steady state, equilibrium is maintained so for any state the total probability flux out is equal to the total probability flux into the state.

$$\underbrace{\pi_i \times \sum_{x_j \in \mathcal{S}, j \neq i} q_{ij}}_{\text{flux out of } x_i} = \underbrace{\sum_{x_j \in \mathcal{S}, j \neq i} (\pi_j \times q_{ji})}_{\text{flux into } x_i}$$

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(If this were not true the distribution over states would change. )



## Global balance equations

Recall that the diagonal elements of the infinitesimal generator matrix  $\mathbf{Q}$  are the negative sum of the other elements in the row, i.e.  $q_{ii} = -\sum_{x_j \in \mathcal{S}, j \neq i} q_{ij}$ .

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Expressing the unknown values  $\pi_i$  as a row vector  $\boldsymbol{\pi}$ , we can write this as a matrix equation:

$$\boldsymbol{\pi} \mathbf{Q} = \mathbf{0}$$

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With these  $n + 1$  equations we can use standard linear algebra techniques to solve the equations and find the  $n$  unknowns,  $\{\pi_i\}$ .



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- The CPUs execute in private memory for a random time before issuing a common memory access request. Assume that this random time is exponentially distributed with parameter  $\lambda$ .
- The common memory access duration is also assumed to be exponentially distributed, with parameter  $\mu$  (the average duration of a common memory access is  $1/\mu$ ).

## Example

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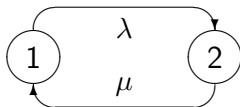
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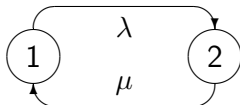
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The system behaviour can be modelled by a 2-state Markov process whose state transition diagram and generator matrix are as shown below:



$$\mathbf{Q} = \begin{pmatrix} -\lambda & \lambda \\ \mu & -\mu \end{pmatrix}$$

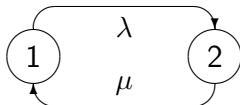
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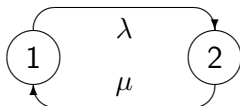


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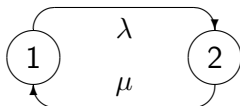
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Thus the steady state probability distribution is

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From this we can deduce, for example, that the probability that the processor is executing in private memory is  $\mu/(\mu + \lambda)$ .

## Solving the global balance equations

- In general the systems of equations will be too large to solve by hand.
- Instead we take advantage of linear algebra packages which can solve matrix equations of the form  $\mathbf{Ax} = \mathbf{b}$ .
- Here
  - $\mathbf{A}$  is an  $N \times N$  matrix,
  - $\mathbf{x}$  is a column vector of  $N$  unknowns, and
  - $\mathbf{b}$  is a column vector of  $N$  values.

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*This problem is resolved by transposing the equation, i.e.  $\mathbf{Q}^T \pi = 0$ , where the right hand side is now a column vector of zeros, rather than a row vector.*

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Now we can use any linear algebra solution package, such as MatLab to solve the resulting equation:

$$\mathbf{Q}_N^T \boldsymbol{\pi} = \mathbf{e}_N$$



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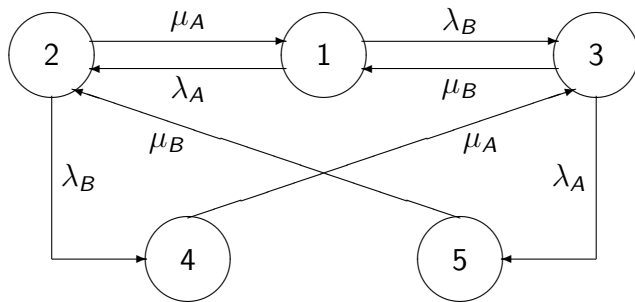
We assume that the processors have different timing characteristics, the private memory access of  $A$  being governed by an exponential distribution with parameter  $\lambda_A$ , the common memory access of  $B$  being governed by an exponential distribution with parameter  $\mu_B$ , etc.

## Example: state space

Now the state space becomes:

- 1  $A$  and  $B$  both executing in their private memories;
- 2  $B$  executing in private memory, and  $A$  accessing common memory;
- 3  $A$  executing in private memory, and  $B$  accessing common memory;
- 4  $A$  accessing common memory,  $B$  waiting for common memory;
- 5  $B$  accessing common memory,  $A$  waiting for common memory;

## Example: state space



## Example: generator matrix

$$\mathbf{Q} = \begin{pmatrix} -(\lambda_A + \lambda_B) & \lambda_A & \lambda_B & 0 & 0 \\ \mu_A & -(\mu_A + \lambda_B) & 0 & \lambda_B & 0 \\ \mu_B & 0 & -(\mu_B + \lambda_A) & 0 & \lambda_A \\ 0 & 0 & \mu_A & -\mu_A & 0 \\ 0 & \mu_B & 0 & 0 & -\mu_B \end{pmatrix}$$

