

## Computer Programming: Skills & Concepts (CP1)

### Searching and sorting

1st November 2010

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### Correctness of a Program

- ▶ How can you show that a program is correct?
- ▶ Similar to a mathematical proof: show that certain statements are true at all times in the program (“Invariants”)
- ▶ Brain teaser: Is it possible to write a program that checks any other program, if it is correct?

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### Power of a number

```
int Power(int n, int k)
/* Assumes k >= 0. Returns n^k: n raised to the power k. */
{
    int p = 1, i = k;
    /* Precondition: i >= 0 */
    while (i > 0) {
        /* Invariant: i >= 0 AND p * n^i == n^k */
        p *= n;
        --i;
    }
    /* p = n^k */
    return p;
}
```

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Example:  $n = 3$ ,  $k = 4$ . The answer should be  $3^4 = 81$ .

The computation progresses as follows. Initially,  $i = k$  and  $p = 1$ . Note that  $p \times n^i$  is invariant!

	$i$	$p$	$p \times n^i$
Initial	4	1	$1 \times 3^4 = 81$
Iteration 1	3	3	$3 \times 3^3 = 81$
Iteration 2	2	9	$9 \times 3^2 = 81$
Iteration 3	1	27	$27 \times 3^1 = 81$
Iteration 4	0	81	$81 \times 3^0 = 81$

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## Searching an array

```
typedef enum {FALSE, TRUE} Bool_t;

Bool_t LinearSearch(int n, int a[], int sKey)
/* Returns TRUE iff (if and only if) sKey is contained
 * in the array, i.e., there exists an index i with 0 <= i < n
 * such that a[i] == sKey.
 */
{
    int i;
    for (i = 0; i < n; ++i) {
        if (a[i] == sKey) return TRUE;
    }
    return FALSE;
}
```

*variant:*

- ▶ Could use return type int with #DEFINE for TRUE, FALSE (see BinarySearch)

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## Binary search

```
int BinarySearch(int n, int a[], int sKey)
/* Assumes the elements of the array a are in ascending order.
 * Returns TRUE iff sKey is contained in the array, i.e.,
 * there exists an index i with 0 <= i < n and a[i] == sKey.
 */
{
    int i, j, m;

    i = 0;
    j = n - 1;
    /* Precondition: a[0] <= a[1] <= ... <= a[n-1] */
```

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## Binary search

Sometimes we quickly want to find an entry in an array.  
It helps if the array is sorted.  
How do you search for a name in a telephone book?

```
while (i < j) {
    /* Invariant: i <= j AND
     * if sKey is in a[0:n-1] then sKey is in a[i:j] */
    m = (i + j)/2;
    if (sKey <= a[m])
        j = m;
    else
        i = m + 1;
}
/* EITHER a[i] == sKey OR sKey is not in a[0:n-1] */
return a[i] == sKey;
}
```

- ▶ Note how we return true/false ...

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## Running time

The (*worst-case*) *running time* of a function (or *algorithm*) is defined to be the maximum number of steps that might be performed by the program as a function of the *input size*.

- ▶ For functions which take an array (of some basic type) as the input, the length of the array ( $n$  in lots of our examples) is usually taken to represent size.
- ▶ The running time of Linear Search is  $\Theta(n)$  (ie, around  $c \cdot n$  for some constant  $c$ ), and the running time of Binary search is  $\Theta(\lg(n))$  (proportional to  $\lg(n)$ ).
- ▶ This *size* is conceptual and *not* what gets measured by the `sizeof` command in C (`sizeof` applied to an array is not length)

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## Sorting

Given an array of integers (or any *comparable* type), re-arrange the array so that the items appear in increasing order.

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## Measuring running time on a machine

```
#include <time.h>
Bool_t flag = FALSE;
int a[880000];
double start, stop, t;
...
start = clock();
flag = LinearSearch(a, 880000, -5);
stop = clock();
t = (stop-start)/CLOCKS_PER_SEC;
printf("Time spent by Linear Search was %lf seconds.\n", t);
...
```

Machines getting faster ... on DICE, this method measures *both* `linsearch` (and `binsearch`) at 0.000000 secs on arrays up to length 400000!

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## Bubble sort

“Proto loop”

```
for (i = n - 1; i >= 1; --i) {
    /* Rearrange the contents of array elements a[0], ..., a[i],
     * so that the largest value appears in element a[i].
     */
}
```

“Method”:

- ▶ Find the largest item, and move it to the end;
- ▶ **repeat** for 2nd largest item, and so on ...

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## Bubble sort (cont'd)

The task of rearranging the contents of array elements  $a[0], a[1], \dots, a[i]$  so that the largest value appears in element  $a[i]$ , may be handled by the following simple loop:

```
for (j = 0; j < i; ++j) {
    if (a[j] > a[j+1]) swap(&a[j], &a[j+1]);
}
```

(The largest value supposedly “bubbles” up the array into its appropriate position.)

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## Running time of Bubble Sort

The (worst case) running time of Bubble Sort is proportional to  $n^2$ . why?

There are better sorting algorithms . . . for example *MergeSort* or *HeapSort* run in time proportional to  $n \lg(n)$ .

( $\lg(n)$  denotes “log to the base-2”)

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## Bubble sort code

```
/* Sorts a[0], a[1], ..., a[n-1] into ascending order. */
void BubbleSort(int a[], int n)
{
    int i, j;
    for (i = n - 1; i >= 1; --i) {
        /* Invariant: The values in locations to the right of a[i]
         * are in their correct resting places: that is, they are
         * the n - i - 1 largest elements, and they are correctly
         * ordered among themselves.
         */
        for (j = 0; j < i; ++j) {
            if (a[j] > a[j+1]) swap(&a[j], &a[j+1]);
        }
    }
}
```

The function used above is the following function from lecture 11.

```
void swap(int *a, int *b)
```

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## A more dramatic example

Sometimes the gap between a good and bad algorithm can be dramatic. Consider the problem of testing whether an  $n$ -bit number is prime.

- ▶ The obvious brute force method requires  $2^{n/2}$  integer divisions
  - ▶ why?
  - ▶ This is completely infeasible if  $n = 200$  (say).
- ▶ On the other hand, a (non-obvious) algorithm for primality testing which take time *polynomial* in  $n$  was discovered in 2002 (Agrawal-Kayal-Saxena)

(Needed for RSA public-key cryptosystem.)

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