#### Number theory

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## Division

#### Definition

If a and b are integers with  $a \neq 0$ , then a divides b if there exists an integer c such that b = ac

#### Theorem

Let a, b, c be integers, where  $a \neq 0$ 

- 1. If a|b and a|c, then a|(b+c)
- 2. If a|b, then a|bc
- 3. If a|b and b|c, then a|c

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#### Proof

- a|b ⇔ ∃k<sub>b</sub>. b = k<sub>b</sub> ⋅ a and a|c ⇔ ∃k<sub>c</sub>. c = k<sub>c</sub> ⋅ a. But then b+c = (k<sub>b</sub> + k<sub>c</sub>) ⋅ a which by definition implies that a|(b+c)
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- 5.  $a|b \Leftrightarrow \exists k_b, b = k_b \cdot a$  and  $b|c \Leftrightarrow \exists k_c, c = k_c \cdot b$ . But then  $c = k_c \cdot k_b \cdot a$  which by definition implies that  $a|c^{(\frac{1}{2}) \cdot (\frac{1}{2})} \stackrel{<}{=} \frac{2}{2} \frac{2}{2}$

# **Division algorithm**

#### Theorem

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If a is an integer and d a positive integer, then there are unique integers q and r, with  $0 \le r < d$ , such that a = dq + r

<u>Proof</u> (by contradiction) Assume  $\exists q_1, q_2, r_1, r_2$  such that  $a = dq_1 + r_1$ ,  $a = dq_2 + r_2$ , and  $(q_1, r_1) \neq (q_2, r_2)$ . But then,

$$d=\frac{r_1-r_2}{q_2-q_1}$$

Now since  $0 \le r_1, r_2 < m$ , it must be that  $-d < r_1 - r_2 < d$ . But since  $q_1, q_2 \in \mathbb{Z}$ , it necessarily is the case that

$$-d < \frac{r_1 - r_2}{q_2 - q_1} < d$$

Which contradicts our hypothesis that  $d = \frac{r_1 - r_2}{q_2 - q_1}$ .

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## **Congruence relation**

#### Definition

If a and b are integers and m is a positive integer, then a is congruent to b modulo m, denoted  $a \equiv b \pmod{m}$ , iff m|(a - b)

# A theorem on congruences

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#### <u>Proof</u>

( $\Leftarrow$ ) If  $a \equiv b \pmod{m}$ , then by the definition of congruence m|(a-b). Hence, there is an integer k such that a-b=km and equivalently a=b+km

(⇒) If there is an integer k such that a = b + km, then km = a - b. Hence, m|(a - b) and  $a \equiv b \pmod{m}$ 

# Congruences of sums, differences, and products

#### Theorem

Let m be a positive integer. If 
$$a \equiv b \pmod{m}$$
 and  $c \equiv d \pmod{m}$ , then  $a + c \equiv b + d \pmod{m}$ ,  $a - c \equiv b - d \pmod{m}$ , and  $ac \equiv bd \pmod{m}$ 

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#### Proof

Since  $a \equiv b \pmod{m}$  and  $c \equiv d \pmod{m}$ , by the Theorem above there are integers s and t with b = a + sm and d = c + tm. Therefore, b + d = (a + sm) + (c + tm) = (a + c) + m(s + t), and bd = (a + sm)(c + tm) = ac + m(at + cs + stm). Hence,  $a + c \equiv b + d \pmod{m}$  and  $ac \equiv bd \pmod{m}$ 

#### Corollary

Let m be a positive integer and let a and b be integers. Then

- $(a+b) \mod m = ((a \mod m) + (b \mod m)) \mod m$
- $ab \mod m = ((a \mod m)(b \mod m)) \mod m$

