

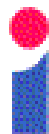


First-Order Logic

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Outline

- Why FOL?
- Syntax and semantics of FOL
- Using FOL
- Wumpus world in FOL

Pros and cons of propositional logic

- Propositional logic is declarative
- Propositional logic allows partial/disjunctive/negated information
 - (unlike most data structures and databases)
- Propositional logic is compositional:
 - The meaning of $B_{1,1} \wedge P_{1,2}$ is derived from that of $B_{1,1}$ and of $P_{1,2}$
- Meaning in propositional logic is context-independent
 - (unlike natural language, where meaning depends on context)
- Propositional logic has very limited expressive power
 - (unlike natural language)
 - for example, we cannot say "pits cause breezes in adjacent squares", except by writing one sentence for **each** square

First-order logic

Whereas propositional logic assumes the world contains facts, first-order logic (like natural language) assumes the world contains:

- **Objects:** people, houses, numbers, colours, football games, wars, ...
- **Relations:** red, round, prime, brother of, bigger than, part of, comes between, ...
- **Functions:** father of, best friend, one more than, plus, ...

Syntax of FOL: Basic elements

- Constants *KingJohn, 2, UoE,...*
- Predicates *Brother, >,...*
- Functions *Sqrt, LeftLegOf,...*
- Variables *x, y, a, b,...*
- Connectives $\neg, \Rightarrow, \wedge, \vee, \Leftrightarrow$
- Equality $=$
- Quantifiers \forall, \exists

Atomic formulae

Atomic formula = *predicate* (*term*₁, ..., *term*_n)
or *term*₁ = *term*₂

Term = *function* (*term*₁, ..., *term*_n)
or *constant* or *variable*

Examples:

- *Brother*(*KingJohn*, *RichardTheLionheart*)
- *>*(*Length*(*LeftLegOf*(*Richard*)), *Length*(*LeftLegOf*(*KingJohn*)))

predicate

functions

constants

Complex formulae

Complex formulae are made from atomic formulae using connectives

$$\neg S, S_1 \wedge S_2, S_1 \vee S_2, S_1 \Rightarrow S_2, S_1 \Leftrightarrow S_2$$

Examples:

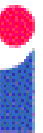
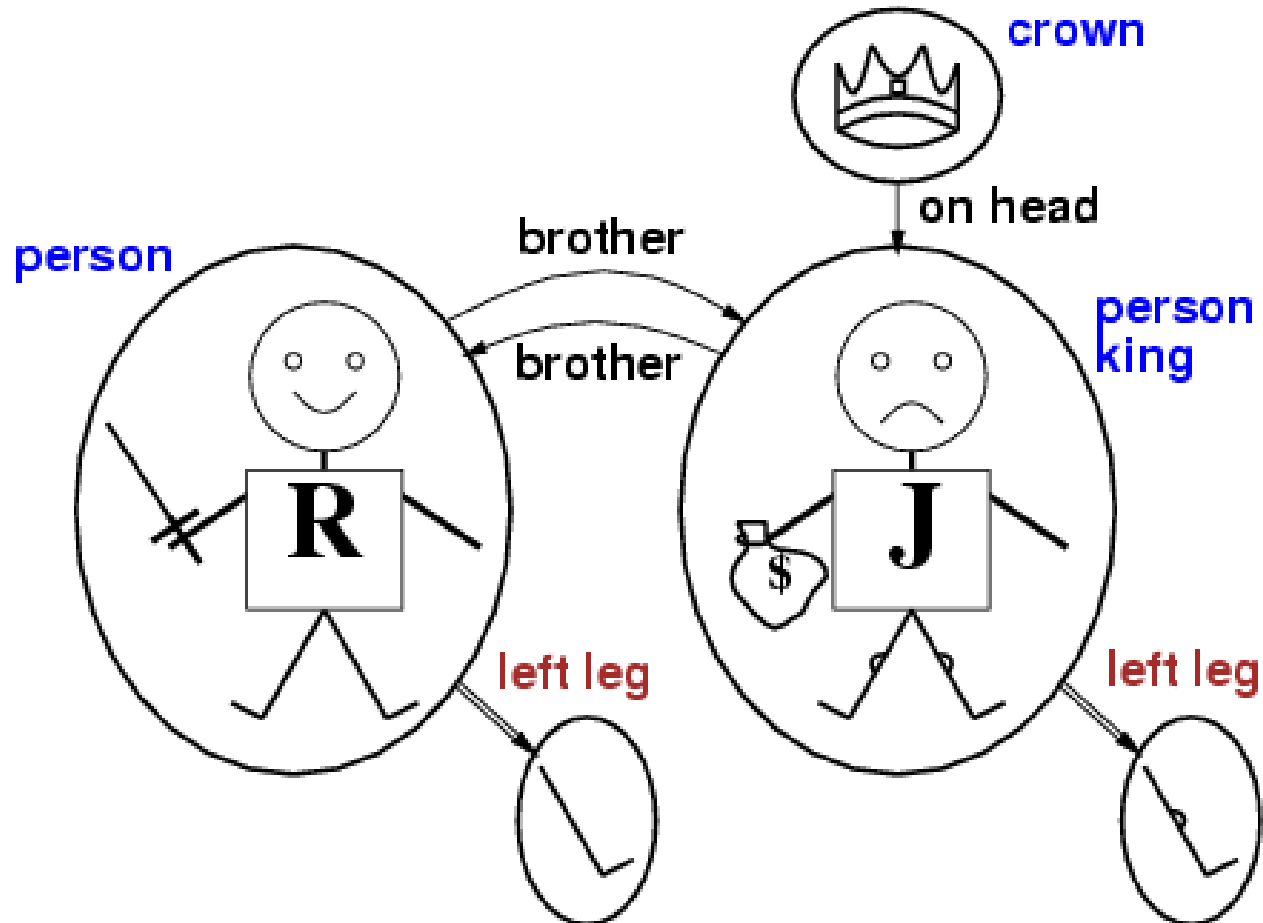
- $Sibling(KingJohn, Richard) \Rightarrow Sibling(Richard, KingJohn)$
- $>(1,2) \vee \leq(1,2)$
- $>(1,2) \wedge \neg >(1,2)$

