# Empirical Methods in Natural Language Processing Lecture 9 Parsing (I): Context-free grammars and chart parsing

Philipp Koehn

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#### The path so far

- Originally, we treated language as a sequence of words  $\rightarrow$  n-gram language models
- Then, we introduced the notion of syntactic properties of words
  → part-of-speech tags
- Now, we look at syntactic relations between words  $\rightarrow$  syntax trees



#### A simple sentence

I like the interesting lecture



#### **Part-of-speech tags**

- I	like	the	interesting	lecture
PRO	VB	DET	JJ	NN



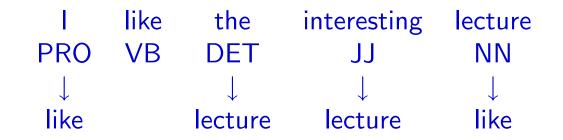
### Syntactic relations

I like the interesting lecture PRO VB DET JJ NN

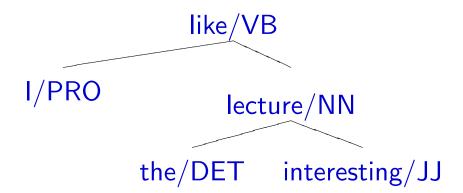
- The adjective *interesting* gives more information about the noun *lecture*
- The determiner *the* says something about the noun *lecture*
- The noun *lecture* is the object of the verb *like*, specifying *what* is being liked
- The pronoun *I* is the subject of the verb *like*, specifying *who* is doing the liking



#### **Dependency structure**



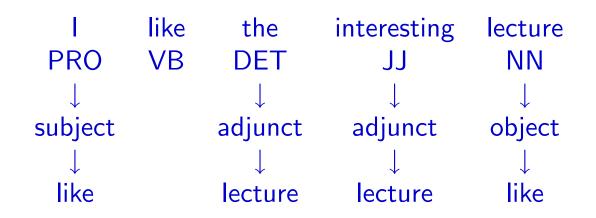
This can also be visualized as a **dependency tree**:





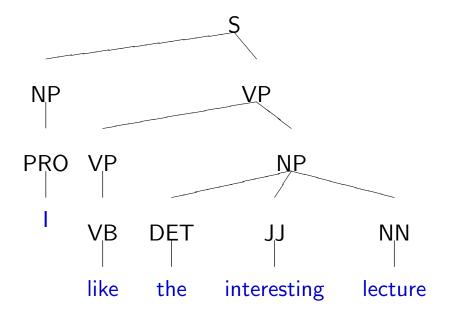
#### **Dependency structure (2)**

The dependencies may also be **labeled** with the type of dependency



#### **Phrase-structure tree**

A popular grammar formalism is **phrase structure grammar** Internal nodes combine leaf nodes into phrases, such as *noun phrases (NP)* 



nformation

ICS



#### Building phrase-structure trees

- Our task for this week: **parsing** 
  - given: an input sentence with part-of-speech tags
  - wanted: the right syntax tree for it
- Formalism: context-free grammars
  - non-terminal nodes such as NP, S appear inside the tree
  - terminal nodes such as *like*, *lecture* appear at the leafs of the tree
  - rules such as  $NP \rightarrow DET JJ NN$

