

Logic Programming

Lecture 3: Recursion, lists, and data structures

Outline for today

- Recursion
 - proof search behavior
 - practical concerns
- List processing
- Programming with terms as data structures

Recursion

- So far we've (mostly) used *nonrecursive* rules
- These are limited:
 - not Turing-complete
 - can't define *transitive closure*
 - e.g. `ancestor`

Recursion (I)

- Nothing to it?
`ancestor(X,Y) :- parent(X,Y).`
`ancestor(X,Y) :- parent(X,Z),`
`ancestor(Z,Y).`
- Just use recursively defined predicate as a goal.
- Easy, right?

Depth first search, revisited

- Prolog tries rules **depth-first** in program order
 - **no matter what**
 - Even if there is an "obvious" solution using later clauses!
- ```
p :- p.
```
- ```
p.
```
- will **always** loop on first rule.

Recursion (2)

- Rule order can matter.

```
ancestor2(X,Y) :- parent(X,Z),  
                    ancestor2(Z,Y).  
  
ancestor2(X,Y) :- parent(X,Y).
```

This may be less efficient (tries to find **longest** path first)

This may also **loop** unnecessarily (if parent were cyclic).

Heuristic: write base case rules first.

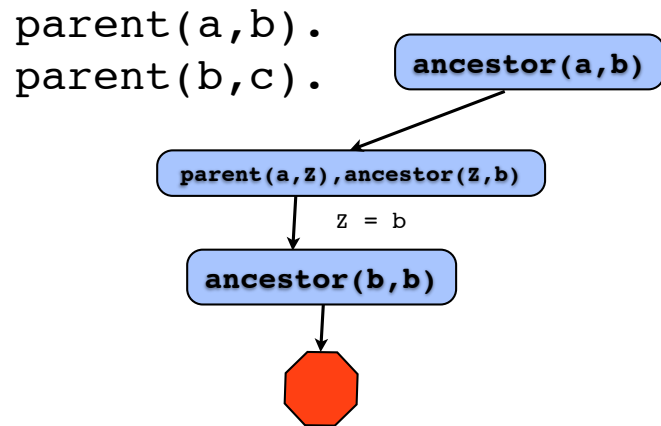
Rule order matters

```
parent(a,b).  
parent(b,c).  
ancestor(a,b)
```

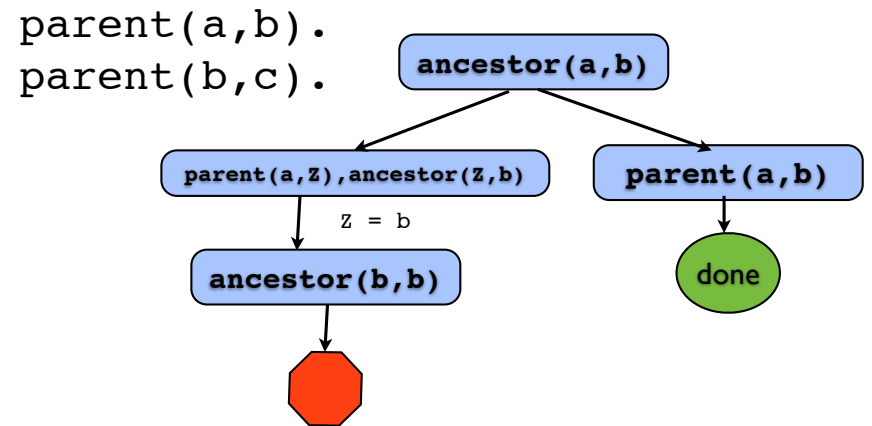
Rule order matters

```
parent(a,b).  
parent(b,c).  
ancestor(a,b)  
parent(a,Z), ancestor(Z,b)
```

