Chart Parsing: the CYK Algorithm

Informatics 2A: Lecture 18

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Grammar Restructuring

Deterministic parsing (e.g., LL(1)) aims to address a limited amount of local ambiguity – the problem of not being able to decide uniquely which grammar rule to use next in a left-to-right analysis of the input string.

By re-structuring the grammar, the parser can make a unique decision, based on a limited amount of look-ahead.

Recursive Descent parsing also demands grammar restructuring, in order to eliminate left-recursive rules that can get it into a hopeless loop.

Left Recursion

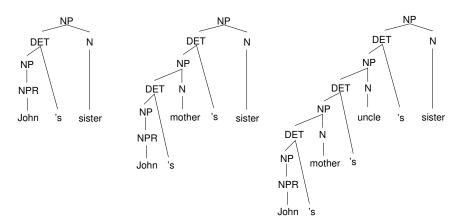
But grammars for natural human languages should be revealing, re-structuring the grammar may destroy this. (Indirectly) left-recursive rules are needed in English.

```
\begin{array}{l} \mathsf{NP} \to \mathsf{DET} \; \mathsf{N} \\ \mathsf{NP} \to \mathsf{NPR} \\ \mathsf{DET} \to \mathsf{NP} \; \mathsf{'s} \end{array}
```

These rules generate NPs with possessive modifiers such as:

John's sister
John's mother's sister
John's mother's uncle's sister
John's mother's uncle's sister's niece

Left Recursion



We don't want to re-structure our grammar rules just to be able to use a particular approach to parsing. Need an alternative.

Problems with Parsing as Search

- A recursive descent parser (top-down) will do badly if there are many different rules for the same LHS. Hopeless for rewriting parts of speech (preterminals) with words (terminals).
- A shift-reduce parser (bottom-up) does a lot of useless work: many phrase structures will be locally possible, but globally impossible. Also inefficient when there is much lexical ambiguity.
- On Both strategies do repeated work by re-analyzing the same substring many times.

We will see how chart parsing solves the re-parsing problem, and also copes well with ambiguity.

Dynamic Programming

With a CFG, a parser should be able to avoid re-analyzing sub-strings because the analysis of any sub-string is independent of the rest of the parse.

The parser's exploration of its search space can exploit this independence if the parser uses dynamic programming.

Dynamic programming is the basis for all chart parsing algorithms.

Parsing as Dynamic Programming

- Given a problem, systematically fill a table of solutions to sub-problems: this is called memoization.
- Once solutions to all sub-problems have been accumulated, solve the overall problem by composing them.
- For parsing, the sub-problems are analyses of sub-strings and correspond to constituents that have been found.
- Sub-trees are stored in a chart (aka well-formed substring table), which is a record of all the substructures that have ever been built during the parse.

Solves **re-parsing problem**: sub-trees are looked up, not re-parsed! Solves **ambiguity problem**: chart implicitly stores all parses!

Depicting a Chart

A chart can be depicted as a matrix:

- Rows and columns of the matrix correspond to the start and end positions of a span (ie, starting right before the first word, ending right after the final one);
- A cell in the matrix corresponds to the sub-string that starts at the row index and ends at the column index.
- It can contain information about the type of constituent (or constituents) that span(s) the substring, pointers to its sub-constituents, and/or predictions about what constituents might follow the substring.

CYK Algorithm

CYK (Cocke, Younger, Kasami) is an algorithm for recognizing and recording constituents in the chart.

- Assumes that the grammar is in Chomsky Normal Form: rules all have form $A \to BC$ or $A \to w$.
- Conversion to CNF can be done automatically.

```
Det Nom
      NP
                                                          NP
                                                                            Det Nom
    Nom
                    N | OptAP Nom
                                                       Nom \rightarrow
                                                                            book
                                                                                                         AP Nom
                                                                                         orange
 \mathsf{Opt}\mathsf{AP} \quad 	o \quad \epsilon \quad | \  \, \mathsf{Opt}\mathsf{Adv} \; \mathsf{A}
                                                          \mathsf{AP} \quad \rightarrow \quad
                                                                            heavy
                                                                                      orange
                                                                                                          Adv A

ightarrow heavy | orange
                                                        \mathsf{A} \quad \rightarrow
                                                                            heavy | orange
      Det \rightarrow
                                                         Det
                                                                   \rightarrow
\mathsf{OptAdv} \quad \rightarrow \quad
                    \epsilon | very
                                                         Adv
                                                                   \rightarrow
                                                                            very
                        book | orange
```

CYK: an example

Let's look at a simple example before we explain the general case.