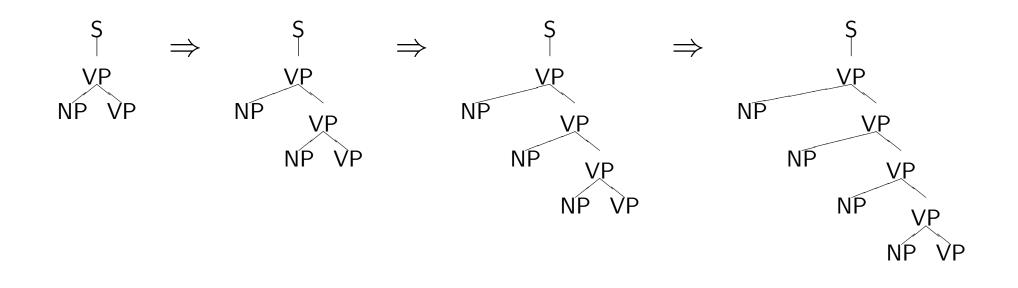


### Applying the rules

Input	Rule	Output	
S	$S \rightarrow NP VP$	NP VP	
NP VP	$NP \to PRO$	PRO VP	
PRO VP	$PRO \to I$	/ VP	
/ VP	$VP \rightarrow VP NP$	/ VP NP	
/ VP NP	$VP \rightarrow VB$	/ VB	
/ VB NP	VB → <i>like</i>	<i>I like</i> NP	
<i>I like</i> NP	$NP \to DET JJ NN$	<i>I like</i> DET JJ NN	
<i>I like</i> DET JJ NN	$DET \to the$	<i>I like the</i> JJ NN	
I like the JJ NN	$JJ \rightarrow interesting$	I like the interesting NN	
I like the interesting NN	$NN \rightarrow \mathit{lecture}$	I like the interesting lecture	

#### Recursion

Rules can be applied **recursively**, for example the rule  $VP \rightarrow NP VP$ 



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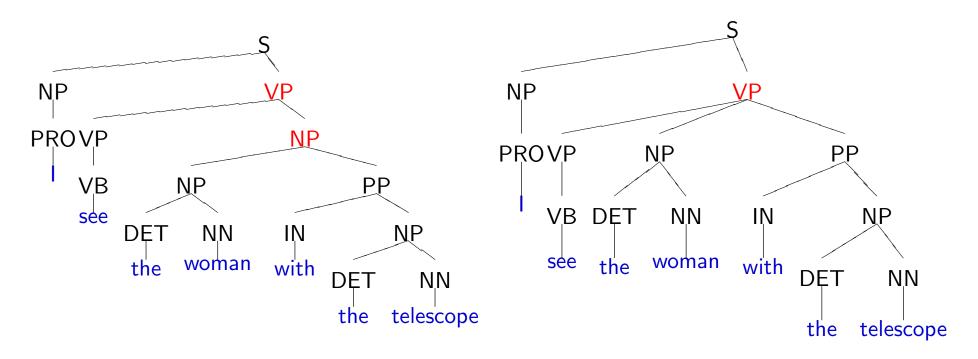


#### **Context-free grammars in context**

- Chomsky hierarchy of formal languages (terminals in caps, non-terminal lowercase)
  - **regular**: only rules of the form  $A \rightarrow a, A \rightarrow B, A \rightarrow Ba$  (or  $A \rightarrow aB$ ) Cannot generate languages such as  $a^n b^n$
  - **context-free**: left-hand side of rule has to be single non-terminal, anything goes on right hand-side. Cannot generate  $a^n b^n c^n$
  - context-sensitive: rules can be restricted to a particular context, e.g.  $\alpha A\beta \rightarrow \alpha aBc\beta$ , where  $\alpha$  and  $\beta$  are strings of terminal and non-terminals
- Moving up the hierarchy, languages are more expressive and parsing becomes computationally more expensive
- Is natural language context-free?

### Why is parsing hard?

**Prepositional phrase attachment:** Who has the *telescope*?

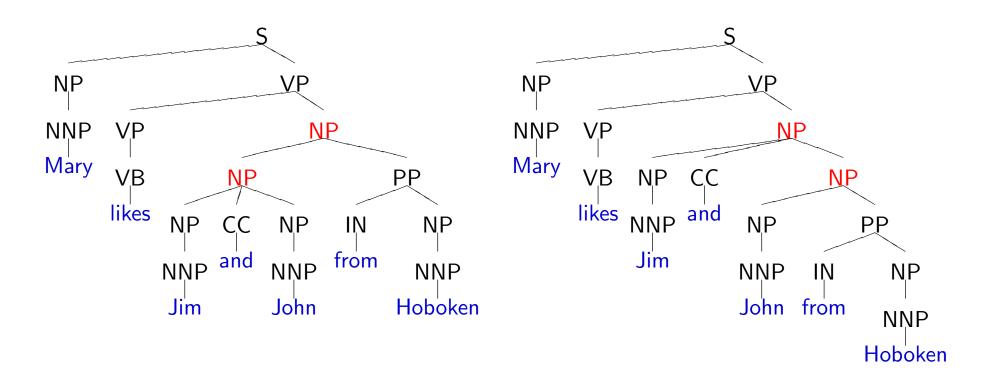


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### Why is parsing hard?

