

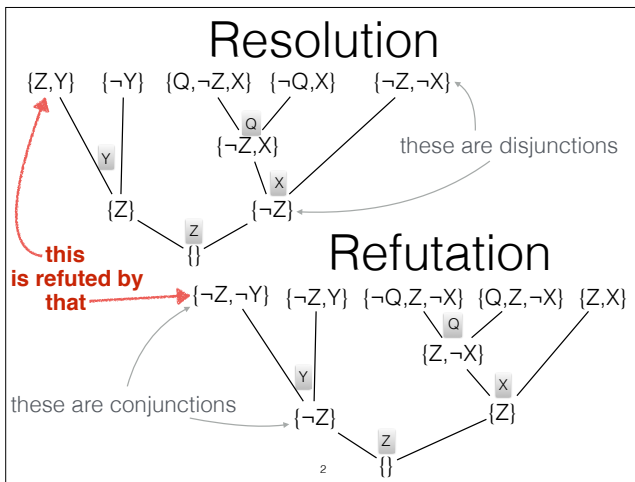
Informatics 1

Lecture 9 Searching for Satisfaction

Michael Fourman

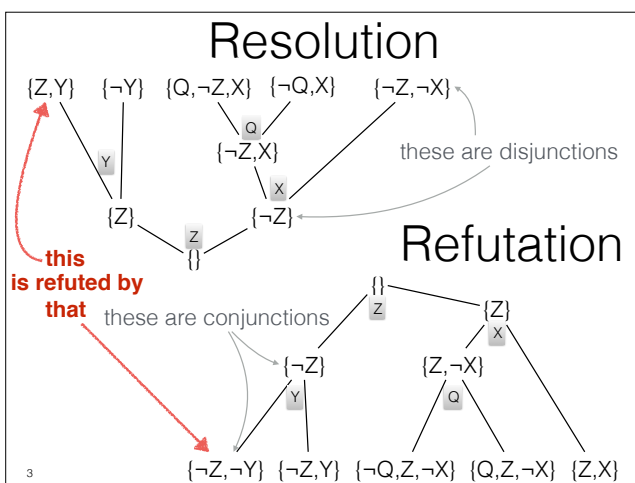


In this lecture we consider formal descriptions of the relationships between a finite number of individuals. We may have different types of individual



From the resolution proof we can derive a refutation.

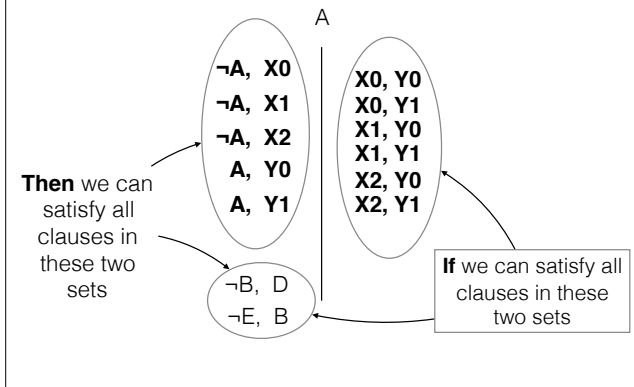
The lower tree demonstrates the fact that whatever values we choose for the variables, we will arrive at a clause that is false for our chosen values. This suffices to show that, no matter what choice of values we make, the conjunction is



We normally grow refutation trees downwards.

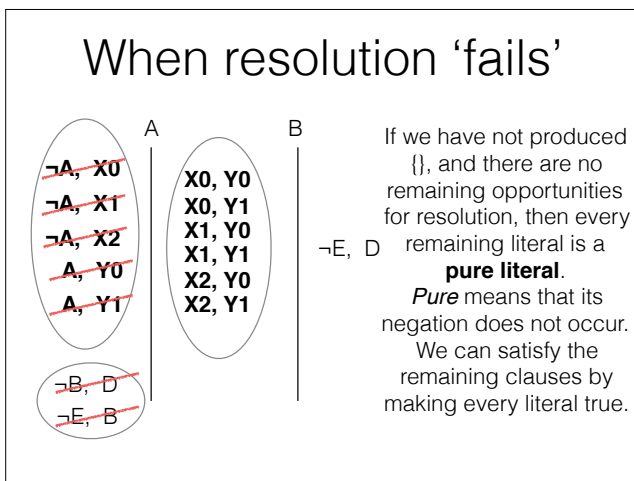
A refutation tree demonstrates the fact that whatever values we choose for the variables, we will arrive at a clause that is false for our chosen values. This suffices to show that, no matter what choice of values we make, the

When resolution 'fails'