(N.B. Converting to CNF sometimes breeds duplication!)

Now let's parse: a very heavy orange book

		1	2	3	4	5
		а	very	heavy	orange	book
0	а					
1	very					
2	heavy					
3	orange					
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very					
2	heavy					
3	orange					
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv			
2	heavy					
3	orange					
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv			
2	heavy			A,AP		
3	orange					
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv	AP		
2	heavy			A,AP		
3	orange					
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv	AP		
2	heavy			A,AP		
3	orange				Nom,A,AP	
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv	AP		
2	heavy			A,AP	Nom	
3	orange				Nom,A,AP	
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det				
1	very		Adv	AP	Nom	
2	heavy			A,AP	Nom	
3	orange				Nom,A,AP	
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	
1	very		Adv	AP	Nom	
2	heavy			A,AP	Nom	
3	orange				Nom,A,AP	
4	book					

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	
1	very		Adv	AP	Nom	
2	heavy			A,AP	Nom	
3	orange				Nom,A,AP	
4	book					Nom

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	
1	very		Adv	AP	Nom	
2	heavy			A,AP	Nom	
3	orange				Nom,A,AP	Nom
4	book					Nom

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	
1	very		Adv	AP	Nom	
2	heavy			A,AP	Nom	Nom
3	orange				Nom,A,AP	Nom
4	book					Nom

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	
1	very		Adv	AP	Nom	Nom
2	heavy			A,AP	Nom	Nom
3	orange				Nom,A,AP	Nom
4	book					Nom

		1	2	3	4	5
		а	very	heavy	orange	book
0	а	Det			NP	NP
1	very		Adv	AP	Nom	Nom
2	heavy			A,AP	Nom	Nom
3	orange				Nom,A,AP	Nom
4	book					Nom

CYK: The general algorithm

```
function CKY-Parse(words, grammar) returns table for j \leftarrow from 1 to Length(words) do table[j-1,j] \leftarrow \{A \mid A \rightarrow words[j] \in grammar\} for i \leftarrow from j-2 downto 0 do for \ k \leftarrow i+1 \ to \ j-1 \ do table[i,j] \leftarrow table[i,j] \cup \\ \{A \mid A \rightarrow BC \in grammar, \\ B \in table[i,k] \\ C \in table[k,j]\}
```

CYK: The general algorithm

function CKY-Parse(words, grammar) returns table for $i \leftarrow \text{from } 1 \text{ to } \text{Length}(words) \text{ do}$ loop over the columns $table[j-1,j] \leftarrow \{A \mid A \rightarrow words[j] \in grammar\} \mid fill bottom cell$ fill row *i* in column *j* for $i \leftarrow$ from j-2 downto 0 do for $k \leftarrow i + 1$ to i - 1 do loop over split locations $table[i, j] \leftarrow table[i, j] \cup | between i and j$ $\{A \mid A \rightarrow BC \in grammar, \}$ Check the grammar $B \in table[i, k]$ for rules that $C \in table[k, j]$ link the constituent in [i, k] with those in [k, j]. For each rule found store LHS in cell [i, j].

A succinct representation of CKY

We have a Boolean table called Chart, such that $\operatorname{Chart}[A,i,j]$ is true if there is a sub-phrase according the grammar that dominates words i through words j

Build this chart recursively, similarly to the Viterbi algorithm:

For
$$j > i + 1$$
:

$$Chart[A, i, j] = \bigvee_{k=i+1}^{j-1} \bigvee_{A \to B} \bigvee_{C} Chart[B, i, k] \land Chart[C, k, j]$$

Seed the chart, for i+1=j: $\operatorname{Chart}[A,i,i+1]=\operatorname{True}$ if there exists a rule $A\to w_{i+1}$ where w_{i+1} is the (i+1)th word in the string

From CYK Recognizer to CYK Parser

- So far, we just have a chart recognizer, a way of determining whether a string belongs to the given language.
- Changing this to a parser requires recording which existing constituents were combined to make each new constituent.
- This requires another field to record the one or more ways in which a constituent spanning (i,j) can be made from constituents spanning (i,k) and (k,j). (More clearly displayed in graph representation, see next lecture.)
- In any case, for a fixed grammar, the CYK algorithm runs in time $O(n^3)$ on an input string of n tokens.
- The algorithm identifies all possible parses.