### Arithmetic modulo m

• Let 
$$\mathbb{Z}_m = \{0, 1, \cdots, m-1\}$$

- Theoperation +<sub>m</sub> is defined as a +<sub>m</sub> b = (a + b) mod m. This is addition modulo m
- The operation  $\cdot_m$  is defined as  $a \cdot_m b = (a \cdot b) \mod m$ . This is multiplication modulo m
- Using these operations is said to be doing arithmetic modulo *m*

 $\begin{array}{l} \hline \text{Example Find } 7+_{11} 9 \ \text{and } 7\cdot_{11} 9 \\ \hline \hline \text{Solution} \ \text{Using the definitions above:} \\ 7+_{11} 9=(7+9) \ \text{mod } 11=16 \ \text{mod } 11=5 \ \text{and} \\ 7\cdot_{11} 9=(7\cdot 9) \ \text{mod } 11=63 \ \text{mod } 11=8 \end{array}$ 

#### Arithmetic modulo m

The operations  $+_m$  and  $\cdot_m$  satisfy many of the same properties as ordinary addition and multiplication

Closure If  $a, b \in \mathbb{Z}_m$ , then  $a +_m b$  and  $a \cdot_m b$  belong to  $\mathbb{Z}_m$ 

Associativity If  $a, b, c \in \mathbb{Z}_m$ , then  $(a +_m b) +_m c = a +_m (b +_m c)$ and  $(a \cdot_m b) \cdot_m c = a \cdot_m (b \cdot_m c)$ 

Commutativity If  $a, b \in \mathbb{Z}_m$ , then  $a +_m b = b +_m a$  and  $a \cdot_m b = b \cdot_m a$ 

Identity elements The elements 0 and 1 are identity elements for addition and multiplication modulo m, respectively. If  $a \in \mathbb{Z}_m$  then  $a +_m 0 = a$  and  $a \cdot_m 1 = a$ 

Additive inverses If  $0 \neq a \in \mathbb{Z}_m$ , then m - a is the additive inverse of a modulo m. Moreover, 0 is its own additive inverse  $a +_m (m - a) = 0$  and  $0 +_m 0 = 0$ Distributivity If  $a, b, c \in \mathbb{Z}_m$ , then

$$a \cdot_m (b +_m c) = (a \cdot_m b) +_m (a \cdot_m c) \text{ and }$$

$$(a +_m b) \cdot_m c = (a \cdot_m c) +_m (b \cdot_m c)$$

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Not all integers have an inverse mod m

#### **Primes**

#### Definition

A positive integer p > 1 is called prime iff the only positive factors of p are 1 and p. Otherwise it is called composite

#### Theorem (Fundamental Theorem of Arithmetic)

Every positive integer greater than 1 can be written uniquely as a prime or as the product of its prime factors, written in order of nondecreasing size.

Proof by induction (see slides on induction)

Example  $765 = 3 \cdot 3 \cdot 5 \cdot 17 = 3^2 \cdot 5 \cdot 17$ 

# There are infinitely many primes - Euclid (325-265 BCE)

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<u>Lemma</u> Every natural number greater than one is either prime or it has a prime divisor

<u>Proof</u> Suppose towards a contradiction that there are only finitely many primes  $p_1, p_2, p_3, \ldots, p_k$ . Consider the number  $q = p_1 p_2 p_3 \ldots p_k + 1$ , the product of all the primes plus one. By hypothesis q cannot be prime because it is strictly larger than all the primes. Thus, by the lemma, it has a prime divisor, p. Because  $p_1, p_2, p_3, \ldots, p_k$  are all the primes, p must be equal to one of them, so p is a divisor of their product. So we have that p divides  $p_1 p_2 p_3 \ldots p_k$ , and p divides q, but that means p divides their difference, which is 1. Therefore  $p \leq 1$ . Contradiction. Therefore there are infinitely many primes.

# The Sieve of Eratosthenes (276-194 BCE)

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How to find all primes between 2 and n?

