

Lecture 2: Data Representation

- Data need to be represented in a convenient way that simplifies:
 - common operations: addition, comparison, etc.
 - How would an algorithm for adding Roman numbers look like?
 - hardware implementation (cheap, fast, reliable computers)
- Data representation is a major part of the software-hardware interface



Lecture outline

- The bit – atomic unit of data
- Representing numbers
- Representing text



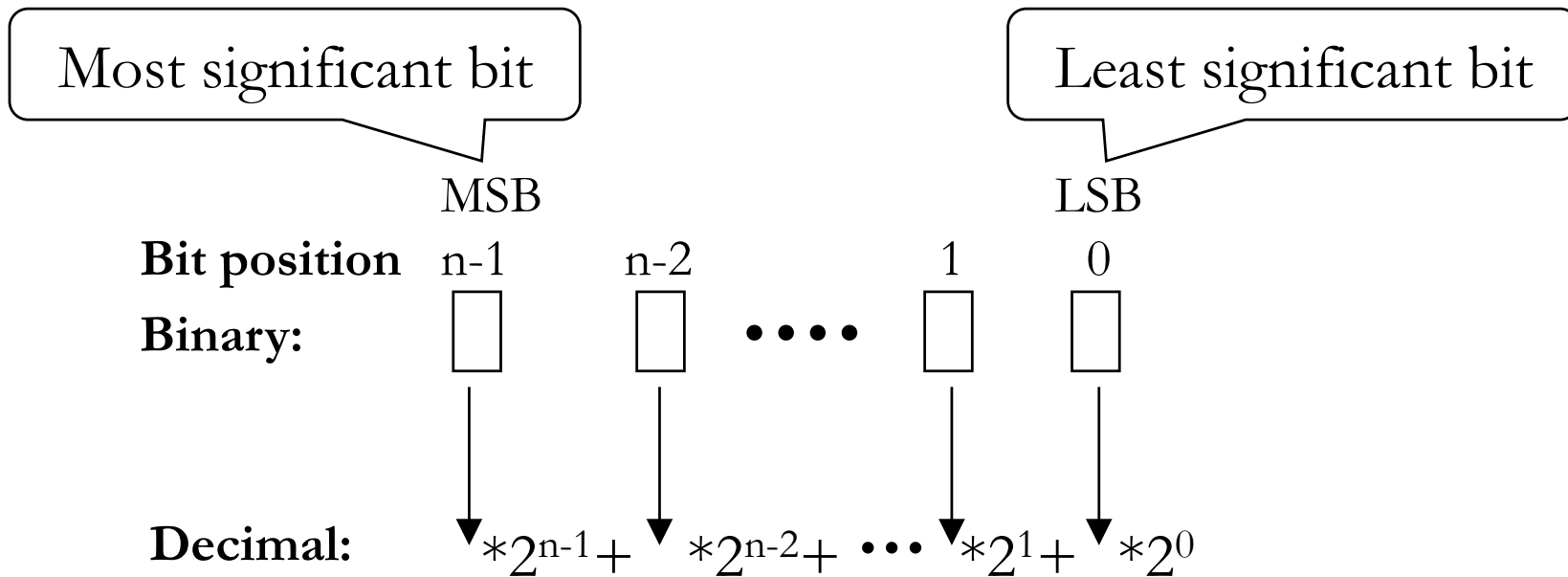
The bit

- Acronym for Binary digiT
- The smallest amount of meaningful info, according to information theory
 - If only 1 value is possible, there is no information
- Disadvantages: too little information per bit, must use many of them
- Advantages: easy to do computation, very reliable, simple circuits



Natural numbers representation

- Positive (unsigned) integers are very simple to represent in binary



Basic operations

- Addition, subtraction with binary numbers is easy:

$$\begin{array}{r} \text{1111} \\ 01101 \\ +01011 \\ \hline 11000 \end{array} \qquad \begin{array}{r} 0010 \\ 01101 \\ -01011 \\ \hline 00010 \end{array}$$

Diagram illustrating binary addition and subtraction with decimal equivalents:

- 13 (01101) + 11 (01011) = 24 (11000)
- 13 (01101) - 11 (01011) = 2 (00010)



Fixed bit-length arithmetic

- Hardware cannot handle infinite long bit sequences
- We end up with a few fixed sized data types
 - **Byte**: always 8 bits
 - **Word**: the ‘natural’ unit of access, usually 32 bits
- **Overflow** happens when a result does not fit
 - Numbers wrap-around when they become too large
 - Comp. arithmetic is modulo 2^n , n =number of bits



What about negative numbers?

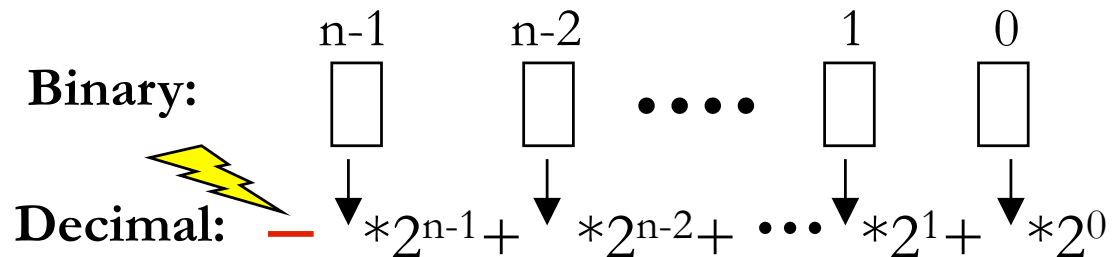
- No special symbols, e.g. +, −, available
- **Sign-magnitude** representation:
 - Use 1st bit (MSB) as the sign: 1-negative, 0-positive
- Complicates the arithmetic operations
 - The actual operation depends on the sign
- There is a better way



