



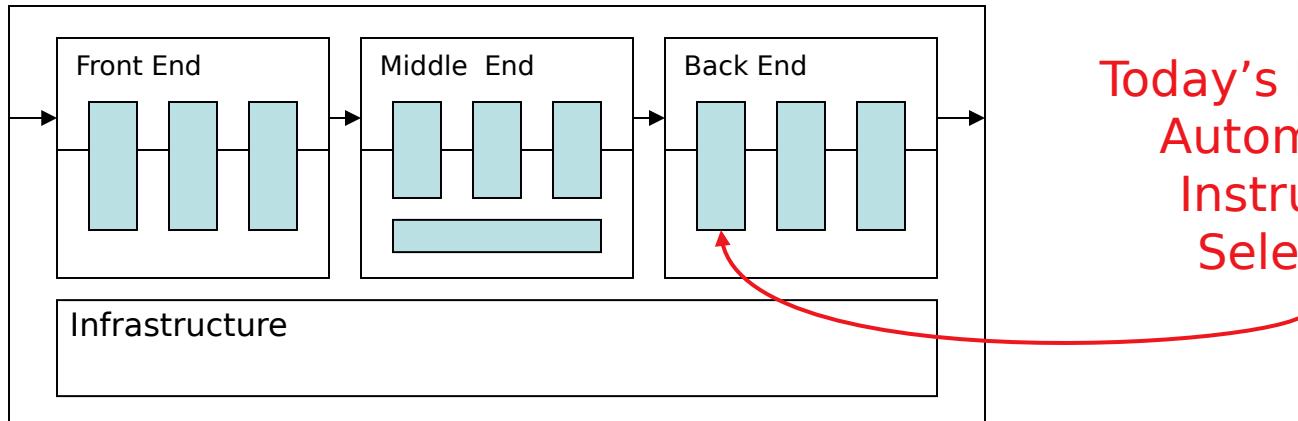
Instruction Selection: Peephole Matching

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The Problem

Writing a compiler is a lot of work

- Would like to reuse components whenever possible
- Would like to automate construction of components



Today's lecture:
Automating
Instruction
Selection

- Front end construction is largely automated
- Middle is largely hand crafted
- (Parts of) back end can be automated

Definitions

Instruction selection

- Mapping *IR* into assembly code
- Assumes a fixed storage mapping & code shape
- Combining operations, using address modes

Instruction scheduling

- Reordering operations to hide latencies
- Assumes a fixed program (*set of operations*)
- Changes demand for registers

Register allocation

- Deciding which values will reside in registers
- Changes the storage mapping, may add false sharing
- Concerns about placement of data & memory operations