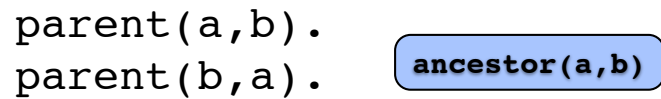
Rule order matters



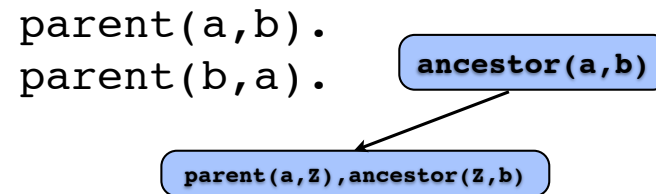
Rule order matters



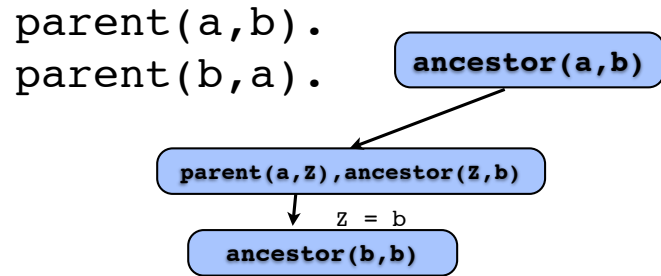
Rule order matters



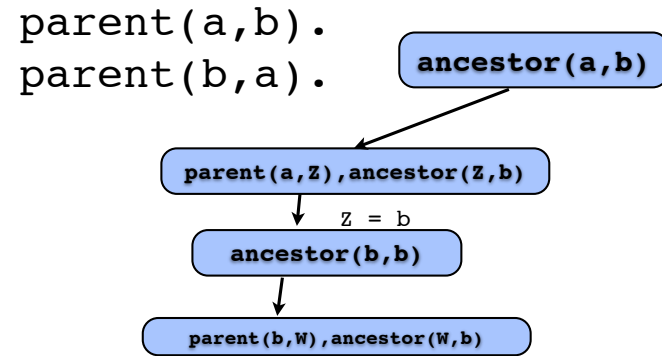
Rule order matters



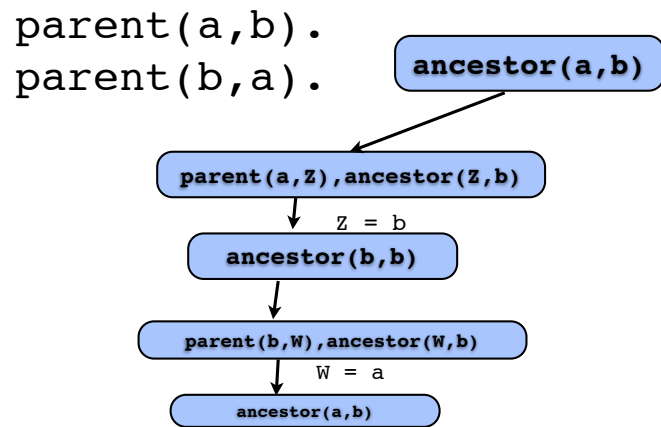
Rule order matters



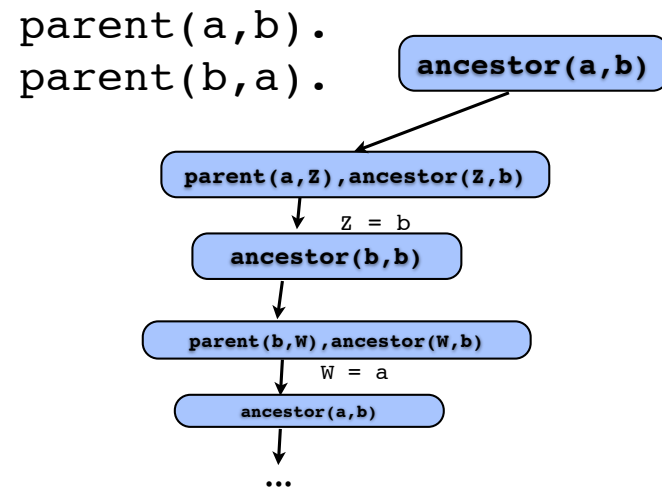
Rule order matters



Rule order matters



Rule order matters



Recursion (3)

- Goal order can matter.

```
ancestor3(X,Y) :- parent(X,Y).
```

```
ancestor3(X,Y) :- ancestor3(Z,Y),  
parent(X,Z).
```

This will list all solutions, then loop.