Two's complement representation

- What is the result of $000 - 001$?
...111

- The MSB has negative weighting:



- Arithmetic operations do not depend on the operands' signs

2's complement quirks

- The MSB is the sign
- Range is asymmetric: -2^{n-1} to $2^{n-1}-1$
- There are two kinds of overflows:
 - Positive overflow produces a negative number
 - Negative **underflow** produces a positive number
- To negate a number
 - Invert all bits ($0 \leftrightarrow 1$) and add 1, at the LSB
 - $-(-2^{n-1})$ overflows!
- $A-B = A + 2$'s complement of B



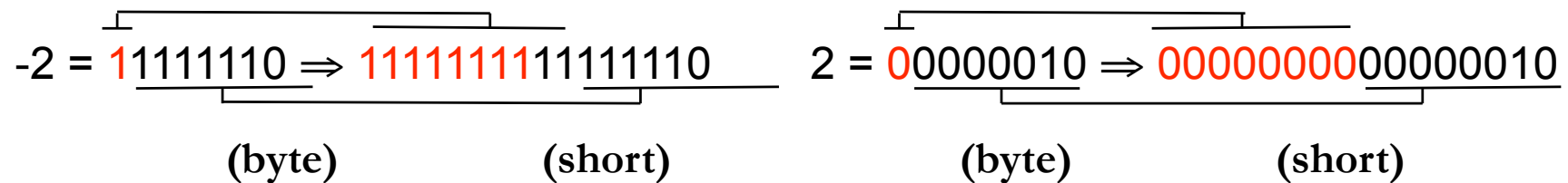
Converting between data types

- Converting from a smaller to a larger representation is done by **sign extension**

Example: from byte to short (16 bits):

2 = 00000010 ⇒ ?????????00000010

-2 = 11111110 ⇒ ?????????11111110



Shifting

- Shifting: move the bits of a data type left or right
 - Data bits falling off the edge are lost
- 0s fill up the empty bit places for left shifts
- For right shifts, two options:
 - Fill with 0: for non-numerical data (or positive integers)
 - Fill with the MSB: for 2's complement numbers
- Shift left by n is equivalent to multiplying by 2^n
- Shift right by n is equivalent to dividing by 2^n
- Example $6 = 00000110 \gg 2 \rightarrow 00000001 = 1$
 $-8 = 11111000 \gg 2 \rightarrow 11111110 = -2$



Hexadecimal notation

- Binary numbers (and other data) are too long and tedious for us to use
- Hexadecimal (base 16) is very commonly used in computer programming
- Hex digits: 0-9 and A-F
 - A=10, B=11, ..., F=15
- Conversion to/from binary is very easy:
Every 4 bits correspond to 1 hex digit:

$$\begin{array}{cccccccc} 1 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \\ \underbrace{\hspace{1.5em}} & \underbrace{\hspace{1.5em}} & & & & & & \\ \text{F(15)} & & 8 & & & & & \end{array} = \text{0xF8}$$

Hex is just a convenience, computers use the binary form



Real numbers - floating point

- Java's **float** (32 bits)
double (64 bits)

- IEEE 754:

– example 0.75 in base 10 \Rightarrow 0.11 in base 2

$(2^{-1} + 2^{-2} = 0.5 + 0.25 = 0.75)$

– Normalized:

mantissa **exponent**

0.11 \Rightarrow 1.1x2⁻¹

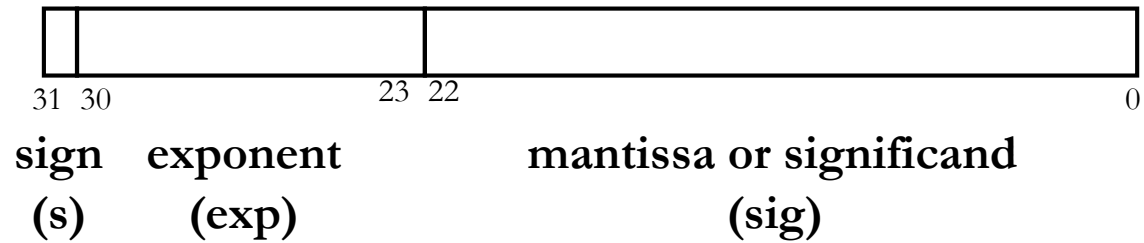
implicit
(always 1)

– example: 25 in base 10 \Rightarrow 11001 in base 2 \Rightarrow 1.1001x2⁴



Floating Point

- 32 bit:



$$(-1)^s \times (1.\text{sig}) \times 2^{\text{exp}-127}$$

e.g.,

$$(0.75)_{10} \rightarrow (0.11)_2 \rightarrow (1.1 \times 2^{-1})_2 \rightarrow 0\ 01111110\ 100000000000000000000000$$

- 64 bit:
 - exponent = 11 bits; significand = 52 bits
- Note: processors usually have specialized floating point units and extra fp registers to perform fp arithmetic



Representing characters, strings

- Characters need to be encoded in binary too
- Operations on characters have simpler requirements than on numbers, so the encoding choice is not crucial
- Most common representation is ASCII
 - Each character is held in a byte
 - E.g. '0' is 0x30, 'A' is 0x41, 'a' is 0x61
- Java uses Unicode which can encode characters from many (all?) languages
 - 16 bits per character required
- Words, sentences, etc. are just **strings** of characters
 - A special character, encoded as 0x00, shows where the string ends (in C)
 - Or the string length is kept with the string itself (in Java)



Summary

- Computers use binary representation
- 2's complement
- Floating point
- Characters and strings