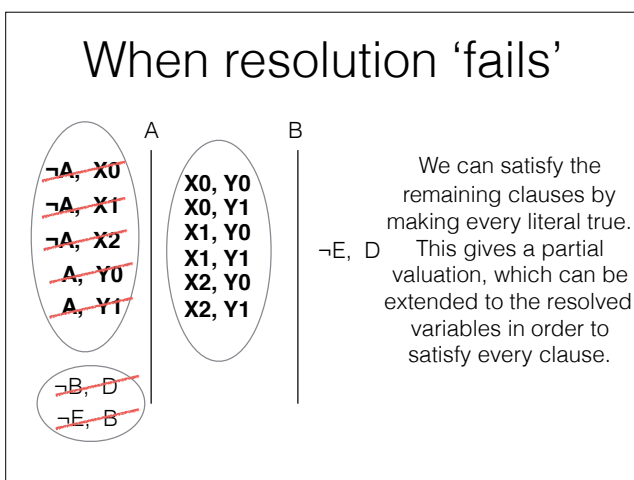
If we can satisfy all the Xs, then making A true will do the trick. If we cannot satisfy X_i then we must be able to satisfy all the Ys, and so making A false will do the trick.

When resolution 'fails'



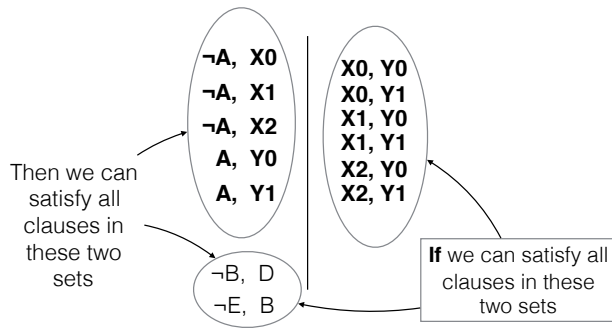
So, once we have resolved all the X, $\neg X$ pairs, we can focus on clauses not mentioning A. Eventually we will either produce {}, or have a set of clauses with no complementary pairs.

When resolution 'fails'



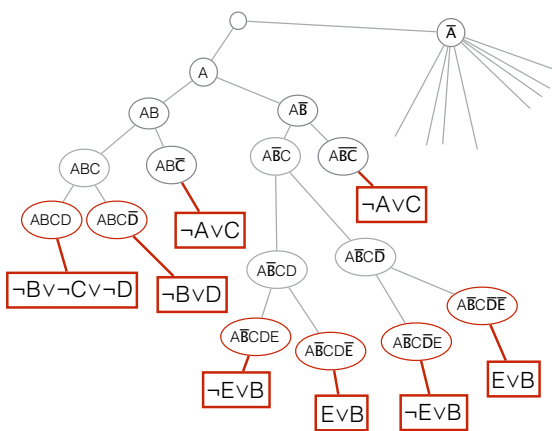
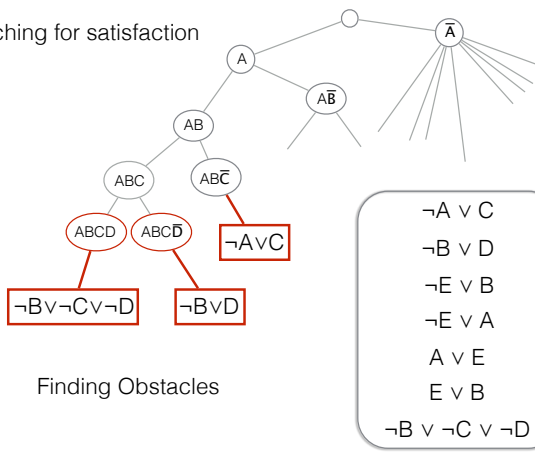
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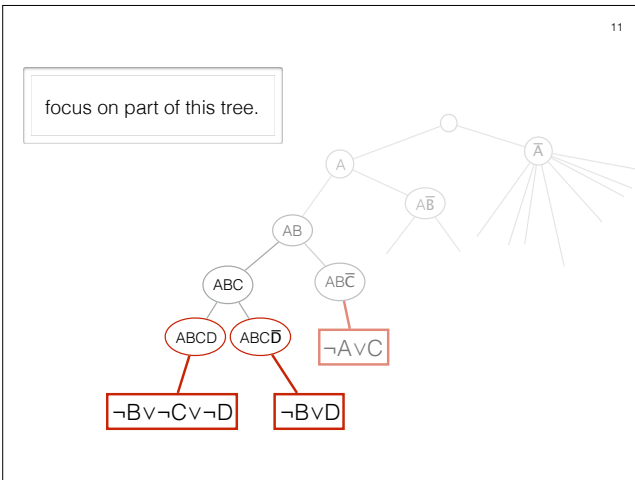
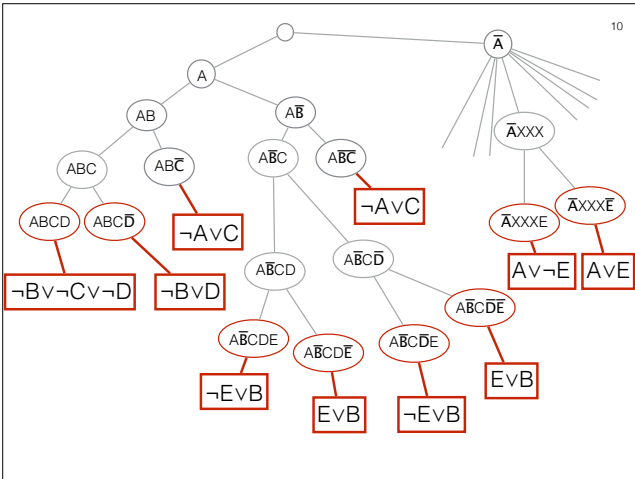
When resolution 'fails'



