## **Information Theory** — **Tutorial 2**

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- 1. Entropy separability: MacKay's book Exercise 4.2, p68.
- 2. Entropy decomposition: MacKay's book Exercise 2.28, p38.
- 3. **Symbol code statistics:** MacKay's book Exercise 5.31, p104. (The ensemble *X* and code  $C_3$  were defined on pp. 92–93: symbols with probabilities  $\mathcal{P}_X = \{1/2, 1/4, 1/8, 1/8\}$  are encoded with codewords  $C_3 = \{0, 10, 110, 111\}$ .)

## 4. Three questions on Huffman coding:

MacKay's book Exercise 5.21, p102. Notation: "codes for  $X^n$ " means codes for blocks of *n* concatenated symbols. For example, when coding  $X^2$  the *source* alphabet is  $A_{X^2} = \{00, 01, 10, 11\}$ .

MacKay Exercise 5.26, p103.

MacKay Exercise 5.29, p103.

5. **Real-valued variables** (*Optional*): We have focussed on discrete random variables, X, taking on values  $\{x_i\}_{i=1}^{I}$ . When compressing many observations of such a variable, the number of bits/symbol required is given by the entropy:

$$H(X) = \sum_{i} P(x_i) \log \frac{1}{P(x_i)}$$

How many bits/outcome are required to encode draws from a unit Gaussian distributed variable, which has probability density  $p(x) = \exp(-x^2/2)/\sqrt{2\pi}$ ? How many bits on average would be required to store the answers to 3 decimal places?