**Scope:** Is *Jim* also from *Hoboken*?



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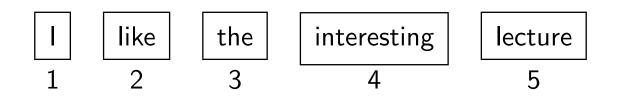
## **CYK** Parsing

- We have input sentence: *I like the interesting lecture*
- We have a set of context-free rules:  $S \rightarrow NP VP, NP \rightarrow PRO, PRO \rightarrow I, VP \rightarrow VP NP, VP \rightarrow VB, VB \rightarrow like,$  $NP \rightarrow DET JJ NN, DET \rightarrow the, JJ \rightarrow, NN \rightarrow lecture$
- Cocke-Younger-Kasami (CYK) parsing
  - a **bottom-up** parsing algorithm
  - uses a **chart** to store intermediate result



#### Example

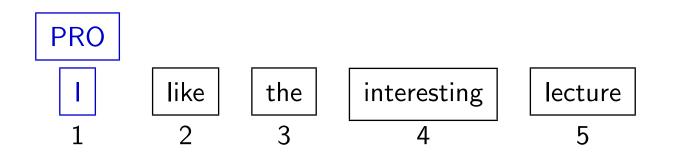
Initialize chart with the words





# Example (2)

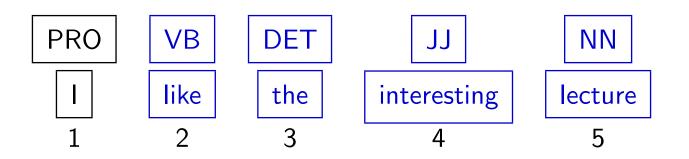
Apply first terminal rule  $PRO \rightarrow I$ 





## Example (3)

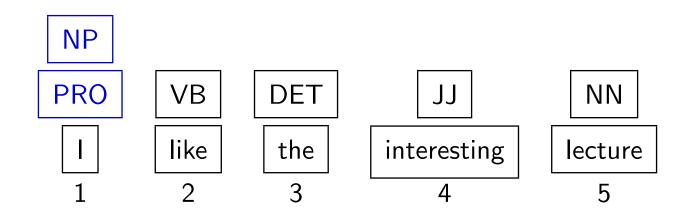
... and so on ...





## Example (4)

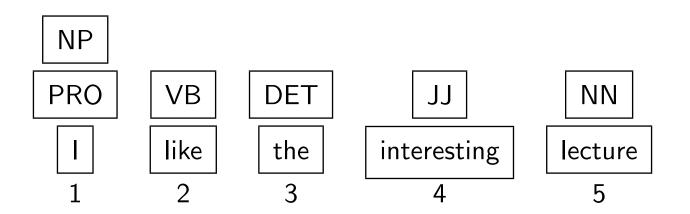
Try to apply a non-terminal rule to the first word The only matching rule is  $\mathsf{NP}\to\mathsf{PRO}$ 





# Example (5)

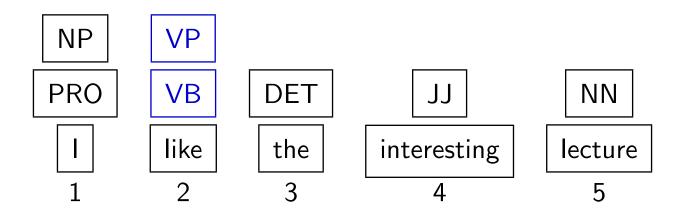
Recurse: try to apply a non-terminal rule to the first word No rule matches





# Example (6)

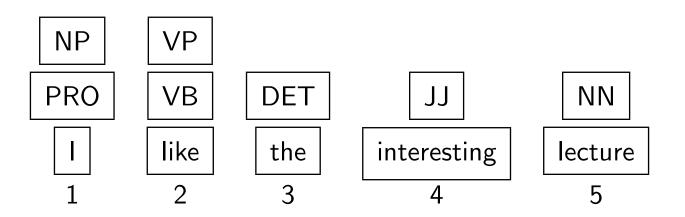
Try to apply a non-terminal rule to the second word The only matching rule is VP  $\rightarrow$  VB No recursion possible, no additional rules match





