

Informatics 1

Functional Programming Lecture 19

Monday 1 December 2008

Complexity

Philip Wadler

University of Edinburgh

Part I

Append

How long does it take to append?

```
(++) :: [a] -> [a] -> [a]
[]    ++ ys  =  ys
(x:xs) ++ ys =  x:(xs ++ ys)
```

```
"abcd" ++ "ef"
= 'a':("bcd" ++ "ef")
= 'a':('b':("cd" ++ "ef"))
= 'a':('b':('c':("d" ++ "ef")))
= 'a':('b':('c':('d':("" ++ "ef"))))
= 'a':('b':('c':('d':"ef")))
= "abcdef"
```

Computing `xs++ys` takes time proportional to the length of `xs`.

Linear vs. quadratic append

Associate to the right

"a"++("b"++("c"++("d"++("e"++[])))) = "abcde"

$$\underbrace{1 + \dots + 1}_{n \text{ times}} = n$$

Associate to the left

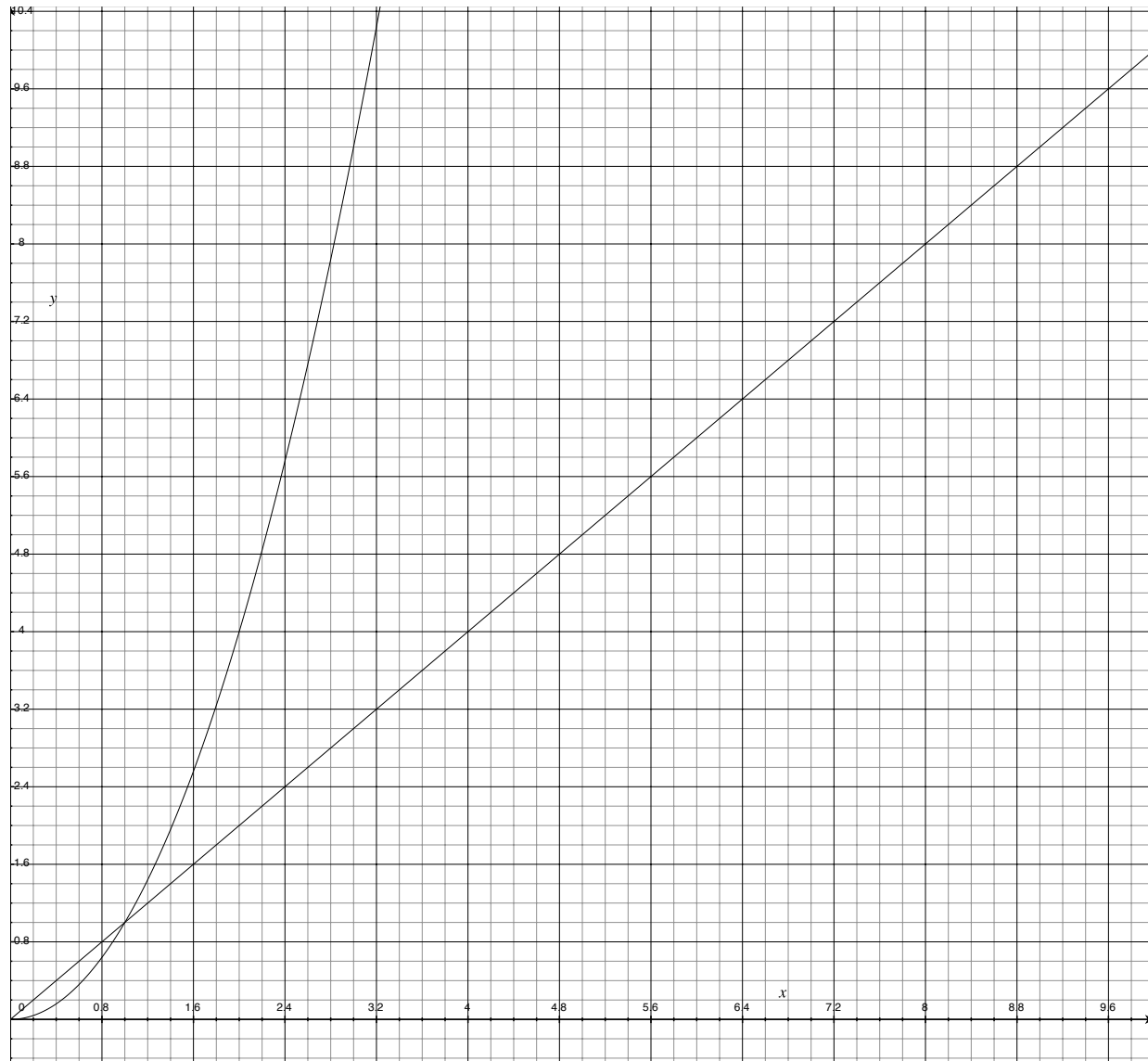
((((([]++"a")++"b")++"c")++"d")++"e") = "abcde"