A **very inefficient** method of determining if a number *n* is prime

Try every integer  $i \leq \sqrt{n}$  and see if *n* is divisible by *i*:

- 1. Write the numbers  $2, \ldots, n$  into a list. Let i := 2
- 2. Remove all strict multiples of *i* from the list
- 3. Let k be the smallest number present in the list s.t. k > i. Then let i := k
- 4. If  $i > \sqrt{n}$  then stop else go to step 2

Testing if a number is prime can be done efficiently in polynomial time [Agrawal-Kayal-Saxena 2002], i.e., polynomial in the number of bits used to describe the input number. Efficient randomized tests had been available previously.

### **Greatest common divisor**

#### Definition

Let  $a, b \in \mathbb{Z} - \{0\}$ . The largest integer d such that d|a and also d|b is called the greatest common divisor of a and b. It is denoted by gcd(a, b)

Example gcd(24, 36) = 12

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Example gcd(24, 36) = 12

#### Definition

The integers a and b are relatively prime (coprime) iff gcd(a, b) = 1

Example 17 and 22 (Note that 22 is not a prime)

### **Gcd by Prime Factorizations**

Suppose that the prime factorizations of a and b are:

$$a = p_1^{a_1} p_2^{a_2} \cdots p_n^{a_n}$$
  $b = p_1^{b_1} p_2^{b_2} \cdots p_n^{b_n}$ 

where each exponent is a nonnegative integer (possibly zero). Then

$$gcd(a, b) = p_1^{\min(a_1, b_1)} p_2^{\min(a_2, b_2)} \cdots p_n^{\min(a_n, b_n)}$$

This number clearly divides *a* and *b*. No larger number can divide both *a* and *b*. Proof by contradiction and the prime factorization of a postulated larger divisor.

Factorization is a very inefficient method to compute gcd. The Euclidian algorithm is much better.

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Euclidian algorithm
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```
algorithm gcd(x,y)
if y = 0
then return(x)
else return(gcd(y,x mod y))
```

Thed Euclidian algorithm relies on the fact that  $\forall x, y \in \mathbb{Z}. x > y \rightarrow gcd(x, y) = gcd(y, x \mod y)$ 

# Euclidian algorithm (proof of correctness)

#### Lemma

Let a = bq + r, where a, b, q, and r are integers. Then gcd(a, b) = gcd(b, r)

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#### Proof

( $\Rightarrow$ ) Suppose that *d* divides both *a* and *b*. Then *d* also divides a - bq = r. Hence, any common divisor of *a* and *b* must also be a common divisor of *b* and *r* ( $\Leftarrow$ ) Suppose that *d* divides both *b* and *r*. Then *d* also divides bq + r = a. Hence, any common divisor of *b* and *r* must also be a common divisor of *a* and *b*.

Therefore, gcd(a, b) = gcd(b, r)

Theorem

Let m, x be positive integers. gcd(m, x) = 1 iff x has a multiplicative inverse modulo m (and it is unique (modulo m)).

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<u>Proof</u> ( $\Rightarrow$ ) Consider the sequence of *m* numbers 0, x, 2x, ...(m-1)x. We first show that these are all distinct modulo *m*.

To verify the above claim, suppose that  $ax \mod m = bx \mod m$ ) for two distinct values a, b in the range  $0 \le a, b \le m-1$ . Then we would have  $(a - b)x \equiv 0 \pmod{m}$ , or equivalently, (a - b)x = kmfor some integer k. But since x and m are relatively prime, it follows that a - b must be an integer multiple of m. This is not possible since a, b are distinct non-negative integers less than m. Now, since there are only m distinct values modulo m, it must then be the case that  $ax \equiv 1 \pmod{m}$  for exactly one a (modulo m). This a is the unique multiplicative inverse.

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# Gcd as a linear combination

Theorem (Bézout's theorem)

If x and y are positive integers, then there exist integers a and b such that gcd(x, y) = ax + by (Proof in exercises of Section 5.2)

Example 
$$2 = \gcd(6, 14) = (-2) \cdot 6 + 1 \cdot 14$$