The Problem

Modern computers (still) have many ways to do anything

Consider register-to-register copy in **ILOC**

- Obvious operation is $i2i \ r_i \Rightarrow r_j$
- Many others exist

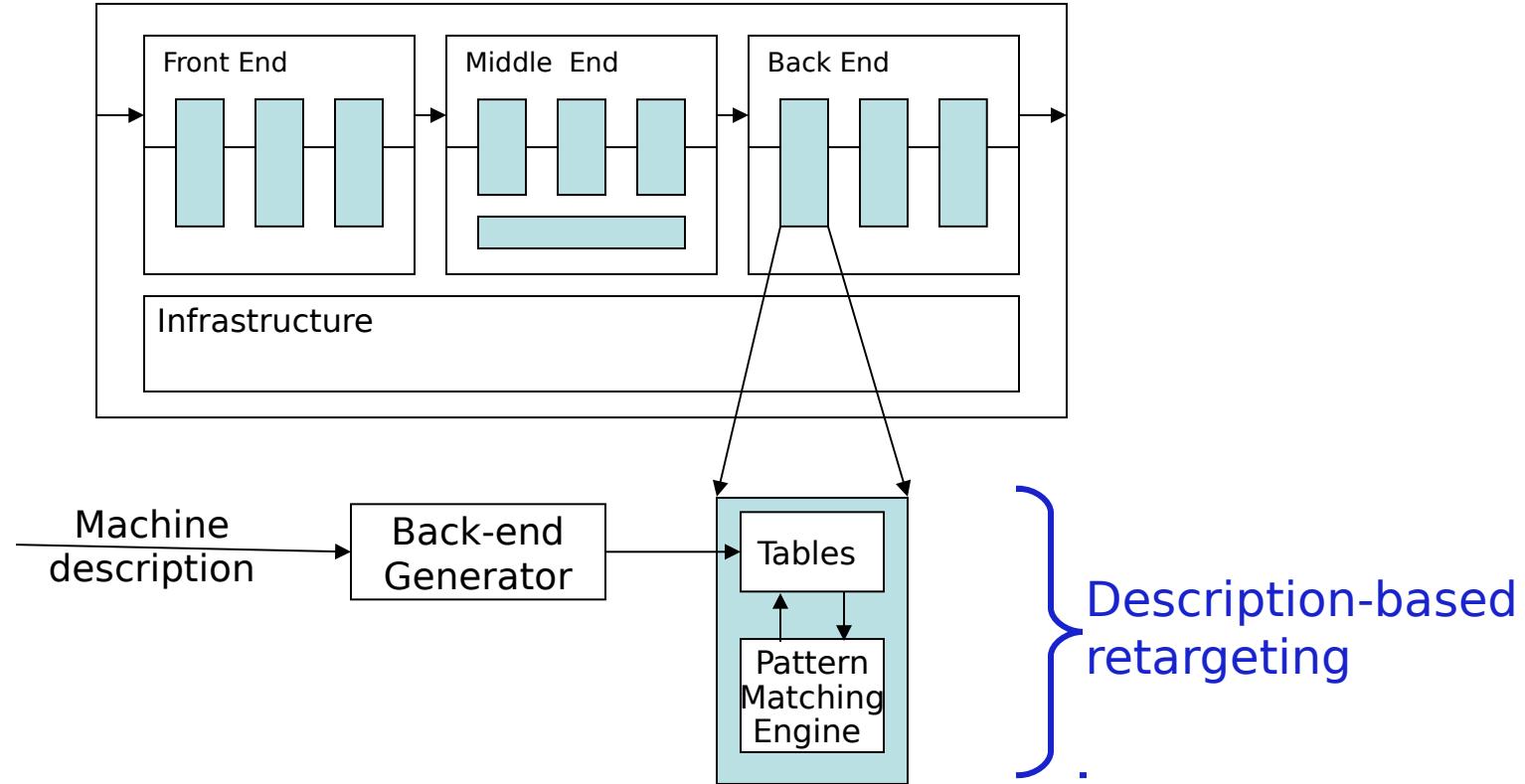
$\text{addI } r_i, 0 \Rightarrow r_j$	$\text{subI } r_i, 0 \Rightarrow r_j$	$\text{lshiftI } r_i, 0 \Rightarrow r_j$
$\text{multI } r_i, 1 \Rightarrow r_j$	$\text{divI } r_i, 1 \Rightarrow r_j$	$\text{rshiftI } r_i, 0 \Rightarrow r_j$
$\text{orI } r_i, 0 \Rightarrow r_j$	$\text{xorI } r_i, 0 \Rightarrow r_j$	$\dots \text{ and others } \dots$

- Human would ignore all of these
- Algorithm must look at all of them & find low-cost encoding
 - Take context into account *(busy functional unit?)*

And **ILOC** is an overly-simplified case

The Goal

Want to automate generation of instruction selectors



The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator ran quickly

How good was the code?