$$\underbrace{0 + 1 + \dots + (n - 1)}_{n \text{ times}} = \frac{n(n - 1)}{2}$$

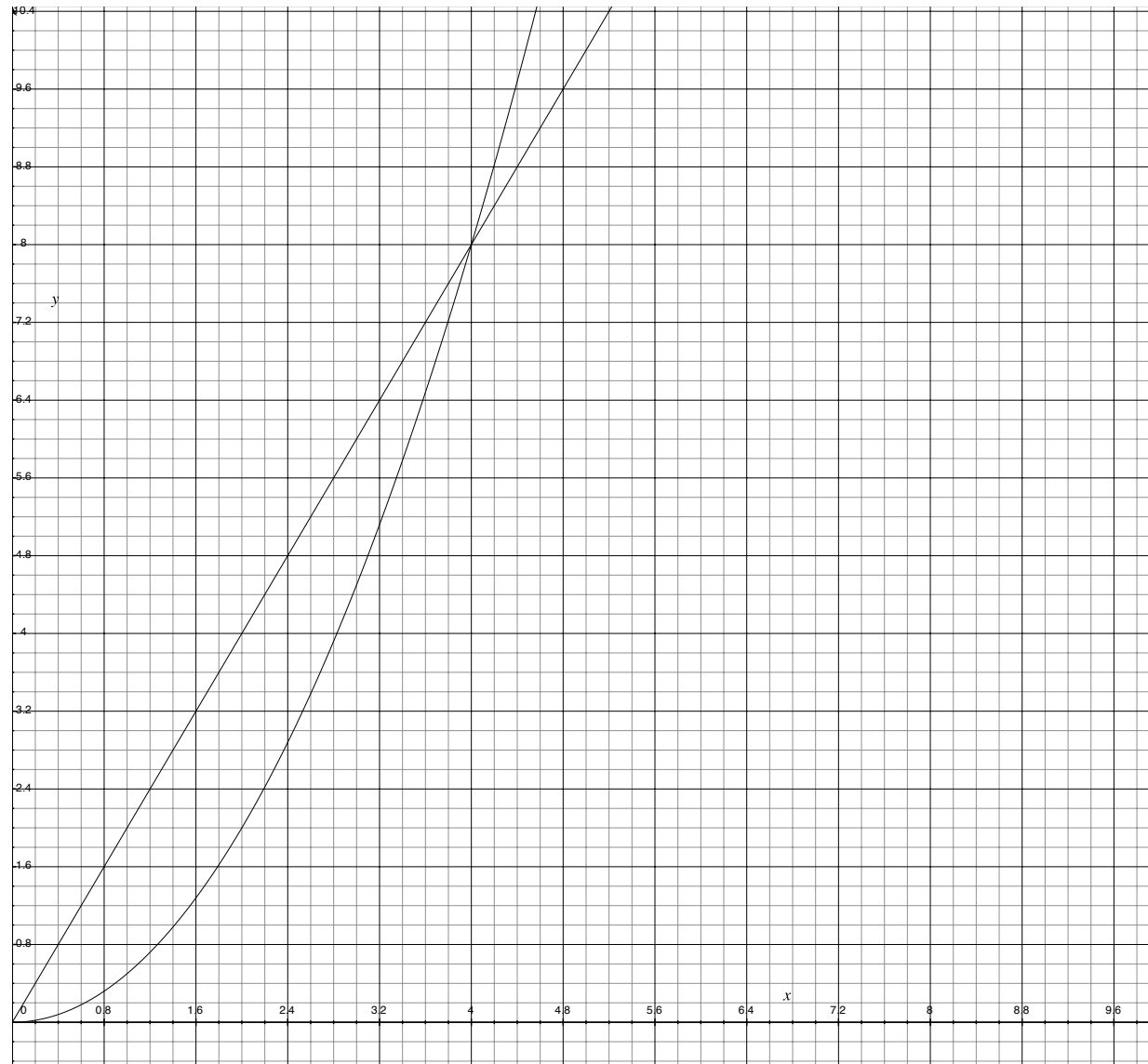
Part II

Complexity

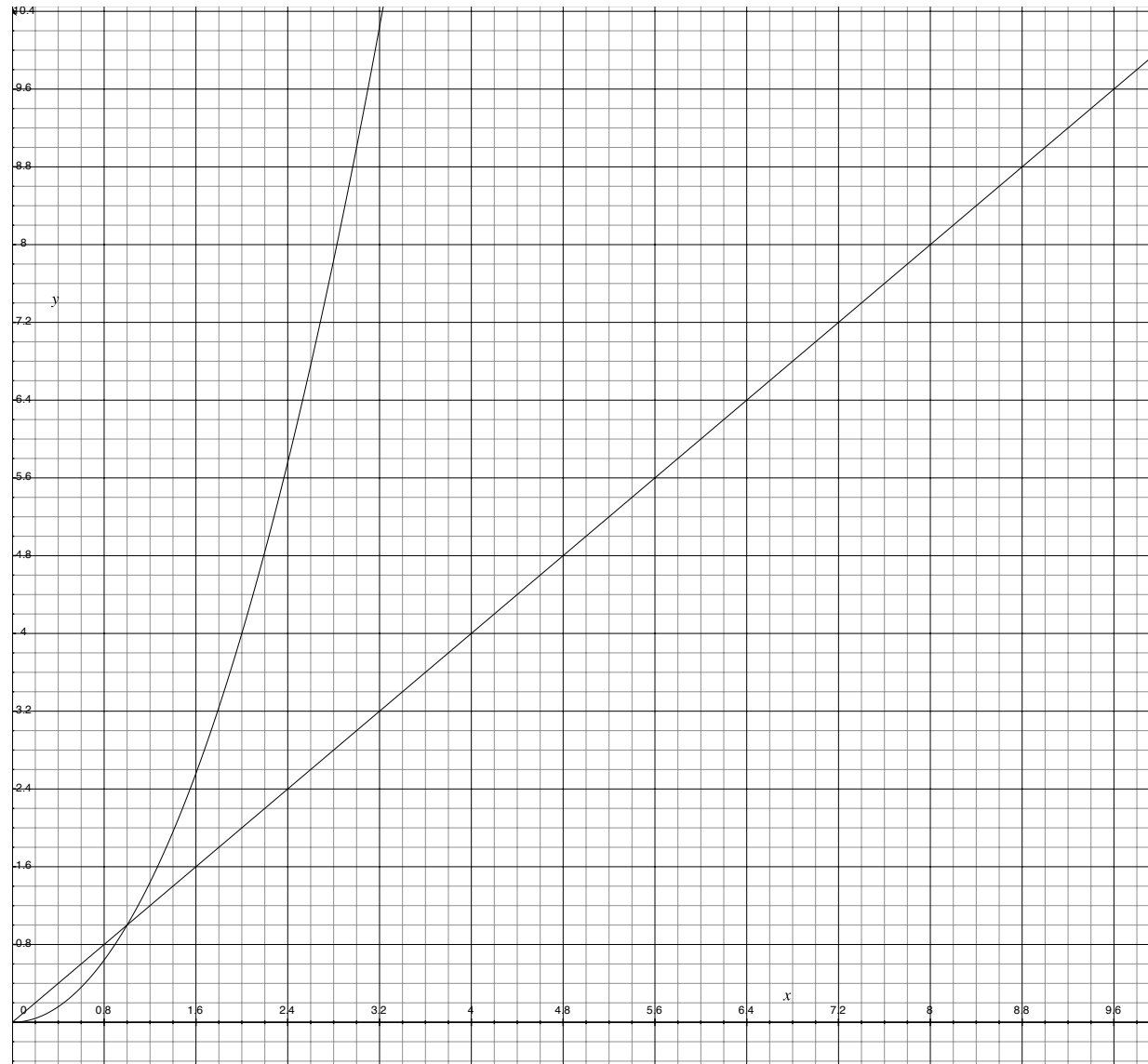
$t = n$ vs $t = n^2$



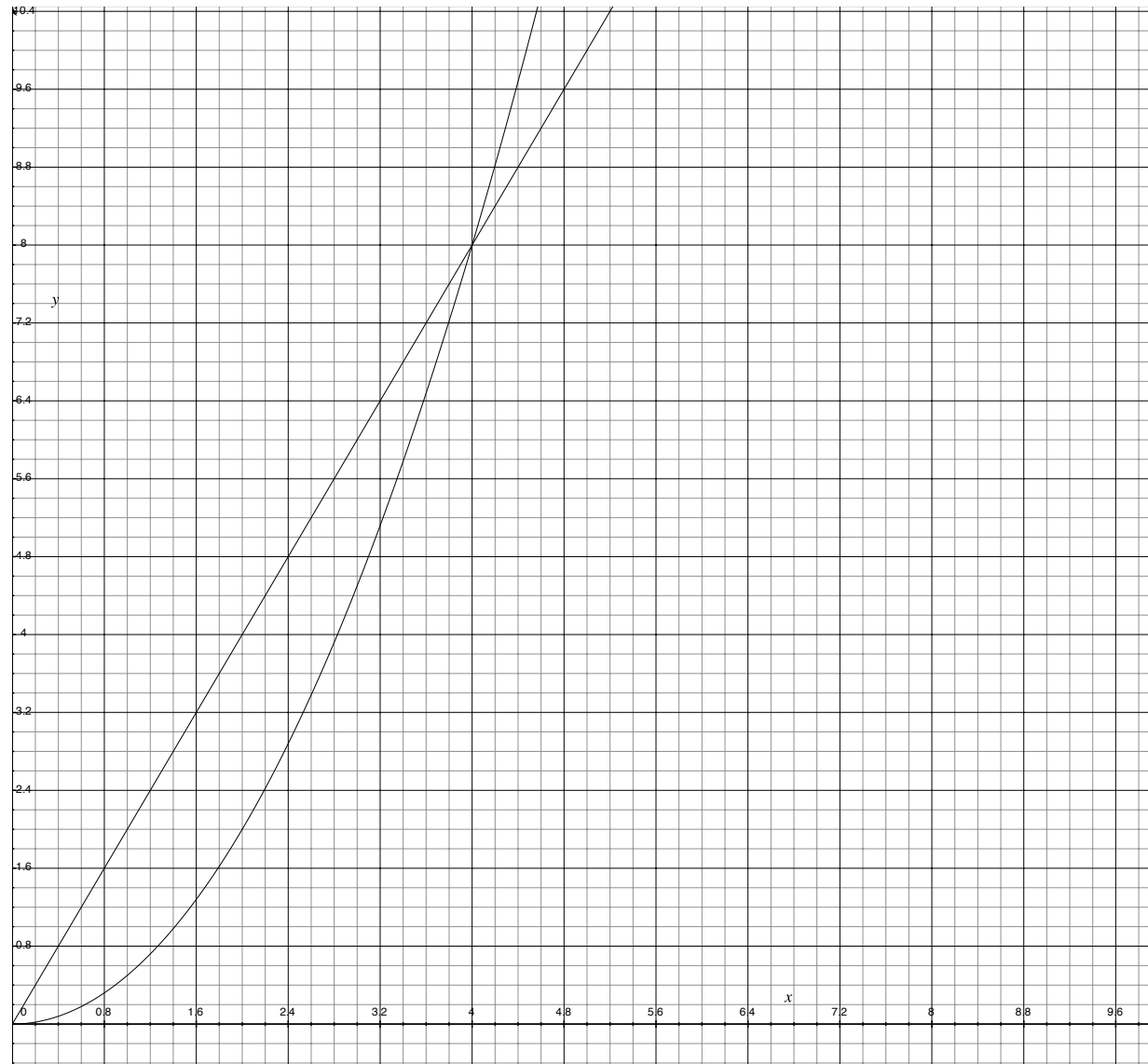
$$t = 2n \text{ vs } t = 0.5n^2$$