CYK-style parse charts

Even without converting a grammar to CNF, we can draw *CYK-style* parse charts:

		1	2	3	4	5
		a	very	heavy	orange	book
0	а	Det			NP	NP
1	very		OptAdv	OptAP	Nom	Nom
2	heavy			A,OptAP	Nom	Nom
3	orange				N,Nom,A,AP	Nom
4	book					N,Nom

(We haven't attempted to show ϵ -phrases here. Could in principle use cells below the main diagonal for this . . .)

However, CYK-style parsing will have run-time worse than $O(n^3)$ if e.g. the grammar has rules $A \to BCD$.

Grammar Rules in CNF

 $\begin{array}{lll} S \rightarrow \textit{NP VP} & \textit{Nominal} \rightarrow \textit{book} | \textit{flight} | \textit{money} \\ S \rightarrow \textit{X1 VP} & \textit{Nominal} \rightarrow \textit{Nominal noun} \\ \textit{X1} \rightarrow \textit{Aux VP} & \textit{Nominal} \rightarrow \textit{Nominal PP} \\ S \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} \\ & \textit{VP} \rightarrow \textit{book} | \textit{include} | \textit{prefer} | \textit{CP} \rightarrow \textit{CP} \rightarrow \textit{CP} | \textit{CP} \rightarrow \textit{CP} | \textit{CP} \rightarrow \textit{C$

S o Verb NP S o X2 VP o X2 PP S o Verb PP X2 o Verb NP YP o Verb NP YP o Verb NP YP o Verb NP YP o VP o

 $NP o Det \ Nominal$ $PP o Preposition \ NP$

 $Verb o book|include|prefer \quad Noun o book|flight|money$

Let's parse Book the flight through Houston!

```
Grammar Rules in CNF
S \rightarrow NP VP
                                      Nominal \rightarrow book flight money
S \rightarrow X1 VP
                                      Nominal → Nominal noun
X1 \rightarrow Aux VP
                                      Nominal → Nominal PP
                                  VP \rightarrow book | include | prefer
S \rightarrow book | include | prefer
S \rightarrow Verb NP
                                      VPVerb \rightarrow NP
S \rightarrow X2
                                      VP \rightarrow X2 PP
S \rightarrow Verb PP
                                      X2 \rightarrow Verb NP
S \rightarrow VP PP
                                    VP \rightarrow Verb NP
NP \rightarrow TWA | Houston
                                 VP \rightarrow VP PP
NP → Det Nominal
                                 PP \rightarrow Preposition NP
Verb \rightarrow book | include | prefer Noun \rightarrow book | flight | money
```

Let's parse Book the flight through Houston!

- 1
_

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]				
	Det			
	[1, 2]			

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]				
	Det			
	[1, 2]			
		Nominal,		
		Noun		
		[2, 3]		

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]				
	Det			
	[1, 2]			
		Nominal,		
		Noun		
		[2, 3]		
			Prep	
			[3, 4]	
			`	

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]				
	Det			
	[1, 2]			
		Nominal,		
		Noun		
		[2, 3]		
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]	[0, 2]			
	Det			
	[1, 2]			
		Nominal,		
		Noun		
		[2, 3]		
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,				
Nominal,				
Noun				
[0, 1]	[0, 2]			
	Det	NP		
	[1, 2]	[1, 3]		
		Nominal,		
		Noun		
		[2, 3]		
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]		
	Det	NP		
	[1, 2]	[1, 3]		
		Nominal,		
		Noun		
		[2, 3]		
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]		
	Det	NP		
	[1, 2]	[1, 3]		
		Nominal,		
		Noun		
		[2, 3]	[2, 4]	
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]		
	Det	NP		
	[1, 2]	[1, 3]	[1,4]	
		Nominal,		
		Noun		
		[2, 3]	[2, 4]	
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]	[0, 4]	
	Det	NP		
	[1, 2]	[1, 3]	[1,4]	
		Nominal,		
		Noun		
		[2, 3]	[2, 4]	
			Prep	
			[3, 4]	
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]	[0, 4]	
	Det	NP		
	[1, 2]	[1, 3]	[1,4]	
		Nominal,		
		Noun		
		[2, 3]	[2, 4]	
			Prep	PP
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]	[0, 4]	
	Det	NP		
	[1, 2]	[1, 3]	[1,4]	
		Nominal,		Nominal
		Noun		
		[2, 3]	[2, 4]	[2, 5]
			Prep	PP
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		
Nominal,		VP,		
Noun		X2		
[0, 1]	[0, 2]	[0, 3]	[0, 4]	
	Det	NP		NP
	[1, 2]	[1, 3]	[1,4]	[1, 5]
		Nominal,		Nominal
		Noun		
		[2, 3]	[2, 4]	[2, 5]
			Prep	PP
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