```
Extended Euclidian algorithm
The Bézout coefficients can be computed as follows:
algorithm extended-gcd(x,y)
if y = 0
then return(x, 1, 0)
else
(d, a, b) := extended-gcd(y, x mod y)
return((d, b, a - (x div y) * b))
```

# The multiplicative group $\mathbb{Z}_m^*$

#### Definition

Let  $\mathbb{Z}_m^* = \{x \mid 1 \le x < m \text{ and } gcd(x, m) = 1\}$ . Together with multiplication modulo m, this is called the multiplicative group modulo m. It is closed, associative, has a neutral element (namely 1) and every element has an inverse.

## Fermat's little theorem

#### Theorem

If p is prime and p  $\not|a$ , then  $a^{p-1} \equiv 1 \pmod{p}$ . Furthermore, for every integer a we have  $a^p \equiv a \pmod{p}$ 

Fermat's little theorem is useful in computing the remainders modulo p of large powers of integers

#### Example Find 7<sup>222</sup> mod 11

By Fermat's little theorem, we know that  $7^{10} \equiv 1 \pmod{11}$ , and so  $(7^{10})^k \equiv 1 \pmod{11}$  for every positive integer k. Therefore,  $7^{222} = 7^{22 = \cdot 10 + 2} = (7^{10})^{22} \cdot 7^2 \equiv 122 \cdot 49 \equiv 5 \pmod{11}$ . Hence,  $7^{222} \mod 11 = 5$ 





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# **RSA:** key generation

- Choose two distinct prime numbers *p* and *q*. Prime integers can be efficiently found using a primality test
- Let n = pq and k = (p-1)(q-1). (In particular,  $k = |\mathbb{Z}_n|$ )
- Choose an integer e such that 1 < e < k and gcd(e, k) = 1;</li>
   i.e. e and k are coprime
- (n, e) is released as the public key
- Let d be the multiplicative inverse of e modulo k, i.e. de ≡ 1 (mod k). (Computed using the extendedE uclidean algorithm)
- (n, d) is the private key and kept secret

# **RSA:** encryption and decryption

Alice transmits her public key (n, e) to Bob and keeps the private key secret

**Encryption** If Bob wishes to send message *M* to Alice.

- 1. He turns M into an integer m, such that  $0 \le m < n$  by using an agreed-upon reversible protocol known as a padding scheme
- 2. He computes the ciphertext *c* corresponding to  $c = m^e \mod n$ . This can be done quickly using the method of exponentiation by squaring.
- 3. Bob transmits c to Alice.

**Decryption** Alice can recover m from c by

- 1. Using her private key exponent d via computing  $m = c^d \mod n$
- 2. Given *m*, she can recover the original message *M* by reversing the padding scheme

#### **RSA: correctness of decryption**

Given that  $c = m^e \mod n$ , is  $m = c^d \mod n$ ?

$$c^d = (m^e)^d \equiv m^{ed} \pmod{n}$$

By construction, d and e are each others multiplicative inverses modulo k, i.e.  $ed \equiv 1 \pmod{k}$ . Also k = (p-1)(q-1). Thus ed - 1 = h(p - 1)(q - 1) for some integer h. We consider m<sup>ed</sup> mod p If *p ∦m* then  $m^{ed} = m^{h(p-1)(q-1)}m = (m^{p-1})^{h(q-1)}m \equiv 1^{h(q-1)}m \equiv m \pmod{p}$ (by Fermat's little theorem) Otherwise  $m^{ed} \equiv 0 \equiv m \pmod{p}$ Symmetrically,  $m^{ed} \equiv m \pmod{q}$ Since p, q are distinct primes, we have  $m^{ed} \equiv m \pmod{pq}$ . Since n = pq, we have  $c^d m^{ed} \equiv m \pmod{n}$