

# Informatics 2D – Reasoning and Agents

Semester 2, 2019–2020

Alex Lascarides  
 alex@inf.ed.ac.uk



Lecture 24 – Exact Inference in Bayesian Networks  
 13th March 2020

## Inference in BNs

- ▶ Basic task: compute posterior distribution for set of **query variables** given some observed **event** (i.e. assignment of values to **evidence variables**)
- ▶ Formally: determine  $\mathbf{P}(X|\mathbf{e})$  given query variables  $\mathbf{X}$ , evidence variables  $\mathbf{E}$  (and non-evidence or **hidden** variables  $\mathbf{Y}$ )
- ▶ Example:  
 $\mathbf{P}(\text{Burglary} | \text{JohnCalls} = \text{true}, \text{MaryCalls} = \text{true}) = \langle 0.284, 0.716 \rangle$
- ▶ First we will discuss exact algorithms for computing posterior probabilities then approximate methods later

## Where are we?

Last time ...

- ▶ Introduced Bayesian networks
- ▶ Allow for compact representation of JPDs
- ▶ Methods for efficient representations of CPTs
- ▶ But how hard is inference in BNs?

Today ...

- ▶ **Inference in Bayesian networks**

## Inference by enumeration

- ▶ We have seen that any conditional probability can be computed from a full JPD by summing terms
- ▶  $\mathbf{P}(X|\mathbf{e}) = \alpha \mathbf{P}(X, \mathbf{e}) = \alpha \sum_{\mathbf{y}} \mathbf{P}(X, \mathbf{e}, \mathbf{y})$
- ▶ Since BN gives complete representation of full JPD, we must be able to answer a query by computing sums of products of conditional probabilities from the BN
- ▶ Consider query  
 $\mathbf{P}(\text{Burglary} | \text{JohnCalls} = \text{true}, \text{MaryCalls} = \text{true}) = \mathbf{P}(B|j, m)$
- ▶  $\mathbf{P}(B|j, m) = \alpha \mathbf{P}(B, j, m) = \alpha \sum_{\mathbf{e}} \sum_{\mathbf{a}} \mathbf{P}(B, \mathbf{e}, \mathbf{a}, j, m)$

## Inference by enumeration

- Recall  $P(x_1, \dots, x_n) = \prod_{i=1}^n P(x_i | \text{parents}(X_i))$
- We can use CPTs to simplify this exploiting BN structure
- For  $Burglary = true$ :

$$P(b|j, m) = \alpha \sum_e \sum_a P(b)P(e)P(a|b, e)P(j|a)P(m|a)$$

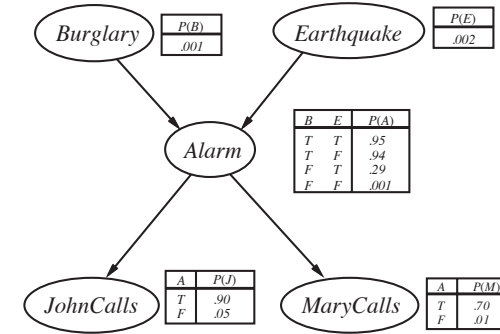
- But we can improve efficiency of this by moving terms outside that don't depend on sums

$$P(b|j, m) = \alpha P(b) \sum_e P(e) \sum_a P(a|b, e)P(j|a)P(m|a)$$

- To compute this, we need to loop through variables in order and multiply CPT entries; for each summation we need to loop over variable's possible values

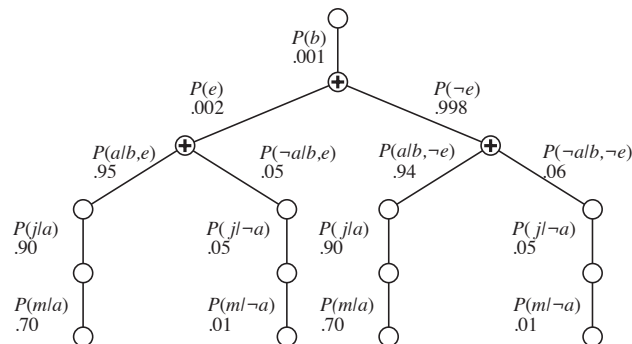
## Example

- New burglar alarm has been fitted, fairly reliable but sometimes reacts to earthquakes
- Neighbours John and Mary promise to call when they hear alarm
- John sometimes mistakes phone for alarm, and Mary listens to loud music and sometimes doesn't hear alarm



## The variable elimination algorithm

- Enumeration method is computationally quite hard.
- You often compute the same thing several times; e.g.  $P(j|a)P(m|a)$  and  $P(j|\neg a)P(m|\neg a)$  for each value of  $e$
- Evaluation of expression shown in the following tree:



## The variable elimination algorithm

- Idea of **variable elimination**: avoid repeated calculations
- Basic idea: store results after doing calculation once
- Works bottom-up by evaluating subexpressions
- Assume we want to evaluate

$$P(B|j, m) = \alpha \underbrace{P(B)}_{f_1(B)} \sum_e \underbrace{P(e)}_{f_2(E)} \sum_a \underbrace{P(a|B, e)}_{f_3(A, B, E)} \underbrace{P(j|a)}_{f_4(A)} \underbrace{P(m|a)}_{f_5(A)}$$