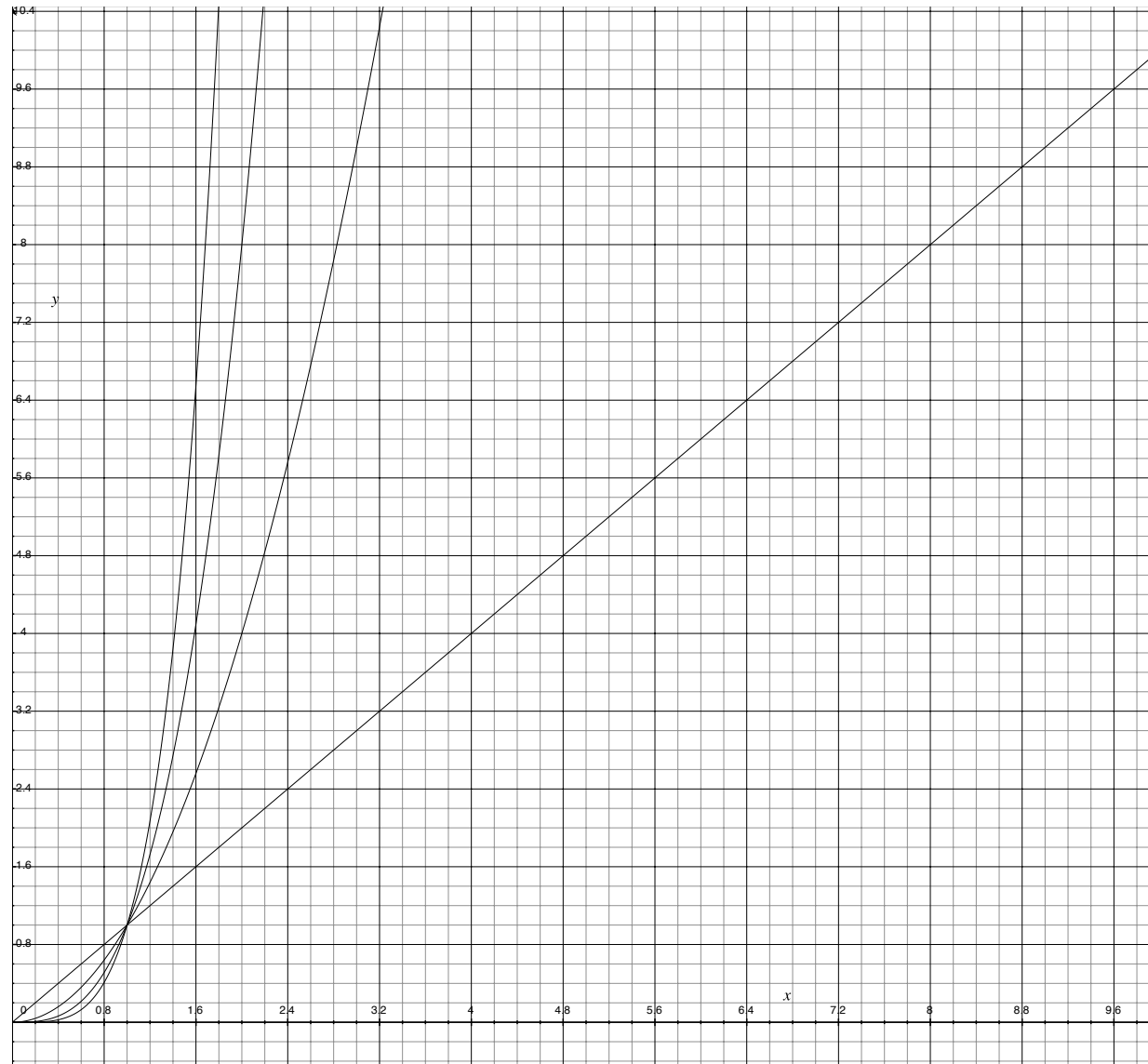
$O(n)$ vs $O(n^2)$



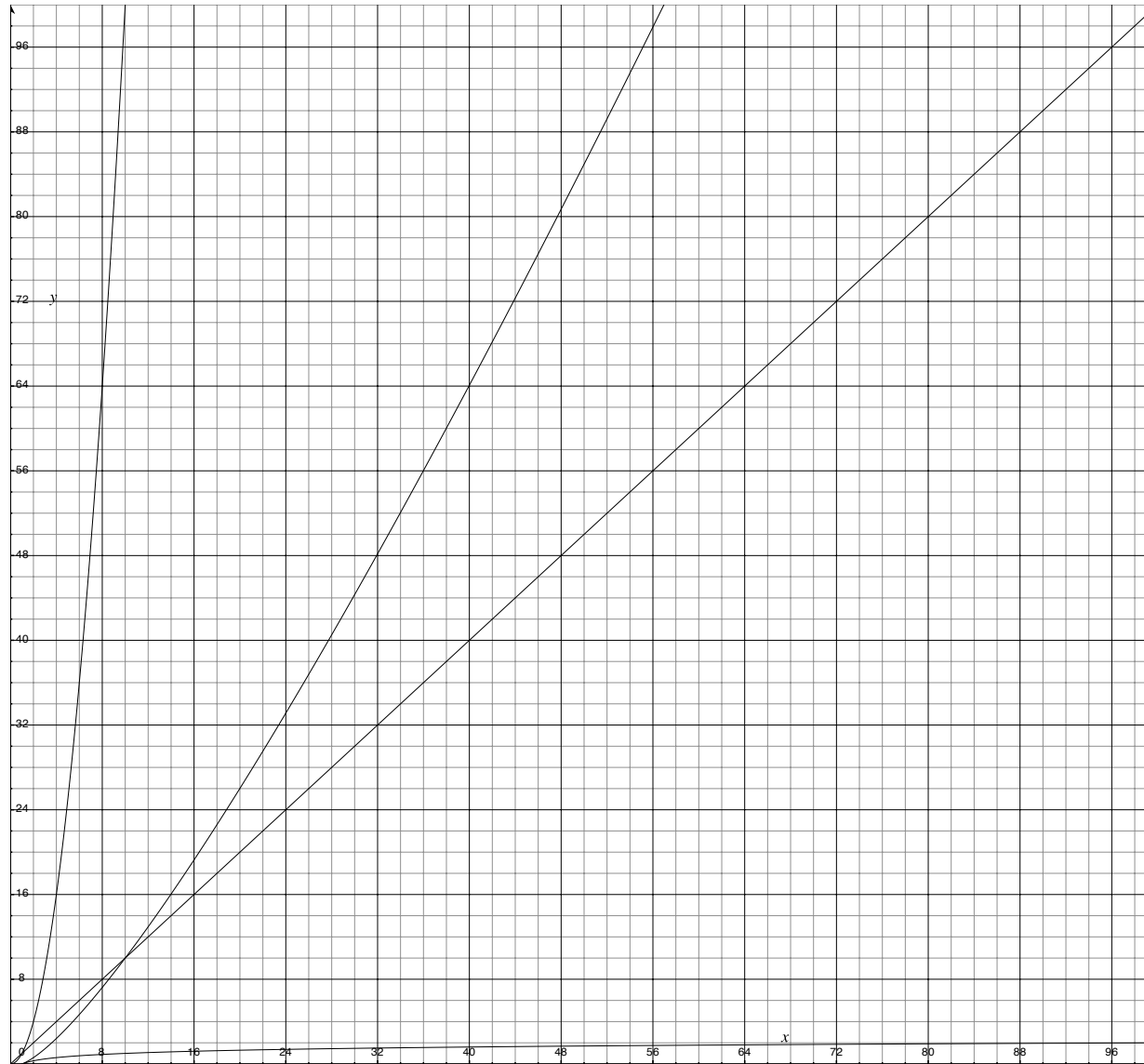
$O(n)$ vs $O(n^2)$



$O(n)$, $O(n^2)$, $O(n^3)$, $O(n^4)$



$O(\log n)$, $O(n)$, $O(n \log n)$, $O(n^2)$



Part III

Fold right and fold left

Fold right and fold left

```