Goal order matters

```
parent(a,b).
```

```
parent(b,c).
```

ancestor(a,b)

Goal order matters

```
parent(a,b).
```

```
parent(b,c).
```

ancestor(a,b)

parent(a,b)

Goal order matters

```
parent(a,b).
```

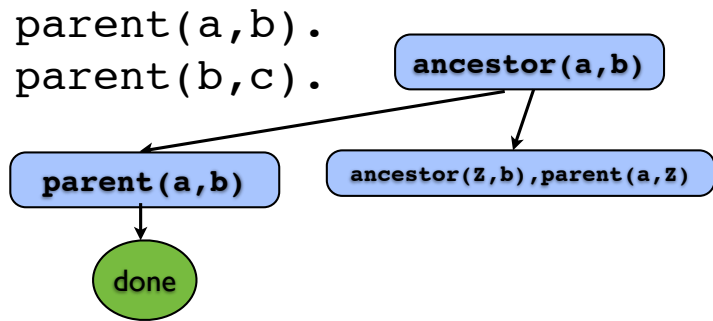
```
parent(b,c).
```

ancestor(a,b)

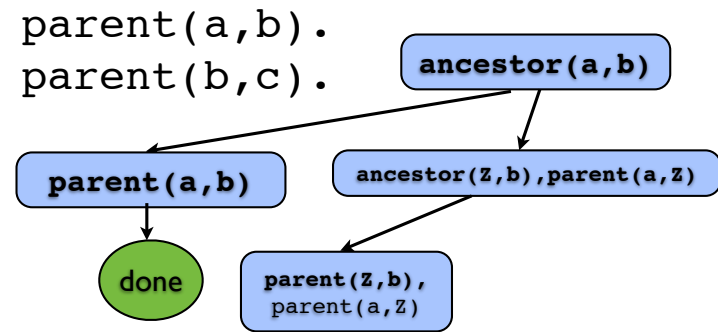
parent(a,b)