If we can satisfy all the Xs, then making A true will do the trick.
 If we cannot satisfy X_i then we must be able to satisfy all the Ys, and so making A false will do the trick.

Searching for satisfaction





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Premises

$\neg B \vee \neg C \vee \neg D$ $\neg B \vee D$

Conclusion

$\neg B \vee \neg C$

A valid inference Any assignment of truth values that makes all the premises true will make the conclusion true.

The conclusion follows from the premises

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Premises

$$\frac{\neg B \vee \neg C \vee \neg D \quad \neg B \vee D}{\neg B \vee \neg C}$$

Conclusion

For any valid inference Any assignment of truth values that makes the conclusion false will make at least one of the premises false.

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Premises

$$\frac{\neg B \vee \neg C \vee \neg D \quad \neg B \vee D}{\neg B \vee \neg C}$$

Conclusion

A special property of this inference If some assignment XYZ of values for ABC makes the conclusion false then the assignments XYZD and XYZD̄ each make one or other of the two premises false.

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Resolution

$$\frac{\neg B \vee \neg C \vee \neg D \quad \neg B \vee D}{\neg B \vee \neg C}$$

$$\frac{\neg B \vee \neg C \quad \neg A \vee C}{\neg B \vee \neg A}$$

Resolution

$$\frac{UvVvWvXv\bar{C} \quad XvYvZvC}{UvVvWvXvYvZ}$$

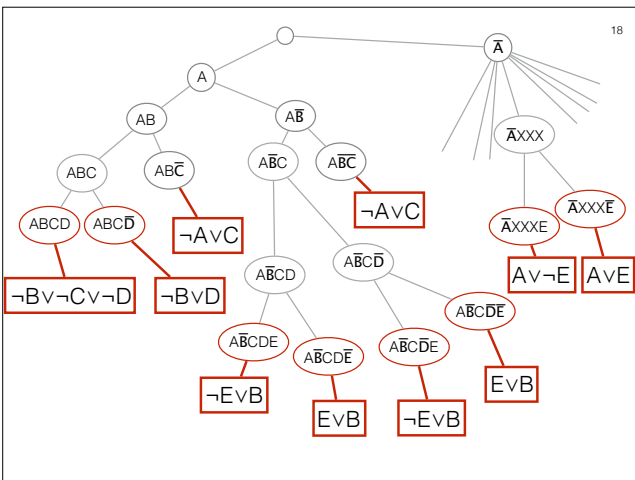
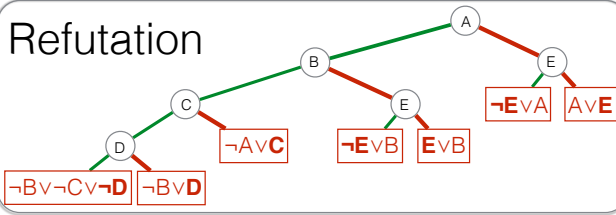
Resolution

$$\frac{\bar{B}v\bar{C}v\bar{D} \quad \bar{B}vD}{\bar{B}v\bar{C} \quad \bar{A}vC} \quad \frac{\bar{E}vB \quad EvB}{B} \quad \frac{\bar{E}vA \quad AvE}{A}$$

$$\frac{\bar{B}v\bar{C} \quad \bar{A}vC}{\bar{B}v\bar{A}} \quad \frac{\bar{E}vB \quad EvB}{B} \quad \frac{\bar{E}vA \quad AvE}{A}$$

$$\frac{\bar{B}v\bar{A} \quad B \quad A}{\perp}$$

Refutation



Ideal! Use the problem to simplify the search

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