foldr (+) 0 [1,2,3,4] = 1+(2+(3+(4+0)))
foldl (+) 0 [1,2,3,4] = (((0+1)+2)+3)+4
```

```
foldr :: (x -> a -> a) -> a -> [x] -> a
foldr f a [] = a
foldr f a (x:xs) = f x (foldr f a xs)
```

```
foldl :: (a -> x -> a) -> a -> [x] -> a
foldl f a [] = a
foldl f a (x:xs) = (foldl f (f a x) xs)
```

Fold right, sum

```
foldr (+) 0 [1..4]
=
1 + foldr (+) 0 [2..4]
=
1 + (2 + foldr (+) 0 [3..4])
=
1 + (2 + (3 + foldr (+) 0 [4..4]))
=
1 + (2 + (3 + (4 + foldr (+) 0 [5..4])))
=
1 + (2 + (3 + (4 + foldr (+) 0 [])))
=
1 + (2 + (3 + (4 + 0)))
=
10
```

Linear time, linear space

Fold right, sum

```
foldl (+) 0 [1..4]
=
foldr (+) (0+1) [2..4]
=
foldr (+) ((0+1)+2) [3..4]
=
foldr (+) (((0+1)+2)+3) [4..4]
=
foldr (+) ((((0+1)+2)+3)+4) [5..4]
=
foldr (+) (((((0+1)+2)+3)+4) [])
=
(((0+1)+2)+3)+4
=
10
```

But this does not reflect the space behaviour!