Semantics of first-order logic

- Formulae are mapped to an **interpretation**
 - An interpretation is called a **model** of a set of formulae when all the formulae are **true** in the interpretation.
- Interpretation contains objects (**domain elements**) and relations between them
- Mapping specifies referents for

constant symbols	↦	objects
predicate symbols	↦	relations
function symbols	↦	functions
- An atomic formula ***predicate***(***term*₁, ..., *term*_{*n*}) is true iff the **objects** referred to by ***term*₁, ..., *term*_{*n*}** are in the **relation** referred to by ***predicate***.**

Interpretations for FOL: Example



Universal quantification

- $\forall \langle \text{variables} \rangle. \langle \text{formula} \rangle$
 - But will often write $\forall x, y. P$ for $\forall x. \forall y.$
 - Example: Everyone at UoE is smart: $\forall x. \text{At}(x, \text{UoE}) \Rightarrow \text{Smart}(x)$
- $\forall x. P$ is **true** in an interpretation m iff P is **true** with x being **each** possible object in the interpretation.
- Roughly speaking, equivalent to the **conjunction** of **instantiations** of P

$$\begin{aligned} & \text{At}(\text{KingJohn}, \text{UoE}) \Rightarrow \text{Smart}(\text{KingJohn}) \\ \wedge & \text{At}(\text{Richard}, \text{UoE}) \Rightarrow \text{Smart}(\text{Richard}) \\ \wedge & \text{At}(\text{UoE}, \text{UoE}) \Rightarrow \text{Smart}(\text{UoE}) \\ \wedge & \dots \end{aligned}$$

A common mistake to avoid

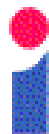
- Typically, \Rightarrow is the main connective with \forall
- Common mistake: using \wedge as the main connective with \forall :

$$\forall x. \text{At}(x, \text{UoE}) \wedge \text{Smart}(x)$$

means “Everyone is at UoE and everyone is smart”

Existential quantification

- $\exists \langle \text{variables} \rangle. \langle \text{formula} \rangle$
 - But will often write $\exists x, y. P$ for $\exists x. \exists y. P$
 - Example: Someone at UoE is smart: $\exists x. \text{At}(x, \text{UoE}) \wedge \text{Smart}(x)$
- $\exists x. P$ is **true** in an interpretation m iff P is **true** with x being **some** possible object in the model.
- Roughly speaking, equivalent to the **disjunction** of **instantiations** of P
 - At(KingJohn, UoE) \wedge Smart(KingJohn)
 - ✓ At(Richard, UoE) \wedge Smart(Richard)
 - ✓ At(UoE, UoE) \wedge Smart(UoE)
 - ✓ ...



Another common mistake to avoid

- Typically, \wedge is the main connective with \exists
- Common mistake: using \Rightarrow as the main connective with \exists :

$$\exists x. \text{At}(x, \text{UoE}) \Rightarrow \text{Smart}(x)$$

is true if there is anyone who is not at UoE!