done

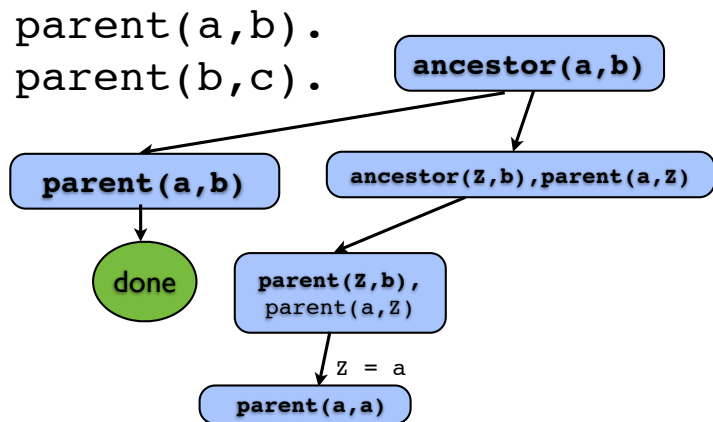
Goal order matters



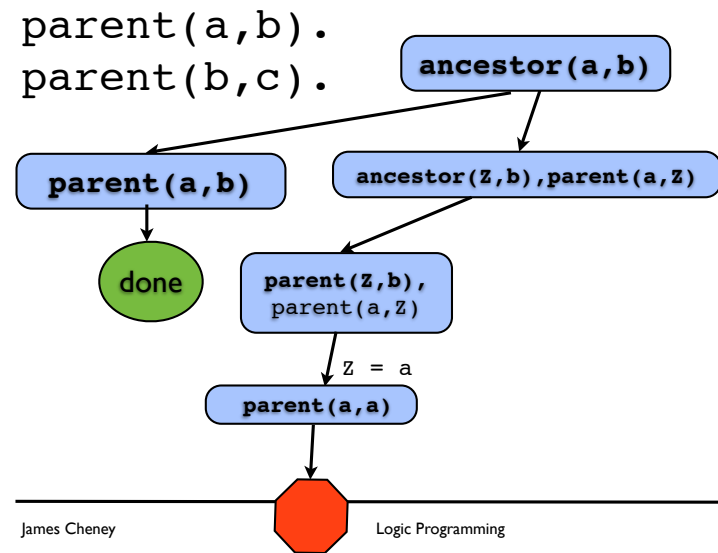
Goal order matters



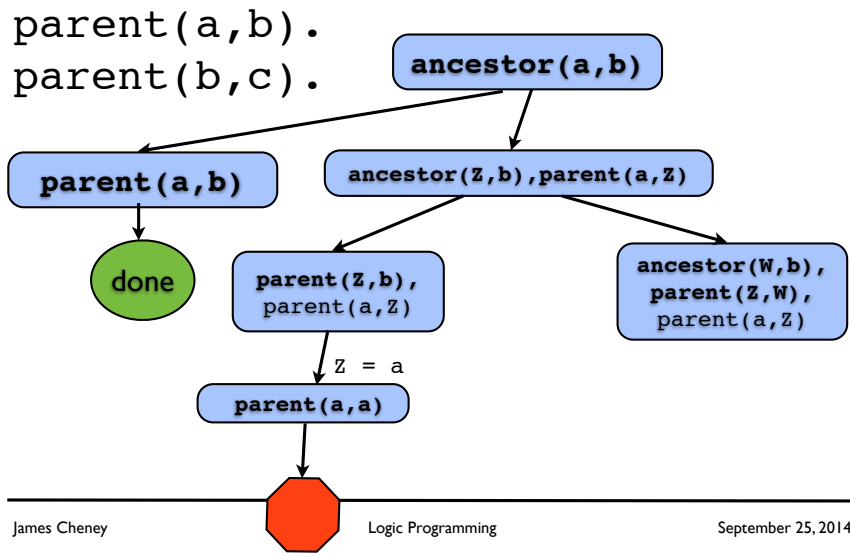
Goal order matters



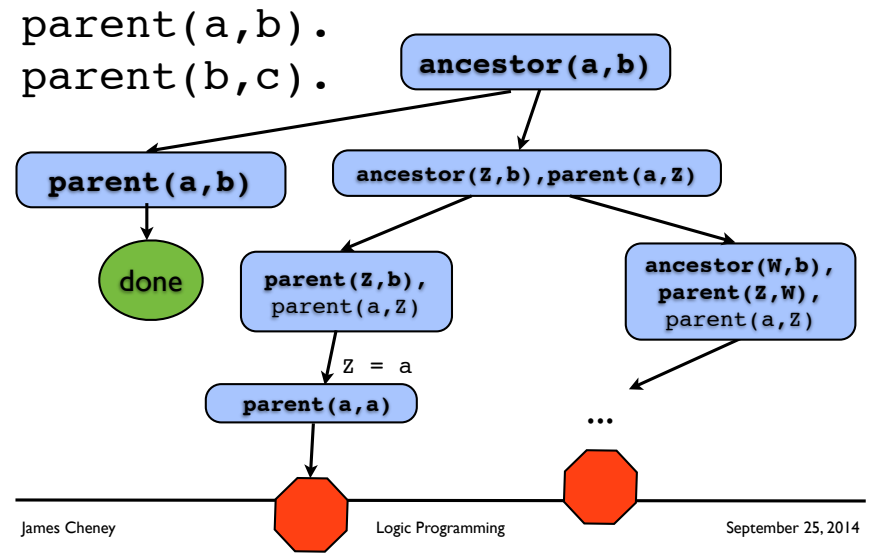
Goal order matters



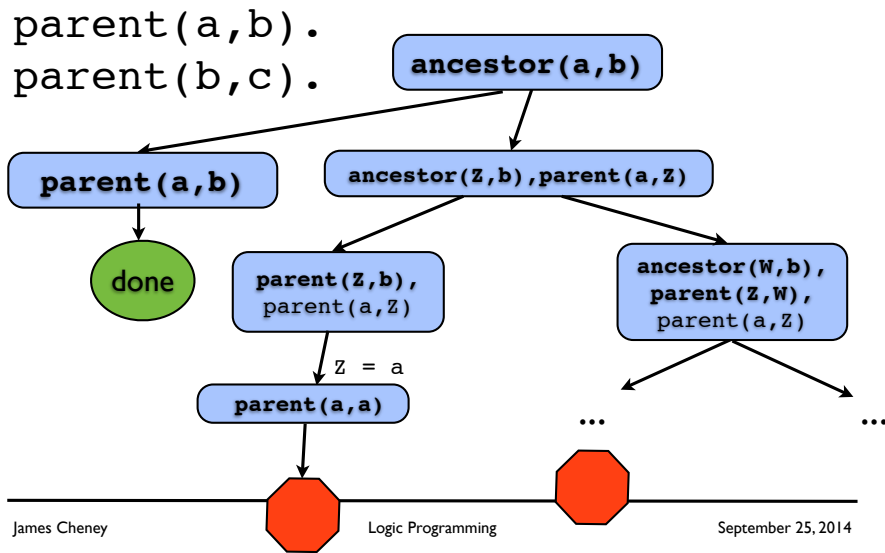
Goal order matters



Goal order matters



Goal order matters



Recursion (4)

- Goal order can matter.


```

ancestor4(X,Y) :- ancestor4(Z,Y),
                  parent(X,Z).
ancestor4(X,Y) :- parent(X,Y).
            
```

This will **always** loop!

Heuristic: try non-recursive goals first.

Goal order matters

`ancestor(X, Y)`

Goal order matters

`ancestor(X, Y)`

`ancestor(X, Z), parent(Z, Y)`

Goal order matters