Book	the	flight	through	Houston
S, VP, Verb,		S,		S ₁ , VP, X2,
Nominal,		VP,		S ₂ , VP,
Noun		X2		S ₃
[0, 1]	[0, 2]	[0, 3]	[0, 4]	[0, 5]
	Det	NP		NP
	[1, 2]	[1, 3]	[1,4]	[1, 5]
		Nominal,		Nominal
		Noun		
		[2, 3]	[2, 4]	[2, 5]
			Prep	PP
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

Visualizing the Chart

Book	the	flight	through	Houston
S, VP, Verb,		S,		S ₁ , VP, X2,
Nominal,		VP,		S ₂ , VP,
Noun		X2		S ₃
[0, 1]	[0, 2]	[0, 3]	[0, 4]	[0, 5]
	Det	NP		NP
	[1, 2]	[1, 3]	[1,4]	[1, 5]
		Nominal,		Nominal
		Noun		
		[2, 3]	[2, 4]	[2, 5]
		•	Prep←	–P₁P
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

Visualizing the Chart

Book	the	flight	through	Houston
S, VP, Verb,		S,		S ₁ , VP, X2,
Nominal,		VP,		S ₂ , VP,
Noun		X2		S ₃
[0, 1]	[0, 2]	[0, 3]	[0, 4]	[0, 5]
	Det	NP		NP
	[1, 2]	[1, 3]	[1,4]	[1, 5]
		Nominal, ←		Nominal
		Noun		
		[2, 3]	[2, 4]	[2, 5]
			Prep	P̈́P
			[3, 4]	[3, 5]
				NP, Proper-
				Noun
				[4, 5]

A Tribute to CKY (part 1/3)

You, my CKY algorithm, dictate every parser's rhythm, if Cocke, Younger and Kasami hadn't bothered, all of our parsing dreams would have been shattered.

You are so simple, yet so powerful, and with the proper semiring and time, you will be truthful, to return the best parse - anything less would be a crime.

With dynamic programming or memoization, you are one of a kind,
I really don't need to mention,
if it werent for you, all syntax trees would be behind.

A Tribute to CKY (part 2/3)

Failed attempts have been made to show there are better, for example, by using matrix multiplication, all of these impractical algorithms didn't matter you came out stronger, insisting on just using summation.

All parsing algorithms to you hail, at least those with backbones which are context-free, you will never become stale, as long as we need to have a syntax tree.

It doesn't matter that the C is always in front, or that the K and Y can swap, you are still on the same hunt, maximizing and summing, nonstop.

A Tribute to CKY (part 3/3)

Every Informatics student knows you intimately, they have seen your variants dozens of times, you have earned that respect legitimately, and you will follow them through their primes.

CKY, going backward and forward, inside and out, it is so straightforward -You are the best, there is no doubt.

Summary

- Parsing as search is inefficient (typically exponential time).
- Alternative: use dynamic programming and memoize sub-analysis in a chart to avoid duplicate work.
- The chart can be visualized as as a matrix.
- The CYK algorithm builds a chart in $O(n^3)$ time. The basic version gives just a recognizer, but it can be made into a parser if more info is recorded in the chart.

Reading: J&M (2nd ed), Chapter. 13, Sections 13.3–13.4

NLTK Book, Chapter. 8 (Analyzing Sentence

Structure), Section 8.4

Next lecture: the Earley parser or dynamic programming for top-

down parsing