Computer Programming: Skills & Concepts (CP1) Searching and sorting

1st November 2010

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Correctness of a Program

- How can you show that a program is correct?
- Similar to a mathematical proof: show that certain statements are true at all times in the program ("Invariants")
- Brain teaser: Is it possible to write a program that checks any other program, if it is correct?

Power of a number

```
int Power(int n, int k)
/* Assumes k >= 0. Returns n^k: n raised to the power k. */
ł
   int p = 1, i = k;
  /* Precondition: i >= 0 */
  while (i > 0) {
     /* Invariant: i >= 0 AND p * n^i == n^k */
     p *= n;
     --i;
  }
  /* p = n^k */
  return p;
}
```

Example: n = 3, k = 4. The answer should be $3^4 = 81$. The computation progesses as follows. Initially, i = k and p = 1. Note that $p \times n^i$ is invariant!

	i	р	$p \times n^i$
Initial	4	1	$1 \times 3^4 = 81$
Iteration 1	3	3	$3 \times 3^3 = 81$
Iteration 2	2	9	$9 \times 3^2 = 81$
Iteration 3	1	27	$27 \times 3^1 = 81$
Iteration 4	0	81	$81 \times 3^0 = 81$

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Searching an array

typedef enum {FALSE, TRUE} Bool_t;

```
Bool_t LinearSearch(int n, int a[], int sKey)
/* Returns TRUE iff (if and only if) sKey is contained
* in the array, i.e., there exists an index i with 0 <= i < n
* such that a[i] == sKey.
*/
{
    int i;
    for (i = 0; i < n; ++i) {
        if (a[i] == sKey) return TRUE;
    }
    return FALSE;
}</pre>
```

variant:

 Could use return type int with #DEFINE for TRUE, FALSE (see BinarySearch)

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Binary search

Sometimes we quickly want to find an entry in an array. It helps if the array is sorted. How do you search for a name in a telephone book?

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Binary search

```
int BinarySearch(int n, int a[], int sKey)
/* Assumes the elements of the array a are in ascending order.
 * Returns TRUE iff sKey is contained in the array, i.e.,
 * there exists an index i with 0 <= i < n and a[i] == sKey.
 */
{
    int i, j, m;
    i = 0;
    j = n - 1;
    /* Precondition: a[0] <= a[1] <= ... <= a[n-1] */</pre>
```

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```
while (i < j) {
    /* Invariant: i <= j AND
    * if sKey is in a[0:n-1] then sKey is in a[i:j] */
    m = (i + j)/2;
    if (sKey <= a[m])
        j = m;
    else
        i = m + 1;
    }
    /* EITHER a[i] == sKey OR sKey is not in a[0:n-1] */
    return a[i] == sKey;</pre>
```

Note how we return true/false ...

}

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Running time

The *(worst-case)* running time of a function (or *algorithm*) is defined to be the maximum number of steps that might be performed by the program as a function of the *input size*.

- ▶ For functions which take an array (of some basic type) as the input, the length of the array (n in lots of our examples) is usually taken to represent size.
- The running time of Linear Search is Θ(n) (ie, around c ⋅ n for some constant c), and the running time of Binary search is Θ(lg(n)) (proportional to lg(n)).
- This size is conceptual and not what gets measured by the sizeof command in C (sizeof applied to an array is not length)

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Measuring running time on a machine

```
#include <time.h>
Bool_t flag = FALSE;
int a[880000];
double start, stop, t;
...
start = clock();
flag = LinearSearch(a, 880000, -5);
stop = clock();
t = (stop-start)/CLOCKS_PER_SEC;
printf("Time spent by Linear Search was %lf seconds.\n", t);
...
```

Machines getting faster ... on DICE, this method measures *both* linsearch (and binsearch) at 0.000000 secs on arrays up to length 400000!

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Sorting

Given an array of integers (or any *comparable* type), re-arrange the array so that the items appear in increasing order.

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Bubble sort

```
"Proto loop"
for (i = n - 1; i >= 1; --i) {
    /* Rearrange the contents of array elements a[0], ..., a[i],
    * so that the largest value appears in element a[i].
    */
}
```

"Method":

- Find the largest item, and move it to the end;
- repeat for 2nd largest item, and so on ...

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Bubble sort (cont'd)

The task of rearranging the contents of array elements $a[0], a[1], \ldots, a[i]$ so that the largest value appears in element a[i], may be handled by the following simple loop:

```
for (j = 0; j < i; ++j) {
    if (a[j] > a[j+1]) swap(&a[j], &a[j+1]);
}
```

(The largest value supposedly "bubbles" up the array into its appropriate position.)

Bubble sort code

```
/* Sorts a[0], a[1], ..., a[n-1] into ascending order. */
void BubbleSort(int a[], int n)
{
   int i, j;
  for (i = n - 1; i \ge 1; --i) {
      /* Invariant: The values in locations to the right of a[i]
       * are in their correct resting places: that is, they are
       * the n - i - 1 largest elements, and they are correctly
       * ordered among themselves.
       */
     for (j = 0; j < i; ++j) {
         if (a[j] > a[j+1]) swap(&a[j], &a[j+1]);
     }
  }
}
```

The function used above is the following function from lecture 11. void swap(int *a, int *b)

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Running time of Bubble Sort

The (worst case) running time of Bubble Sort is proportional to n^2 . why?

There are better sorting algorithms ... for example MergeSort or HeapSort run in time proportional to $n \lg(n)$.

 $(\lg(n) \text{ denotes "log to the base-2"})$

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A more dramatic example

Sometimes the gap between a good and bad algorithm can be dramatic. Consider the problem of testing whether an n-bit number is prime.

- The obvious brute force method requires $2^{n/2}$ integer divisions
 - why?
 - This is completely infeasible if n = 200 (say).
- On the other hand, a (non-obvious) algorithm for primality testing which take time *polynomial* in *n* was discovered in 2002 (Agrawal-Kayal-Saxena)

(Needed for RSA public-key cryptosystem.)