## Example: modified generator matrix

$$\mathbf{Q}_N^T = \begin{pmatrix} -(\lambda_A + \lambda_B) & \mu_A & \mu_B & 0 & 0 \\ \lambda_A & -(\mu_A + \lambda_B) & 0 & 0 & \mu_B \\ \lambda_B & 0 & -(\mu_B + \lambda_A) & \mu_A & 0 \\ 0 & \lambda_B & 0 & -\mu_A & 0 \\ 1 & 1 & 1 & 1 & 1 \end{pmatrix}$$

## Example: steady state probability distribution

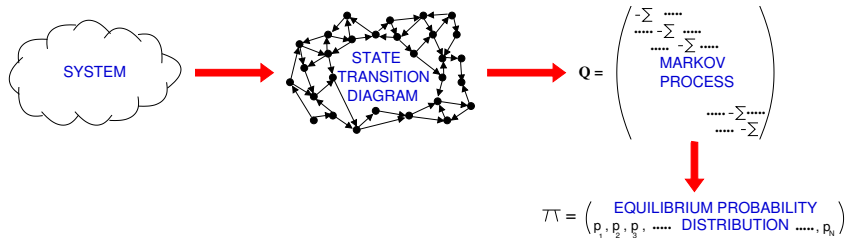
If we choose the following values for the parameters:

$$\lambda_A = 0.05 \quad \lambda_B := 0.1 \quad \mu_A = 0.02 \quad \mu_B = 0.05$$

solving the matrix equation, and rounding figures to 4 significant figures, we obtain:

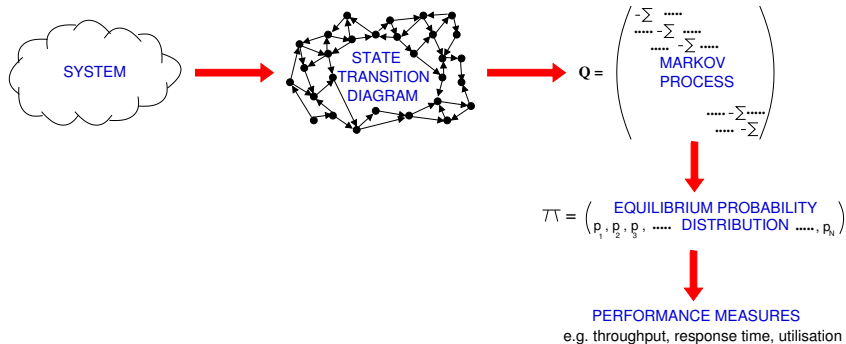
$$\boldsymbol{\pi} = (0.0693, 0.0990, 0.1683, 0.4951, 0.1683)$$

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- **state-based measures**, e.g. utilisation;
- **rate-based measures**, e.g. throughput;
- other measures which fall outside the above categories, e.g. **response time**.

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If we consider the multiprocessor example, the utilisation of the common memory,  $U_{mem}$ , is the total probability that the model is in one of the states in which the common memory is in use:

$$U_{mem} = \pi_2 + \pi_3 + \pi_4 + \pi_5 = 93.07\%$$



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For example, if we consider jobs waiting for the common memory to be queued in that subsystem, then the average number of jobs in the common memory,  $N_{mem}$ , is:

$$N_{mem} = (1 \times \pi_2) + (1 \times \pi_3) + (2 \times \pi_4) + (2 \times \pi_5) = 1.594$$

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**Rate-based measures** are those which correspond to the predicted rate at which some event occurs.

This will be the **product of the rate of the event, and the probability that the event is enabled**, i.e. the probability of being in one of the states from which the event can occur.

## Example: rate-based measures

In order to calculate the throughput of the common memory, we need the average number of accesses from either processor which it satisfies in unit time.

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$X_{mem}$  is thus calculated as:

$$X_{mem} = (\mu_A \times (\pi_2 + \pi_4)) + (\mu_B \times (\pi_3 + \pi_5)) = 0.0287$$

or, approximately one access every **35 milliseconds**.

## Other measures

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For example, applying Little's Law to the common memory we see that

$$W_{mem} = N_{mem}/X_{mem} = 1.594/0.0287 = 55.54 \text{ milliseconds}$$

# Assumptions

## Stochastic Hypothesis

*“The behaviour of a real system during a given period of time is characterised by the probability distributions of a stochastic process.”*

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- The Markov/memoryless assumption — future behaviour is only dependent on the current state, not on the past history — **is** a reasonable assumption for computer and communication systems, if we choose our states carefully.
- We generally assume that the Markov process is **finite, time homogeneous and irreducible**.

## Exercise

- Consider the multiprocessor example, but with three processors,  $A$ ,  $B$  and  $C$  sharing the common memory instead of two.
- List the states of the system, and draw the state transition diagram for this case.
- What is the difficulty in doing this and what further information do you need?
- Solution will be presented online later in the week.