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# Data Intensive Linguistics — Lecture 2

## Introduction (II): Probability and Information Theory

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## Recap

- Given word counts we can estimate a probability distribution:

$$P(w) = \frac{\text{count}(w)}{\sum_{w'} \text{count}(w')}$$

- Another useful concept is conditional probability

$$p(w_2|w_1)$$

- Chain rule:

$$p(w_1, w_2) = p(w_1) p(w_2|w_1)$$

- Bayes rule:

$$p(x|y) = \frac{p(y|x) p(x)}{p(y)}$$

# Expectation

- We introduced the concept of a random variable  $X$

$$\text{prob}(X = x) = p(x)$$

- Example: Roll of a dice. There is a  $\frac{1}{6}$  chance that it will be 1, 2, 3, 4, 5, or 6.

- We define the **expectation**  $E(X)$  of a random variable as:

$$E(X) = \sum_x p(x) x$$

- Roll of a dice:

$$E(X) = \frac{1}{6} \times 1 + \frac{1}{6} \times 2 + \frac{1}{6} \times 3 + \frac{1}{6} \times 4 + \frac{1}{6} \times 5 + \frac{1}{6} \times 6 = 3.5$$

# Variance

- **Variance** is defined as

$$\text{Var}(X) = E((X - E(X))^2) = E(X^2) - E^2(X)$$

$$\text{Var}(X) = \sum_x p(x) (x - E(X))^2$$

- Intuitively, this is a measure how far events diverge from the mean (expectation)
- Related to this is **standard deviation**, denoted as  $\sigma$ .

$$\text{Var}(X) = \sigma^2$$

$$E(X) = \mu$$

## Variance (2)

- Roll of a dice:

$$\begin{aligned} \text{Var}(X) &= \frac{1}{6}(1 - 3.5)^2 + \frac{1}{6}(2 - 3.5)^2 + \frac{1}{6}(3 - 3.5)^2 \\ &\quad + \frac{1}{6}(4 - 3.5)^2 + \frac{1}{6}(5 - 3.5)^2 + \frac{1}{6}(6 - 3.5)^2 \\ &= \frac{1}{6}((-2.5)^2 + (-1.5)^2 + (-0.5)^2 + 0.5^2 + 1.5^2 + 2.5^2) \\ &= \frac{1}{6}(6.25 + 2.25 + 0.25 + 0.25 + 2.25 + 6.25) \\ &= 2.917 \end{aligned}$$

## Standard distributions

- **Uniform:** all events equally likely
  - $\forall x, y : p(x) = p(y)$
  - example: roll of one dice
- **Binomial:** a series of trials with only two outcomes
  - probability  $p$  for each trial, occurrence  $r$  out of  $n$  times:  
$$b(r; n, p) = \binom{n}{r} p^r (1 - p)^{n-r}$$
  - a number of coin tosses

## Standard distributions (2)

- **Normal**: common distribution for continuous values
  - value in the range  $[-\text{inf}, x]$ , given expectation  $\mu$  and standard deviation  $\sigma$ :  
$$n(x; \mu, \sigma) = \frac{1}{\sqrt{2\pi}\sigma} e^{-(x-\mu)^2/(2\sigma^2)}$$
  - also called **Bell curve**, or **Gaussian**
  - examples: heights of people, IQ of people, tree heights, ...

## Estimation revisited

- We introduced last lecture an estimation of probabilities based on frequencies:

$$P(w) = \frac{\text{count}(w)}{\sum_{w'} \text{count}(w')}$$

- Alternative view: Bayesian: what is the most likely model given the data

$$p(M|D)$$

- Model and data are viewed as random variables
  - model  $M$  as random variable
  - data  $D$  as random variable