Fold left, sum

```
foldl (+) 0 [1..4]
=
foldl (+) (0+1) [2..4]
=
foldl (+) 1 [2..4]
=
foldl (+) (1+2) [3..4]
=
foldl (+) 3 [3..4]
=
foldl (+) (3+3) [4..4]
=
foldl (+) 6 [4..4]
=
foldl (+) (6+4) [5..4]
=
foldl (+) 10 []
=
10
```

Linear time, constant space

Fold right, append

```
foldr (++) [] ["a", "b", "c", "d"]
=
"a" ++ foldr (++) [] ["b", "c", "d"]
=
"a" ++ ("b" ++ foldr (++) [] ["c", "d"])
=
"a" ++ ("b" ++ ("c" ++ foldr (++) [] ["d"]))
=
"a" ++ ("b" ++ ("c" ++ ("d" ++ foldr (++) [] [])))
=
"a" ++ ("b" ++ ("c" ++ ("d" ++ [])))
=
"abcd"
```

Linear time, linear space

Fold left, append

```
foldl (++) [] ["a", "b", "c", "d"]
=
foldl (++) ([] ++ "a") ["b", "c", "d"]
=
foldl (++) "a" ["b", "c", "d"]
=
foldl (++) ("a" ++ "b") ["c", "d"]
=
foldl (++) "ab" ["c", "d"]
=
foldl (++) ("ab" ++ "c") ["d"]
=
foldl (++) "abc" ["d"]
=
foldl (++) ("abc" ++ "d") []
=
foldl (+) "abcd" []
=
"abcd"
```

Linear time, constant space

Part IV

Sort

Insert

```
insert :: Int -> [Int] -> [Int]
insert x [] = [x]
insert x (y:ys) | x > y = y : insert x ys
                | otherwise = x : y : ys
```

```
insert 3 [1,2,4]
=
insert 3 (1:2:4:[])
=
1 : insert 3 (2:4:[])
=
1 : 2 : insert 3 (4:[])
=
1 : 2 : 3 : 4 : []
=
[1,2,3,4]
```

Insertion sort

```
isort :: [Int] -> [Int]
isort []      = []
isort (x:xs) = insert x (isort xs)
```

```
    isort [3,1,4,2]
=
    insert 3 (isort [1,4,2])
=
    insert 3 (insert 1 (isort [4,2]))
=
    insert 3 (insert 1 (insert 4 (isort [2])))
=
    insert 3 (insert 1 (insert 4 (insert 2 [])))
=
    insert 3 (insert 1 (insert 4 [2]))
=
    insert 3 (insert 1 [2,4])
=
    insert 3 [1,2,4]
=
    [1,2,3,4]
```

Insertion sort

```
insert :: Ord a => a -> [a] -> [a]
insert x [] = [x]
insert x (y:ys) | x > y = y : insert x ys
                 | otherwise = x : y : ys
```

constant time $O(1)$, best case

linear time $O(n)$, average case

```
isort :: Ord a => [a] -> [a]
isort [] = []
isort (x:xs) = insert x (isort xs)
```

linear time $O(n)$, best case

quadratic time $O(n^2)$, average case

Insertion sort, higher order

```
insert :: Ord a => a -> [a] -> [a]
insert x ys =
    takeWhile (x>) ys ++ [x] ++ dropWhile (x>) ys
```

```
isort :: Ord a => [a] -> [a]
isort xs = foldr insert [] xs
```