Tree	Treewalk Code	Desired Code								
<pre>graph TD; x[x] --> IDENT1["IDENT
<a,ARP,4>"]; x --> IDENT2["IDENT
<b,ARP,8>"]</pre>	<table><tbody><tr><td>loadl 4 $\Rightarrow r_5$</td></tr><tr><td>loadAO $r_0, r_5 \Rightarrow r_6$</td></tr><tr><td>loadl 8 $\Rightarrow r_7$</td></tr><tr><td>loadAO $r_0, r_7 \Rightarrow r_8$</td></tr><tr><td>mult $r_6, r_8 \Rightarrow r_9$</td></tr></tbody></table>	loadl 4 $\Rightarrow r_5$	loadAO $r_0, r_5 \Rightarrow r_6$	loadl 8 $\Rightarrow r_7$	loadAO $r_0, r_7 \Rightarrow r_8$	mult $r_6, r_8 \Rightarrow r_9$	<table><tbody><tr><td>loadAI $r_0, 4 \Rightarrow r_5$</td></tr><tr><td>loadAI $r_0, 8 \Rightarrow r_6$</td></tr><tr><td>mult $r_5, r_6 \Rightarrow r_7$</td></tr></tbody></table>	loadAI $r_0, 4 \Rightarrow r_5$	loadAI $r_0, 8 \Rightarrow r_6$	mult $r_5, r_6 \Rightarrow r_7$
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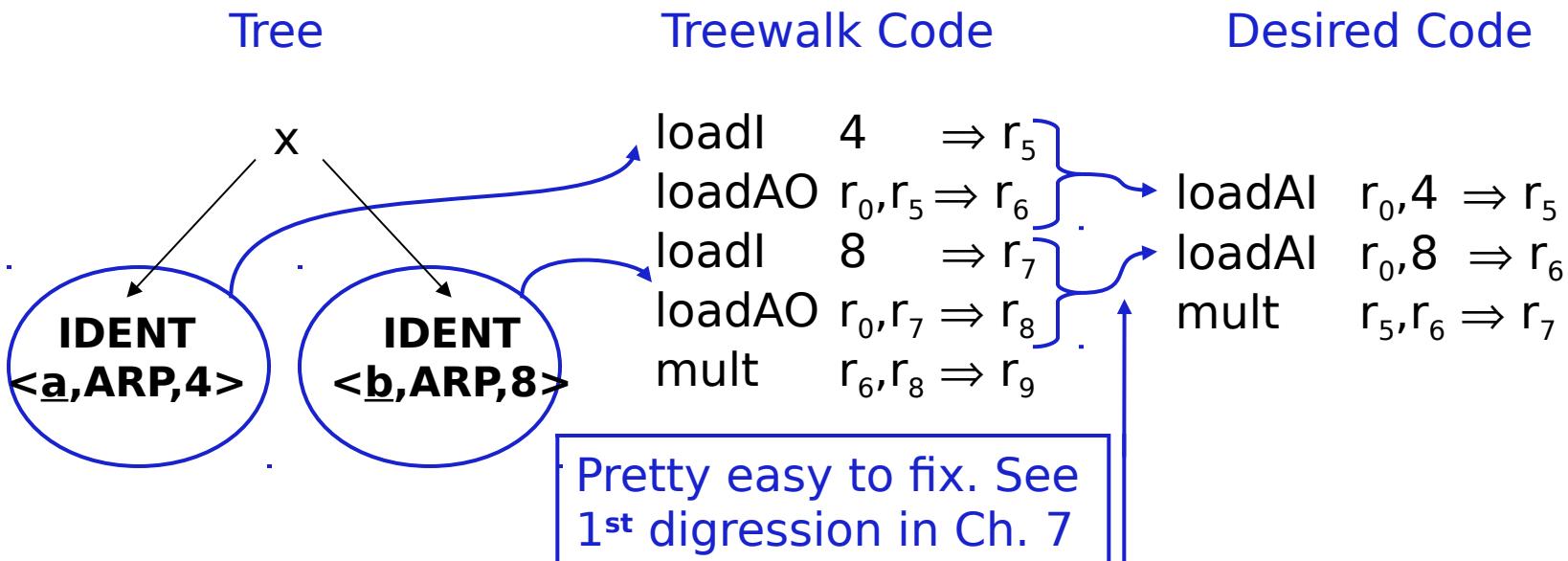
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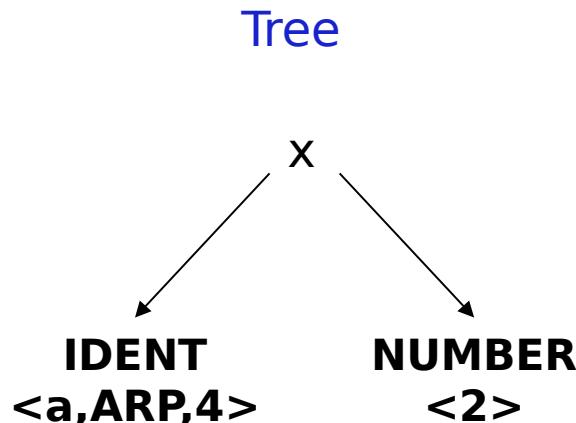
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
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How good was the code?



Treewalk Code

loadl	4	$\Rightarrow r_5$
loadAO	r_0, r_5	$\Rightarrow r_6$
loadl	2	$\Rightarrow r_7$
mult	r_6, r_7	$\Rightarrow r_8$

