

Data Intensive Linguistics — Lecture 9

Parsing (I): Context-free grammars and chart parsing

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A simple sentence

I like the interesting lecture

Syntactic relations

I like the interesting lecture
 PRO VB DET JJ NN

- The adjective *interesting* gives more information about the noun *lecture*
- The determiner *the* says something about the noun *lecture*
- The noun *lecture* is the object of the verb *like*, specifying *what* is being liked
- The pronoun *I* is the subject of the verb *like*, specifying *who* is doing the liking

Dependency structure (2)

The dependencies may also be **labeled** with the type of dependency

I	like	the	interesting	lecture
PRO	VB	DET	JJ	NN
↓		↓	↓	↓
subject		adjunct	adjunct	object
↓		↓	↓	↓
like		lecture	lecture	like

The path so far

- Originally, we treated language as a *sequence of words*
 → n-gram language models
- Then, we introduced the notion of *syntactic properties of words*
 → part-of-speech tags
- Now, we look at *syntactic relations* between words
 → syntax trees

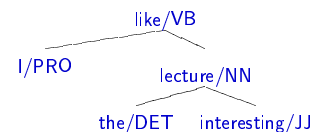
Part-of-speech tags

I like the interesting lecture
 PRO VB DET JJ NN

Dependency structure

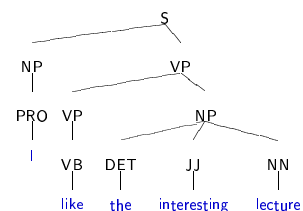
I	like	the	interesting	lecture
PRO	VB	DET	JJ	NN
↓		↓	↓	↓
like		lecture	lecture	like

This can also be visualized as a **dependency tree**:



Phrase-structure tree

A popular grammar formalism is **phrase structure grammar**. Internal nodes combine leaf nodes into phrases, such as *noun phrases (NP)*

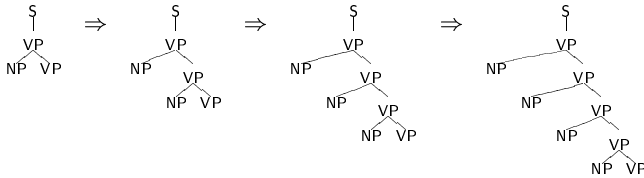


Building phrase-structure trees

- Our task for this week: **parsing**
 - given*: an input sentence with part-of-speech tags
 - wanted*: the right syntax tree for it
- Formalism: **context-free grammars**
 - non-terminal nodes** such as *NP*, *S* appear inside the tree
 - terminal nodes** such as *like*, *lecture* appear at the leaves of the tree
 - rules** such as $NP \rightarrow DET\ JJ\ NN$

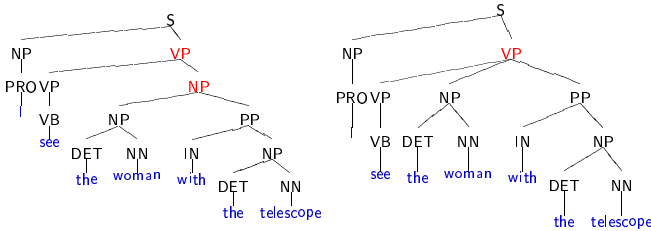
Recursion

Rules can be applied **recursively**, for example the rule $VP \rightarrow NP\ VP$



Why is parsing hard?

Prepositional phrase attachment: Who has the *telescope*?



CYK Parsing

- We have input sentence: *I like the interesting lecture*
- We have a set of context-free rules:
 - $S \rightarrow NP\ VP$, $NP \rightarrow PRO$, $PRO \rightarrow I$, $VP \rightarrow VP\ NP$, $VP \rightarrow VB$, $VB \rightarrow like$
 - $NP \rightarrow DET\ JJ\ NN$, $DET \rightarrow the$, $JJ \rightarrow interesting$, $NN \rightarrow lecture$
- Cocke-Younger-Kasami (CYK)** parsing
 - a **bottom-up** parsing algorithm
 - uses a **chart** to store intermediate result

Applying the rules

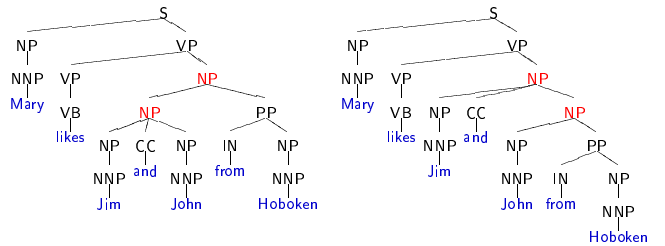
Input	Rule	Output
S	$S \rightarrow NP\ VP$	NP VP
NP VP	$NP \rightarrow PRO$	PRO VP
PRO VP	$PRO \rightarrow I$	I VP
I VP	$VP \rightarrow VP\ NP$	I VP NP
I VP NP	$VP \rightarrow VB$	I VB
I VB NP	$VB \rightarrow like$	I like NP
I like NP	$NP \rightarrow DET\ JJ\ NN$	I like DET JJ NN
I like DET JJ NN	$DET \rightarrow the$	I like the JJ NN
I like the JJ NN	$JJ \rightarrow interesting$	I like the interesting NN
I like the interesting NN	$NN \rightarrow lecture$	I like the interesting lecture

Context-free grammars in context

- Chomsky hierarchy** of formal languages (terminals in caps, non-terminal lowercase)
 - regular**: only rules of the form $A \rightarrow a$, $A \rightarrow B$, $A \rightarrow Ba$ (or $A \rightarrow aB$)
Cannot generate languages such as $a^n b^n$
 - context-free**: left-hand side of rule has to be single non-terminal, anything goes on right hand-side. Cannot generate $a^n b^n c^n$
 - context-sensitive**: rules can be restricted to a particular context, e.g. $\alpha A \beta \rightarrow \alpha A c \beta$, where α and β are strings of terminal and non-terminals
- Moving up the hierarchy, languages are more expressive and parsing becomes computationally more expensive
- Is natural language context-free?