## Bayesian estimation

- Reformulation of  $p(M|D)$  using Bayes rule:

$$p(M|D) = \frac{p(D|M) p(M)}{p(D)}$$

$$\operatorname{argmax}_M p(M|D) = \operatorname{argmax}_M p(D|M) p(M)$$

- $p(M|D)$  answers the question: What is the most likely model given the data
- $p(M)$  is a prior that prefers certain models (e.g. simple models)
- The frequentist estimation of word probabilities  $p(w)$  is the same as Bayesian estimation with a uniform prior (no bias towards a specific model), hence it is also called the **maximum likelihood estimation**

# Entropy

- An important concept is **entropy**:

$$H(X) = \sum_x -p(x) \log_2 p(x)$$

- A measure for the degree of disorder

## Entropy example

One event

$$p(a) = 1$$

$$\begin{aligned} H(X) &= -1 \log_2 1 \\ &= 0 \end{aligned}$$

## Entropy example

2 equally likely events:

$$p(a) = 0.5$$

$$p(b) = 0.5$$

$$H(X) = -0.5 \log_2 0.5 - 0.5 \log_2 0.5$$

$$= -\log_2 0.5$$

$$= 1$$

## Entropy example

4 equally likely events:

$$p(a) = 0.25$$

$$p(b) = 0.25$$

$$p(c) = 0.25$$

$$p(d) = 0.25$$

$$\begin{aligned} H(X) &= -0.25 \log_2 0.25 - 0.25 \log_2 0.25 \\ &\quad - 0.25 \log_2 0.25 - 0.25 \log_2 0.25 \\ &= -\log_2 0.25 \\ &= 2 \end{aligned}$$

## Entropy example

4 equally likely events, one more likely than the others:

$$p(a) = 0.7$$

$$p(b) = 0.1$$

$$p(c) = 0.1$$

$$p(d) = 0.1$$

$$H(X) = -0.7 \log_2 0.7 - 0.1 \log_2 0.1$$

$$- 0.1 \log_2 0.1 - 0.1 \log_2 0.1$$

$$= -0.7 \log_2 0.7 - 0.3 \log_2 0.1$$

$$= -0.7 \times -0.5146 - 0.3 \times -3.3219$$

$$= 0.36020 + 0.99658$$

$$= 1.35678$$

## Entropy example

4 equally likely events, one much more likely than the others:

$$\begin{aligned} (X) \quad & H(X) = -0.97 \log_2 0.97 - 0.01 \log_2 0.01 \\ & \quad - 0.01 \log_2 0.01 - 0.01 \log_2 0.01 \\ & = -0.97 \log_2 0.97 - 0.03 \log_2 0.01 \\ & = -0.97 \times -0.04394 - 0.03 \times -6.6439 \\ & = 0.04262 + 0.19932 \\ & = 0.24194 \end{aligned}$$
$$\begin{aligned} p(a) &= 0.97 \\ p(b) &= 0.01 \\ p(c) &= 0.01 \\ p(d) &= 0.01 \end{aligned}$$

## Intuition behind entropy

- A good model has low entropy

→ it is more certain about outcomes

- For instance a translation table

$e$	$f$	$p(e f)$
the	der	0.8
that	der	0.2

is better than

$e$	$f$	$p(e f)$
the	der	0.02
that	der	0.01
...	...	...

- A lot of statistical estimation is about reducing entropy



## Information theory and entropy

- Assume that we want to encode a sequence of events  $X$
- Each event is encoded by a sequence of bits
- For example
  - Coin flip: heads = 0, tails = 1
  - 4 equally likely events: a = 00, b = 01, c = 10, d = 11
  - 3 events, one more likely than others: a = 0, b = 10, c = 11
  - Morse code: e has shorter code than q
- Average number of bits needed to encode  $X \geq$  entropy of  $X$

## The entropy of English

- We already talked about the probability of a word  $p(w)$
- But words come in sequence. Given a number of words in a text, can we guess the next word  $p(w_n|w_1, \dots, w_{n-1})$ ?
- Example: Newspaper article

## Entropy for letter sequences

Assuming a model with a limited window size

Model	Entropy
0th order	4.76
1st order	4.03
2nd order	2.8
human, unlimited	1.3