Desired Code

loadAl	$r_0, 4$	$\Rightarrow r_5$
multl	$r_5, 2$	$\Rightarrow r_7$

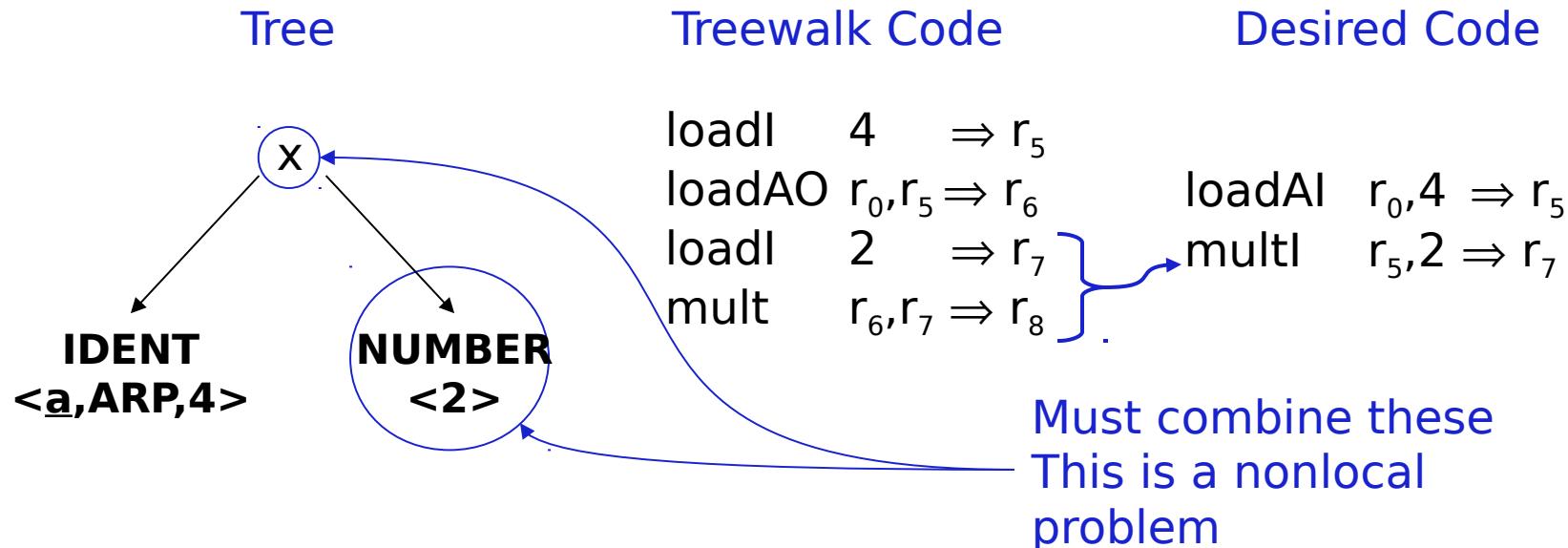
The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator ran quickly

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The Big Picture

Need pattern matching techniques

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How good was the code?

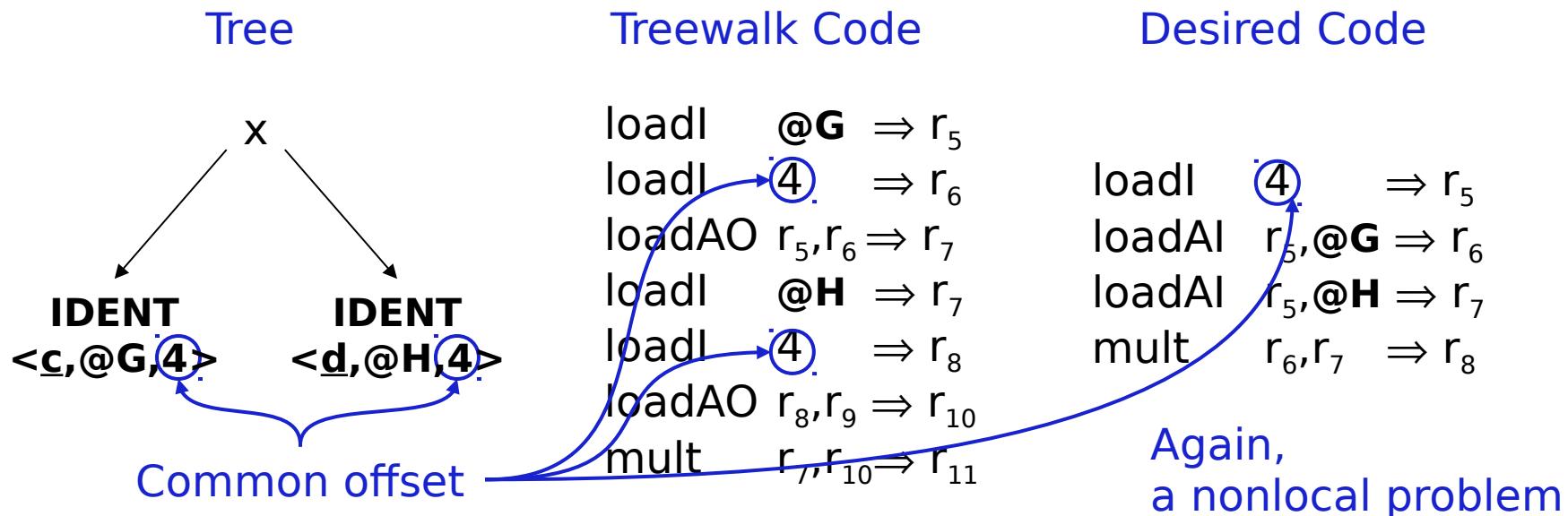
Tree	Treewalk Code	Desired Code
<pre>graph TD; x[x] --> IDENT1[IDENT<c,@G,4>]; x --> IDENT2[IDENT<d,@H,4>]</pre>	<pre>loadl @G ⇒ r₅ loadl 4 ⇒ r₆ loadAO r₅,r₆ ⇒ r₇ loadl @H ⇒ r₇ loadl 4 ⇒ r₈ loadAO r₈,r₉ ⇒ r₁₀ mult r₇,r₁₀ ⇒ r₁₁</pre>	<pre>loadl 4 ⇒ r₅ loadAI r₅,@G ⇒ r₆ loadAI r₅,@H ⇒ r₇ mult r₆,r₇ ⇒ r₈</pre>

