

Computer Programming: Skills & Concepts (CP1)

Intro to Practical 3: Travelling Salesman Problem

9th November 2010

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Complexity of problems

We have already encountered problems with different complexity:

- ▶ search through unsorted array: linear (ie, $O(n)$)
- ▶ binary search through sorted array: \log (ie, $O(\lg(n))$)
- ▶ BubbleSort: $O(n^2)$
- ▶ MergeSort: $O(n \lg(n))$

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Travelling Salesman Problem (TSP)

A well-known theoretical and practical problem:

- ▶ a salesman has to visit a number of cities
- ▶ what is the shortest route to visit all cities and return home?

Properties of the problem:

- ▶ hard to solve for large number of cities
- ▶ instance of a *NP-complete* problem

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NP-complete?

- ▶ For some problems, no polynomial time solution is known — $O(n^c)$ for some constant c . One class of these problems is called NP-complete (NP = non-polynomial).
- ▶ There may be polynomial solutions, but nobody found them so far.
- ▶ If *efficient* solution of a problem is not possible, we resort to *heuristics* that give us approximate solutions.

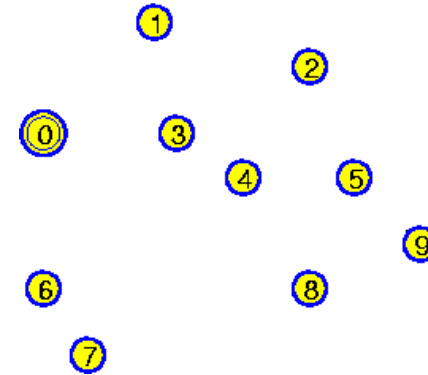
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Other NP-hard problems

- ▶ Knapsack problem: given a set of whole numbers a_1, \dots, a_n , and an upper bound K find a subset of the numbers whose sum is of *maximum value*, subject to being no more than K .
eg, for 2, 4, 9, 11, 14 and $K = 25$, the subset is $\{2, 9, 14\}$
- ▶ Minesweeper: is a given configuration "possible"?

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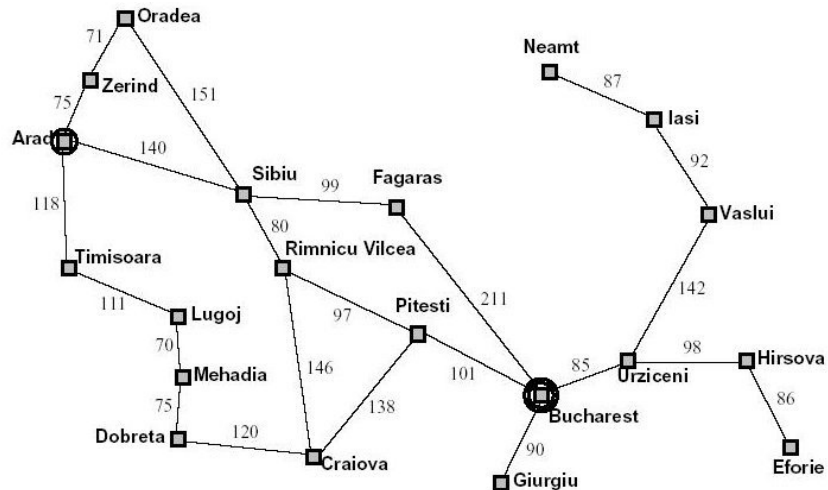
Simplified: Euclidean TSP



All connections are straight lines. How do we find the shortest path?

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Example TSP: Romania



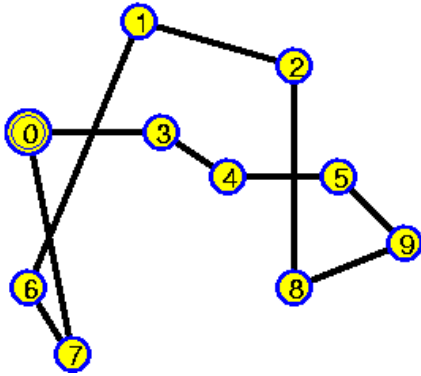
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Greedy heuristic

- ▶ start at some point
- ▶ go to closest not visited city

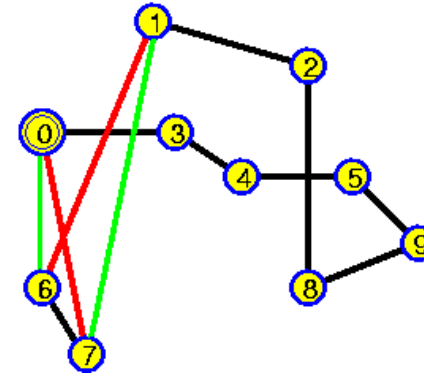
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Greedy heuristic: result



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Swap 6,7

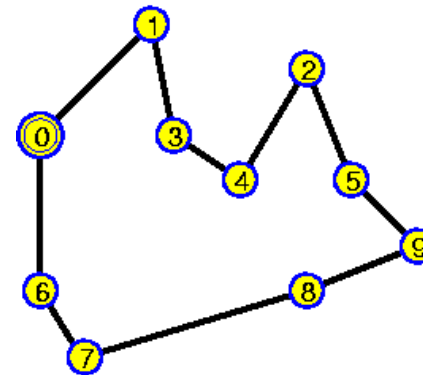


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Improving the solution

- ▶ Swap neighboring cities, if it shortens path

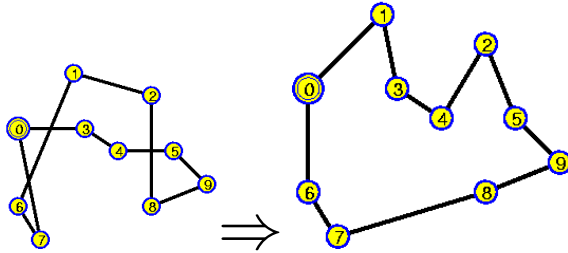
Locally Optimal solution



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Other improvements?



What other improvements can be made?

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Practical 3

- ▶ Part A: capture positions of cities (from mouse clicks), and store them all in an array. Write a function to compute the length of a given tour.
- ▶ Part B: implement swap heuristic.
- ▶ Part C: implement 2-opt heuristic (more powerful).
- ▶ Part D: implement greedy heuristic.
- ▶ Part E: do better, with almost no extra work?

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