- We've annotated each part with a **factor**.
- A factor is a **matrix**, indexed with its argument variables. E.g:
  - Factor  $f_5(A)$  corresponds to  $P(m|a)$  and depends just on  $A$  because  $m$  is fixed (it's a  $2 \times 1$  matrix).

$$f_5(A) = \langle P(m|a), P(m|\neg a) \rangle$$

- $f_3(A, B, E)$  is a  $2 \times 2 \times 2$  matrix for  $P(a|B, e)$

## The variable elimination algorithm

$$\mathbf{P}(B|j, m) = \alpha \mathbf{f}_1(B) \times \sum_e \mathbf{f}_2(E) \sum_a \mathbf{f}_3(A, B, E) \times \mathbf{f}_4(A) \times \mathbf{f}_5(A)$$

- Summing out  $A$  produces a  $2 \times 2$  matrix (via **pointwise product**):

$$\begin{aligned} \mathbf{f}_6(B, E) &= \sum_a \mathbf{f}_3(A, B, E) \times \mathbf{f}_4(A) \times \mathbf{f}_5(A) \\ &= (\mathbf{f}_3(a, B, E) \times \mathbf{f}_4(a) \times \mathbf{f}_5(a)) + \\ &\quad (\mathbf{f}_3(\neg a, B, E) \times \mathbf{f}_4(\neg a) \times \mathbf{f}_5(\neg a)) \end{aligned}$$

- So now we have  $\mathbf{P}(B|j, m) = \alpha \mathbf{f}_1(B) \times \sum_e \mathbf{f}_2(E) \times \mathbf{f}_6(B, E)$
- Sum out  $E$  in the same way:  $\mathbf{f}_7(B) = (\mathbf{f}_2(e) \times \mathbf{f}_6(B, e)) + (\mathbf{f}_2(\neg e) \times \mathbf{f}_6(B, \neg e))$
- Using  $\mathbf{f}_1(B) = \mathbf{P}(B)$ , we can finally compute

$$\mathbf{P}(B|j, m) = \alpha \mathbf{f}_1(B) \times \mathbf{f}_7(B)$$

- Remains to define pointwise product and summing out

## An example

- Pointwise product yields product for union of variables in its arguments:

$$\mathbf{f}(X_1 \dots X_i, Y_1 \dots Y_j, Z_1 \dots Z_k) = \mathbf{f}_1(X_1 \dots X_i, Y_1 \dots Y_j) \mathbf{f}_2(Y_1 \dots Y_j, Z_1 \dots Z_k)$$

A	B	$\mathbf{f}_1(A, B)$	B	C	$\mathbf{f}_2(B, C)$	A	B	C	$\mathbf{f}(A, B, C)$
T	T	0.3	T	T	0.2	T	T	T	$0.3 \times 0.2$
T	F	0.7	T	F	0.8	T	T	F	$0.3 \times 0.8$
F	T	0.9	F	T	0.6	T	F	T	$0.7 \times 0.6$
F	F	0.1	F	F	0.4	T	F	F	$0.7 \times 0.4$
						F	T	T	$0.9 \times 0.2$
						F	T	F	$0.9 \times 0.8$
						F	F	T	$0.1 \times 0.6$
						F	F	F	$0.1 \times 0.4$

- For example  $\mathbf{f}(T, T, F) = \mathbf{f}_1(T, T) \times \mathbf{f}_2(T, F)$

## An example

- Summing out is similarly straightforward
- Trick: any factor that does not depend on the variable to be summed out can be moved outside the summation process
- For example

$$\begin{aligned} &\sum_e \mathbf{f}_2(E) \times \mathbf{f}_3(A, B, E) \times \mathbf{f}_4(A) \times \mathbf{f}_5(A) \\ &= \mathbf{f}_4(A) \times \mathbf{f}_5(A) \times \sum_e \mathbf{f}_2(E) \times \mathbf{f}_3(A, B, E) \end{aligned}$$

- Matrices are only multiplied when we need to sum out a variable from the accumulated product

## Another Example: $\mathbf{P}(J|b) = \langle P(j|b), P(\neg j|b) \rangle$

$$\begin{aligned} \mathbf{P}(J|b) &= \alpha \sum_e \sum_a \sum_m \mathbf{P}(J, b, e, a, m) && \text{prod., marg.} \\ &= \alpha \sum_e \sum_a \sum_m P(b)P(e)P(a|b, e)\mathbf{P}(J|a)P(m|a) && \text{cond. indep.} \\ &= \alpha' \sum_e \underbrace{P(e)}_{\mathbf{f}_1(E)} \sum_a \underbrace{P(a|b, e)}_{\mathbf{f}_2(A, E)} \underbrace{\mathbf{P}(J|a)}_{\mathbf{f}_3(J, A)} \underbrace{\sum_m P(m|a)}_{=1} && \text{move terms} \\ &= \alpha' \sum_e \mathbf{f}_1(E) \sum_a \mathbf{f}_2(A, E) \mathbf{f}_3(J, A) \\ &= \alpha' \sum_e \mathbf{f}_1(E) \mathbf{f}_4(J, E) \\ &= \alpha' \mathbf{f}_5(J) \end{aligned}$$

Can eliminate all variables that aren't ancestors of query or evidence variables!

## Summary

- ▶ Inference in Bayesian Networks
- ▶ Exact methods: enumeration, variable elimination algorithm
- ▶ Computationally intractable in the worst case
- ▶ Next time: **Approximate inference in Bayesian Networks**