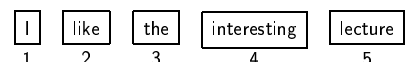
Why is parsing hard?

Scope: Is *Jim* also from *Hoboken*?



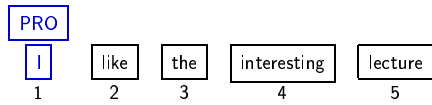
Example

Initialize chart with the words



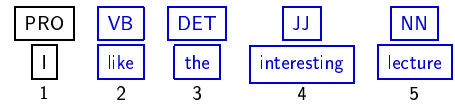
Example (2)

Apply first terminal rule $PRO \rightarrow I$



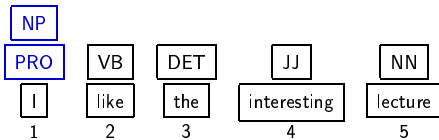
Example (3)

... and so on ...



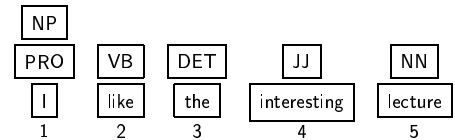
Example (4)

Try to apply a non-terminal rule to the first word
The only matching rule is $NP \rightarrow PRO$



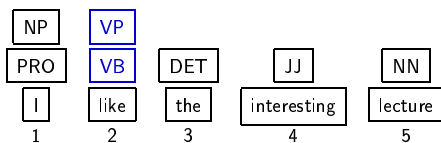
Example (5)

Recurse: try to apply a non-terminal rule to the first word
No rule matches



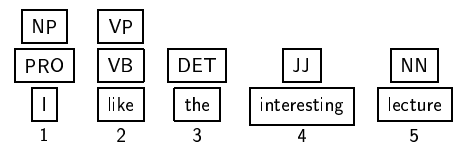
Example (6)

Try to apply a non-terminal rule to the second word
The only matching rule is $VP \rightarrow VB$
No recursion possible, no additional rules match



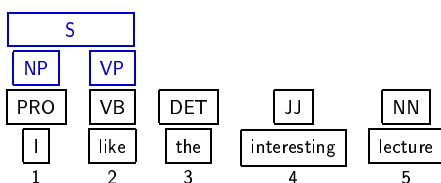
Example (7)

Try to apply a non-terminal rule to the third word
No rule matches



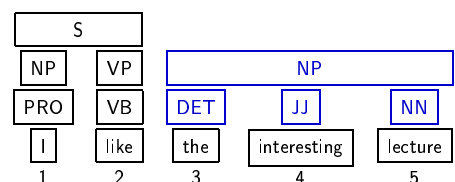
Example (8)

Try to apply a non-terminal rule to the first two words
The only matching rule is $S \rightarrow NP VP$
No other rules match for **spans** of two words



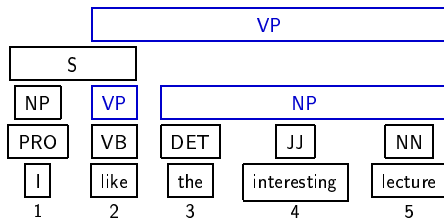
Example (9)

One rule matches for a span of three words: $NP \rightarrow DET JJ NN$



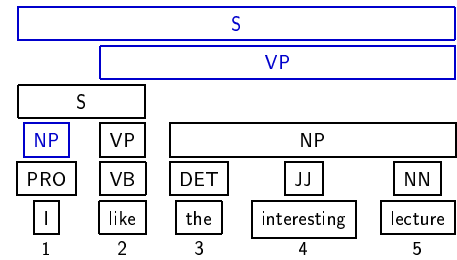
Example (10)

One rule matches for a span of four words: $VP \rightarrow VP NP$



Example (11)

One rule matches for a span of five words: $S \rightarrow NP VP$



CYK algorithm for binarized grammars

- for all words w_i : // terminal rules
- for all rules $A \rightarrow w_i$: add new chart entry A at span $[i, i]$
- for $length = 1$ to sentence length n // non-terminal rules
- for $start = 1$ to $n - (length - 1)$
 - $end = start + length - 1$
 - for $middle = start$ to $end - 1$: // binary rules
 - for all non-terminals X in $[start, middle]$:
 - for all non-terminals Y in $[middle + 1, end]$:
 - for all rules $A \rightarrow X Y$:
 - add new chart entry A at position $[start, end]$
 - for all non-terminals X in $[start, end]$: // unary rules
 - for all rules $A \rightarrow X$:
 - add new chart entry A at position $[start, end]$