The Big Picture

Need pattern matching techniques

- Must produce good code *(some metric for good)*
- Must run quickly

Our treewalk code generator met the second criteria
How did it do on the first ?



How do we perform this kind of matching ?

Tree-oriented IR suggests pattern matching on trees

- Tree-patterns as input, matcher as output
- Each pattern maps to a target-machine instruction sequence
- Use dynamic programming or bottom-up rewrite systems

Linear IR suggests using some sort of string matching

- Strings as input, matcher as output
- Each string maps to a target-machine instruction sequence
- Use text matching (Aho-Corasick) or peephole matching

In practice, both work well; matchers are quite different

Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load

Original code

storeAI $r_1 \Rightarrow r_0, 8$
loadAI $r_0, 8 \Rightarrow r_{15}$

Improved code

storeAI $r_1 \Rightarrow r_0, 8$
i2i $r_1 \Rightarrow r_{15}$

Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load
- Simple algebraic identities

Original code

addl $r_2, 0 \Rightarrow r_7$
mult $r_4, r_7 \Rightarrow r_{10}$

Improved code

mult $r_4, r_2 \Rightarrow r_{10}$

Peephole Matching

- Basic idea
- Compiler can discover local improvements locally
 - Look at a small set of adjacent operations
 - Move a “peephole” over code & search for improvement
- Classic example was store followed by load
- Simple algebraic identities
- Jump to a jump

Original code

$$\begin{array}{ll} \text{jumpl} & \rightarrow L_{10} \\ L_{10}: \text{jumpl} & \rightarrow L_{11} \end{array}$$

Improved code

$$L_{10}: \text{jumpl} \rightarrow L_{11}$$

Peephole Matching

Implementing it

- Early systems used limited set of hand-coded patterns
- Window size ensured quick processing

Modern peephole instruction selectors

(Davidson)

- Break problem into three tasks



- Apply symbolic interpretation & simplification systematically

Peephole Matching

Expander

- Turns IR code into a low-level IR (LLIR) such as RTL*
- Operation-by-operation, template-driven rewriting
- LLIR form includes all direct effects (e.g., *setting cc*)
- Significant, albeit constant, expansion of size



*RTL = Register transfer language

Peephole Matching

Simplifier

- Looks at LLIR through window and rewrites it
- Uses forward substitution, algebraic simplification, local constant propagation, and dead-effect elimination
- Performs local optimization within window



- This is the heart of the peephole system
 - Benefit of peephole optimization shows up in this step

Peephole Matching

Matcher

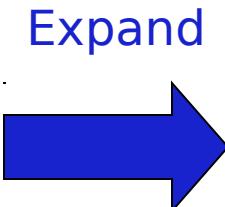
- Compares simplified LLIR against a library of patterns
- Picks low-cost pattern that captures effects
- Must preserve LLIR effects, may add new ones (*e.g., set cc*)
- Generates the assembly code output



Example

Original IR Code

OP	Arg ₁	Arg ₂	Result
mult	2	Y	t ₁
sub	x	t ₁	w



LLIR Code

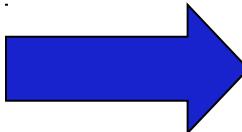
```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Example

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
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r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
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MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(r0 + @y)  
r14 ← 2 × r13  
r17 ← MEM(r0 + @x)  
r18 ← r17 - r14  
MEM(r0 + @w) ← r18
```

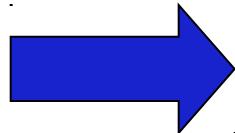
Example

LLIR Code

```
r13 ← MEM(r0 + @y)  
r14 ← 2 × r13  
r17 ← MEM(r0 + @x)  
r18 ← r17 - r14
```

```
MEM(r0 + @w) ← r18
```

Match



ILoc Code

```
loadAl r0,@y ⇒ r13  
multI 2 × r13 ⇒ r14  
loadAl r0,@x ⇒ r17  
sub r17 - r14 ⇒ r18  
storeAl r18 ⇒ r0,@w
```

- Introduced all memory operations & temporary names
- Turned out pretty good code