`ancestor(X, Y)`

`ancestor(X, Z), parent(Z, Y)`

`ancestor(X, W), parent(W, Z), parent(Z, Y)`

Goal order matters

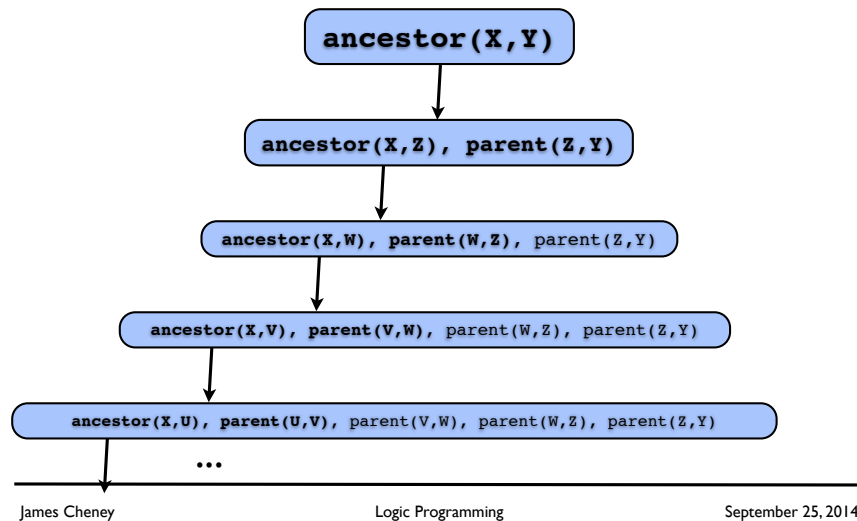
`ancestor(X, Y)`

`ancestor(X, Z), parent(Z, Y)`

`ancestor(X, W), parent(W, Z), parent(Z, Y)`

`ancestor(X, V), parent(V, W), parent(W, Z), parent(Z, Y)`

Goal order matters



Recursion and terms

- Terms can be **arbitrarily nested**
- **Example:** unary natural numbers
`nat(z).`
`nat(s(N)) :- nat(N).`
- To do interesting things we need recursion

Addition and subtraction

- **Example:** addition
`add(z, N, N).`
`add(s(N), M, s(P)) :- add(N, M, P).`
- Can run in reverse to find all M, N with M+N=P
- Can use to define `leq`
`leq(M, N) :- add(M, _, N).`