Properties of quantifiers

- $\forall x. \forall y.$ is the same as $\forall y. \forall x.$
- $\exists x. \exists y.$ is the same as $\exists y. \exists x.$
- $\exists x. \forall y.$ is **not** the same as $\forall y. \exists x.$
 - $\exists x. \forall y. \text{Loves}(x,y)$

“There is a person who loves everyone in the world”
 - $\forall y. \exists x. \text{Loves}(x,y)$

“Everyone in the world is loved by at least one person”
- **Quantifier duality**: each can be expressed using the other:
 - $\forall x. \text{Likes}(x, \text{IceCream}) \equiv \neg \exists x. \neg \text{Likes}(x, \text{IceCream})$
 - $\exists x. \text{Likes}(x, \text{Broccoli}) \equiv \neg \forall x. \neg \text{Likes}(x, \text{Broccoli})$

Equality

- $term_1 = term_2$ is true under a given interpretation **if and only if** $term_1$ and $term_2$ refer to the same object.

- Example. Definition of *Sibling* in terms of *Parent*:

$$\forall x,y. \textit{Sibling}(x,y) \Leftrightarrow (\neg(x = y) \wedge \exists m,f. \neg(m = f) \wedge \textit{Parent}(m,x) \wedge \textit{Parent}(f,x) \wedge \textit{Parent}(m,y) \wedge \textit{Parent}(f,y))$$

Using FOL

Example: The kinship domain:

- Brothers are siblings

$$\forall x,y. \textit{Brother}(x,y) \Rightarrow \textit{Sibling}(x,y)$$

- One's mother is one's female parent

$$\forall m,c. \textit{Mother}(c) = m \Leftrightarrow (\textit{Female}(m) \wedge \textit{Parent}(m,c))$$

- “Sibling” is symmetric

$$\forall x,y. \textit{Sibling}(x,y) \Leftrightarrow \textit{Sibling}(y,x)$$

Using FOL

The set domain:

- $\forall s. \text{Set}(s) \Leftrightarrow (s = \{\}) \vee (\exists x, s_2. \text{Set}(s_2) \wedge s = \{x|s_2\})$
- $\neg \exists x, s. \{x|s\} = \{\}$
- $\forall x, s. x \in s \Leftrightarrow s = \{x|s\}$
- $\forall x, s. x \in s \Leftrightarrow [\exists y, s_2. (s = \{y|s_2\} \wedge (x = y \vee x \in s_2))]$
- $\forall s_1, s_2. s_1 \subseteq s_2 \Leftrightarrow (\forall x. x \in s_1 \Rightarrow x \in s_2)$
- $\forall s_1, s_2. (s_1 = s_2) \Leftrightarrow (s_1 \subseteq s_2 \wedge s_2 \subseteq s_1)$
- $\forall x, s_1, s_2. x \in (s_1 \cap s_2) \Leftrightarrow (x \in s_1 \wedge x \in s_2)$
- $\forall x, s_1, s_2. x \in (s_1 \cup s_2) \Leftrightarrow (x \in s_1 \vee x \in s_2)$

Interacting with FOL KBs

- Suppose a wumpus-world agent is using an FOL KB and perceives a smell and a breeze (but no glitter) at $t=5$:

`Tell(KB, Percept([Smell,Breeze,None], 5))`

`Ask(KB, $\exists a$. BestAction(a, 5))`

- i.e., does the KB entail some best action at $t=5$?
- Answer: Yes, $\{a/Shoot\}$ ← substitution (binding list)

- Given a sentence S and a substitution σ ,
 - $S\sigma$ denotes the result of “plugging” σ into S ; e.g.,

$S = \text{Smarter}(x, y)$

$\sigma = \{x/Obama, y/Palin\}$

$S\sigma = \text{Smarter}(Obama, Palin)$

- `Ask(KB, S)` returns some/all σ such that $KB \models S\sigma$

Knowledge base for the wumpus world

- Perception

- $\forall t, s, b. \text{Percept}([s, b, \text{Glitter}], t) \Rightarrow \text{Glitter}(t)$

- Reflex

- $\forall t. \text{Glitter}(t) \Rightarrow \text{BestAction}(\text{Grab}, t)$

Deducing hidden properties

- $\forall x,y,a,b. \text{Adjacent}([x,y],[a,b]) \Leftrightarrow [a,b] \in \{[x+1,y], [x-1,y],[x,y+1],[x,y-1]\}$
- $\forall s,t. \text{At}(\text{Agent},s,t) \wedge \text{Breeze}(t) \Rightarrow \text{Breezy}(s)$
- Squares are breezy near a pit:
 - **Diagnostic** rule: infer cause from effect
 $\forall s. \text{Breezy}(s) \Rightarrow \exists r. \text{Adjacent}(r,s) \wedge \text{Pit}(r)$
 - **Causal** rule: infer effect from cause
 $\forall r. \text{Pit}(r) \Rightarrow [\forall s. \text{Adjacent}(r,s) \Rightarrow \text{Breezy}(s)]$

Summary

- First-order logic:
 - objects and relations are semantic primitives
 - syntax: constants, functions, predicates, equality, quantifiers
- Increased expressive power: sufficient to define wumpus world