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

Steps of the Simplifier *(3-operation window)*

LLIR Code

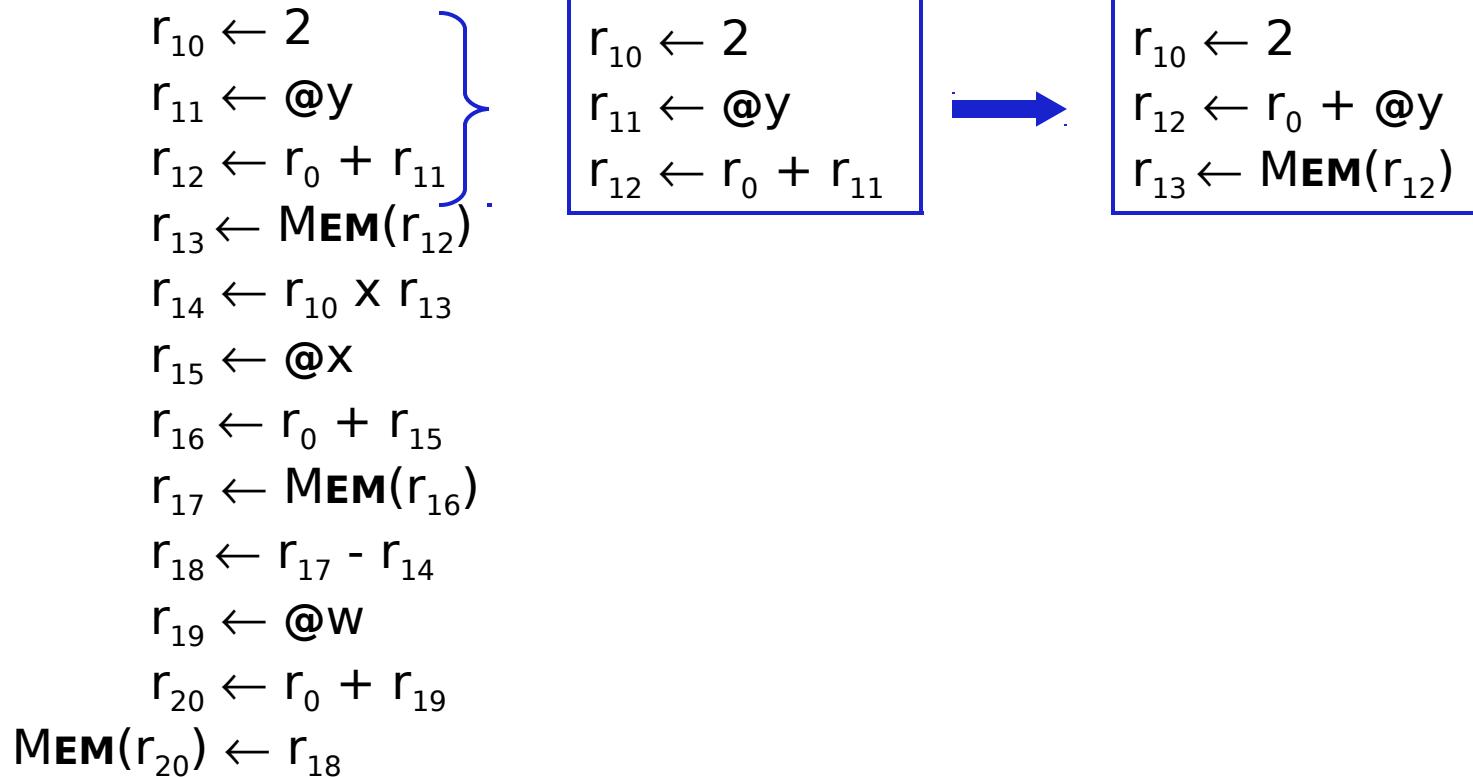
```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

The diagram illustrates a 3-operation window. The first three assignments ($r_{10} \leftarrow 2$, $r_{11} \leftarrow @y$, $r_{12} \leftarrow r_0 + r_{11}$) are grouped by a brace on the left and enclosed in a blue box on the right. This indicates that the window has shifted one position to the right after these three operations.

Steps of the Simplifier

(3-operation window)

LLIR Code