# Example (7)

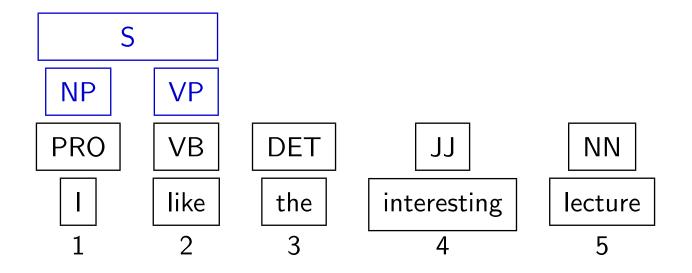
Try to apply a non-terminal rule to the third word No rule matches





## Example (8)

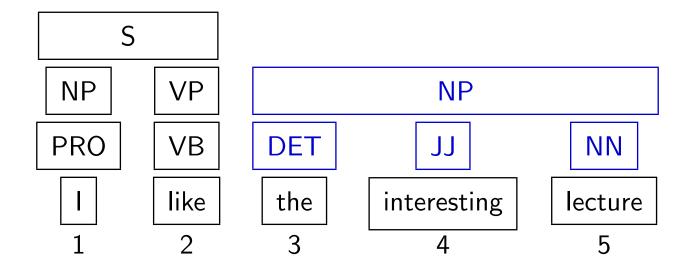
Try to apply a non-terminal rule to the first two words The only matching rule is  $S \rightarrow NP VP$ No other rules match for **spans** of two words





# Example (9)

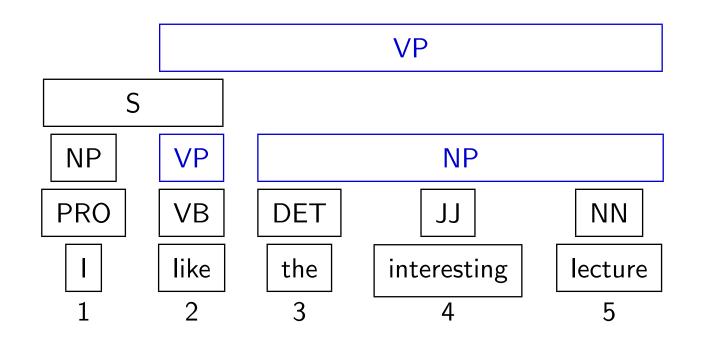
One rule matches for a span of three words:  $\mathsf{NP} \to \mathsf{DET}\ \mathsf{JJ}\ \mathsf{NN}$ 





# Example (10)

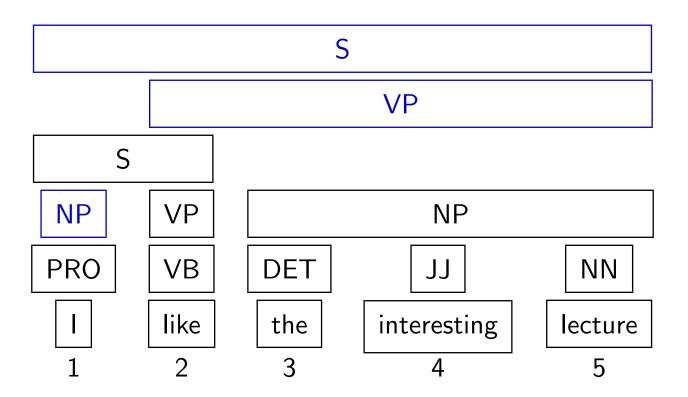
One rule matches for a span of four words:  $\mathsf{VP} \to \mathsf{VP} \ \mathsf{NP}$ 





# Example (11)

One rule matches for a span of five words:  $\mathsf{S} \to \mathsf{NP} \ \mathsf{VP}$ 





# **CYK** algorithm for binarized grammars

- for all words  $w_i$ : // terminal rules - for all rules  $A \rightarrow w_i$ : add new chart entry A at span [i, i]- for length = 1 to sentence length n / / non-terminal rules - for start = 1 to n - (length - 1)end = start + length - 1- for middle = start to end - 1: // binary rules for all non-terminals X in [start, middle]: for all non-terminals Y in [middle + 1, end]: for all rules  $A \to X Y^{\cdot}$ add new chart entry A at position [start, end]- for all non-terminals X in [start, end]: // unary rules for all rules  $A \to X$ :

add new chart entry A at position [start, end]