Multiplication

- Can multiply two numbers:
`multiply(z, N, z).`
`multiply(s(N), M, P) :-`
 `multiply(N, M, Q), add(M, Q, P).`
`square(M) :- multiply(N, N, M).`

List processing

- Recall built-in **list syntax**

```
list([]).
```

```
list([X|L]) :- list(L).
```

- **Example:** list append

```
append([],L,L).
```

```
append([X|L],M,[X|N]) :- append(L,M,N).
```

Append in action

- Forward direction

```
?- append([1,2],[3,4],X).
```

- Backward direction

```
?- append(X,Y,[1,2,3,4]).
```

Mode annotations

- Notation `append(+,+, -)`
 - "if you call `append` with first two arguments ground then it will make the third argument ground"
- Similarly, `append(-,-,+)`
 - "if you call `append` with last argument ground then it will make the first two arguments ground"
- Not "code", but often used in documentation
 - "?" annotation means either + or -

List processing (2)

- When is something a member of a list?

```
mem(X,[X|_]).
```

```
mem(X,[_|L]) :- mem(X,L).
```

- Typical modes

```
mem(+,+)
```

```
mem(-,+)
```

List processing(3)

- Removing an element of a list

```
remove(X, [X|L], L).
```

```
remove(X, [Y|L], [Y|M]) :- remove(X, L, M).
```

- Typical mode

```
remove(+, +, -)
```

List processing (4)

- Zip, or "pairing" corresponding elements of two lists

```
zip([], [], []).
```

```
zip([X|L], [Y|M], [(X,Y)|N]) :- zip(L, M, N).
```

- Typical modes:

```
zip(+, +, -).
```

```
zip(-, -, +). % "unzip"
```

List flattening

- Write **flatten** predicate `flatten/2`
 - Given a list of (lists of ..) lists
 - Produces a list containing all elements in order

List flattening

- Write **flatten** predicate `flatten/2`
 - Given a list of (lists of ..) lists
 - Produces a list containing all elements in order

```
flatten([], []).
```

```
flatten([X|L], M) :- flatten(X, Y1),
```

```
flatten(L, Y2),
```

```
append(Y1, Y2, M).
```

```
flatten(X, [X]) :- ???.
```

List flattening

- Write **flatten** predicate `flatten/2`
 - Given a list of (lists of ..) lists
 - Produces a list containing all elements in order

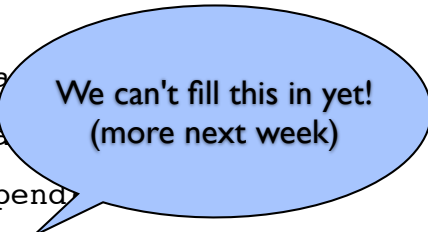
```
flatten([], []).
```

```
flatten([X|L],M) :- fla
```

```
fla
```

```
append
```

```
flatten(X,[X]) :- ???.
```



We can't fill this in yet!
(more next week)

Records/structs

- We can use terms to define data structures
`pb([entry(james, '123-4567'), ...])`
- and operations on them

```
pb_lookup(pb(B),P,N) :-
```

```
member(entry(P,N),B).
```

```
pb_insert(pb(B),P,N,pb([entry(P,N)|B])).
```

```
pb_remove(pb(B),P,pb(B2)) :-
```

```
remove(entry(P,_),B,B2).
```

Trees

- We can define (binary) trees with data:

```
tree(leaf).
```

```
tree(node(X,T,U)) :- tree(T), tree(U).
```

Tree membership

- Define **membership in tree**

```
mem_tree(X,node(X,T,U)).
```

```
mem_tree(X,node(Y,T,U)) :-
```

```
mem_tree(X,T) ;
```

```
mem_tree(X,U).
```

Preorder traversal

- Define **preorder**

```
preorder(leaf, []).
```

```
preorder(node(X,T,U), [X|N]) :-
```

```
    preorder(T,L),
```

```
    preorder(U,M),
```

```
    append(L,M,N).
```

- What happens if we run this in reverse?

Next time

- Nonlogical features
 - Expression evaluation
 - I/O
 - "Cut" (pruning proof search)
- Further reading:
 - Learn Prolog Now, ch. 3-4