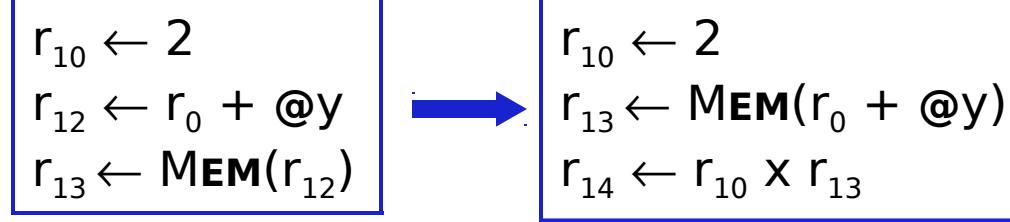


Steps of the Simplifier

(3-operation window)

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow r_0 + r_{11}$
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} \times r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow r_0 + r_{15}$
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow r_0 + r_{19}$
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$



Steps of the Simplifier

(3-operation window)

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

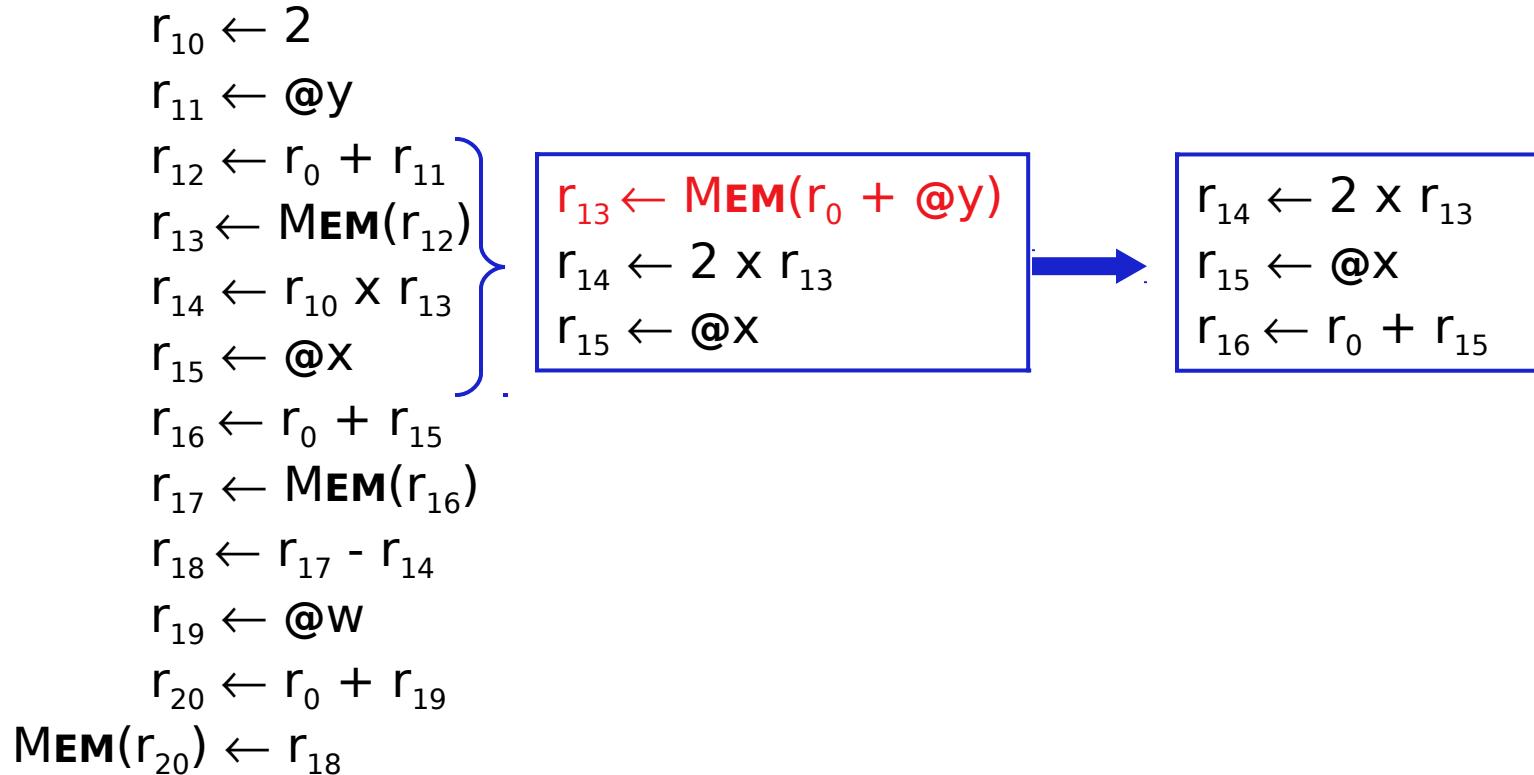
1st op rolling out of window