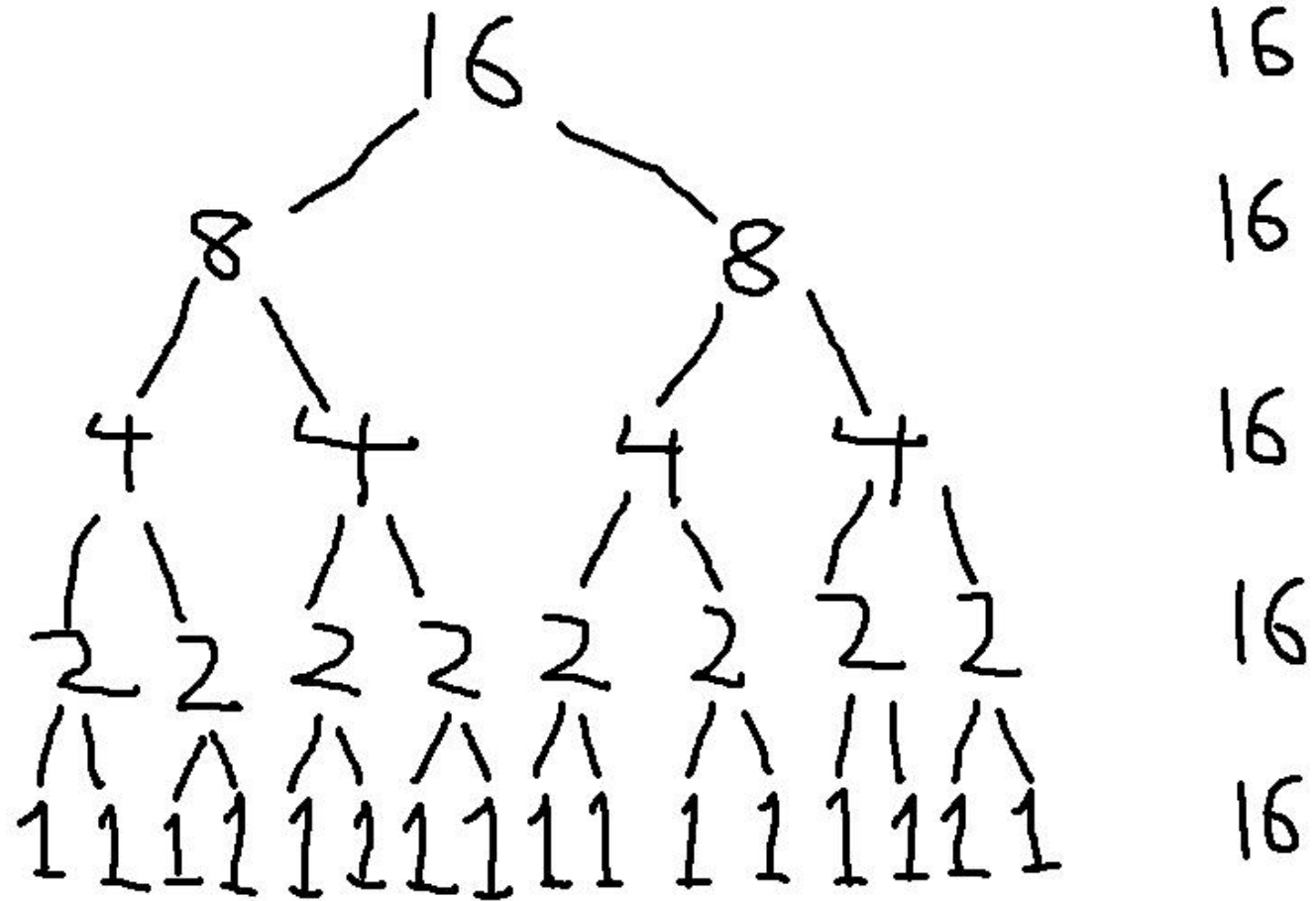
Quicksort

```
qsort :: Ord a => [a] -> [a]
qsort []      = []
qsort (x:ys)  =
  qsort [ y | y <- ys, y < x ] ++
  [x] ++
  qsort [ y | y <- ys, y >= x ]
```

$O(n \log n)$, average case

$O(n^2)$, worst case

Where the logs come from



Merge sort

```
merge :: Ord a => [a] -> [a] -> [a]
merge [] ys          = ys
merge (x:xs) []      = x:xs
merge (x:xs) (y:ys)
  | x < y            = x : merge xs (y:ys)
  | otherwise        = y : merge (x:xs) ys
```

```
split :: [a] -> ([a],[a])
split []          = ([], [])
split [x]         = ([x], [])
split (x:y:zs)   = (x:xs, y:ys)
  where
    (xs,ys) = split zs
```

```
msort :: Ord a => [a] -> [a]
msort []         = []
msort [x]        = [x]
msort zs         = merge (msort xs) (msort ys)
  where
    (xs,ys) = split zs
```

Merge sort

$O(n \log n)$, always