```
r10 ← 2  
r13 ← MEM(r0 + @y)  
r14 ← r10 × r13
```

```
r13 ← MEM(r0 + @y)  
r14 ← 2 × r13  
r15 ← @x
```

Steps of the Simplifier *(3-operation window)*

LLIR Code



Steps of the Simplifier *(3-operation window)*

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow r_0 + r_{11}$
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} \times r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow r_0 + r_{15}$
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow r_0 + r_{19}$
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

```
graph LR; A["r14 ← 2 × r13  
r15 ← @x  
r16 ← r0 + r15"] --> B["r14 ← 2 × r13  
r16 ← r0 + @x  
r17 ← MEM(r16)"]
```

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

$r_{14} \leftarrow 2 \times r_{13}$
 $r_{16} \leftarrow r_0 + @x$
 $r_{17} \leftarrow \text{MEM}(r_{16})$



$r_{14} \leftarrow 2 \times r_{13}$
 $r_{17} \leftarrow \text{MEM}(r_0 + @x)$
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Steps of the Simplifier *(3-operation window)*

LLIR Code

$r_{10} \leftarrow 2$
 $r_{11} \leftarrow @y$
 $r_{12} \leftarrow r_0 + r_{11}$
 $r_{13} \leftarrow \mathbf{MEM}(r_{12})$
 $r_{14} \leftarrow r_{10} \times r_{13}$
 $r_{15} \leftarrow @x$
 $r_{16} \leftarrow r_0 + r_{15}$
 $r_{17} \leftarrow \mathbf{MEM}(r_{16})$
 $r_{18} \leftarrow r_{17} - r_{14}$
 $r_{19} \leftarrow @w$
 $r_{20} \leftarrow r_0 + r_{19}$
 $\mathbf{MEM}(r_{20}) \leftarrow r_{18}$

The diagram illustrates the simplification process. A brace groups the operations $r_{14} \leftarrow 2 \times r_{13}$, $r_{15} \leftarrow @x$, $r_{16} \leftarrow r_0 + r_{15}$, and $r_{17} \leftarrow \mathbf{MEM}(r_0 + @x)$. A blue box encloses this group, with a red annotation $r_{14} \leftarrow 2 \times r_{13}$ indicating a modification. An arrow points from this box to another blue box containing the simplified code $r_{17} \leftarrow \mathbf{MEM}(r_0 + @x)$, $r_{18} \leftarrow r_{17} - r_{14}$, and $r_{19} \leftarrow @w$.

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
r16 ← r0 + r15
r17 ← MEM(r16)
r18 ← r17 - r14
r19 ← @w
r20 ← r0 + r19
MEM(r20) ← r18
```

The diagram illustrates the simplification process. A brace on the left groups operations 16-19. A red box highlights the assignment $r_{17} \leftarrow \text{MEM}(r_0 + @x)$. An arrow points from this red box to a blue box containing the intermediate steps $r_{18} \leftarrow r_{17} - r_{14}$ and $r_{19} \leftarrow @w$. Another arrow points from this blue box to a final blue box containing the simplified code $r_{18} \leftarrow r_{17} - r_{14}$, $r_{19} \leftarrow @w$, and $r_{20} \leftarrow r_0 + r_{19}$.

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18.
```



```
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19
```



```
r18 ← r17 - r14  
r20 ← r0 + @w  
MEM(r20) ← r18
```

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
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r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

{ r₁₈ ← r₁₇ - r₁₄
 r₂₀ ← r₀ + @w
 MEM(r₂₀) ← r₁₈



 r₁₈ ← r₁₇ - r₁₄
 MEM(r₀ + @w) ← r₁₈

Steps of the Simplifier *(3-operation window)*

LLIR Code

```
r10 ← 2
r11 ← @y
r12 ← r0 + r11
r13 ← MEM(r12)
r14 ← r10 × r13
r15 ← @x
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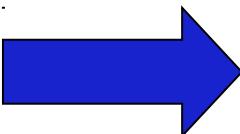
$r_{18} \leftarrow r_{17} - r_{14}$
 $\text{MEM}(r_0 + @w) \leftarrow r_{18}$

Example

LLIR Code

```
r10 ← 2  
r11 ← @y  
r12 ← r0 + r11  
r13 ← MEM(r12)  
r14 ← r10 × r13  
r15 ← @x  
r16 ← r0 + r15  
r17 ← MEM(r16)  
r18 ← r17 - r14  
r19 ← @w  
r20 ← r0 + r19  
MEM(r20) ← r18
```

Simplify



LLIR Code

```
r13 ← MEM(r0 + @y)  
r14 ← 2 × r13  
r17 ← MEM(r0 + @x)  
r18 ← r17 - r14  
MEM(r0 + @w) ← r18
```

Making It All Work

Details

- LLIR is largely machine independent (RTL)
- Target machine described as LLIR → ASM pattern
- Actual pattern matching
 - Use a hand-coded pattern matcher (gcc)
 - Turn patterns into grammar & use LR parser (VPO)
- Several important compilers use this technology
- It seems to produce good portable instruction selectors

Key strength appears to be late low-level optimization