

## Computer Programming: Skills & Concepts (CP) Searching and sorting

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CP Lect 17 – slide 1 – Monday 13 November 2017

## Binary search

Sometimes we quickly want to find an entry in an array.

It helps if the array is sorted.

How do you search for a name in a telephone book?

Computers aren't so clever, so we do a simplified version:

Repeatedly chop the array in half to close in on where the element must be. E.g., to search for 17 in:

2 3 5 7 11 13 17 19 23 29

Find the mid-point:  $17 > 11$ , so narrow to right half:

Find the mid-point:  $17 \leq 19$ , so narrow to left half:

Find the mid-point:  $17 \leq 17$ , so narrow to left half:

(yes, we could stop here because we've found it...)

Find the mid-point:  $17 > 13$ , so narrow to right half:

Now we're left with an array of size 1, so either its element is 17 and we've found it, or 17 isn't there.

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## Searching an array

```
typedef enum {FALSE, TRUE} Bool_t;

Bool_t LinearSearch(int n, int a[], int sKey)
/* Returns TRUE iff (if and only if) sKey is contained in
 * the array, i.e., there exists an index i with  $0 \leq i < n$ 
 * such that  $a[i] == sKey$ .
 */
{
    int i;
    for (i = 0; i < n; ++i) {
        if (a[i] == sKey) return TRUE;
    }
    return FALSE;
}
```

variant:

- ▶ Could use return type int with #DEFINE for TRUE, FALSE (see BinarySearch)

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## Binary search

```
int BinarySearch(int n, int a[], int sKey)
/* Assumes: elements of array a are in ascending order.
 * Returns TRUE iff sKey is contained in the array, i.e.,
 * there exists an index i with  $0 \leq i < n$  and  $a[i] == sKey$ .
 */
{
    /* Precondition:  $(n > 0)$ 
     AND  $a[0] \leq a[1] \leq \dots \leq a[n-1]$  */

    int i, j, m;
    /* i will be the start of the sub-array
     * we're currently chopping;
     * j will be the end of it (its last element);
     * m will be the mid-point of it.
     */

    i = 0;
    j = n - 1;
}
```

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```

/* Invariant just before (re-)entering loop: i <= j AND
 * if sKey is in a[0:n-1] then sKey is in a[i:j] */
while (i < j) {
    m = (i + j)/2;
    if (sKey <= a[m]) {
        j = m;
    }
    else {
        i = m + 1;
    }
}
/* After exiting loop:
 * (i >= j), by i, j updates, means (i == j).
 * now EITHER a[i] == sKey OR sKey is not in a[0:n-1] */
return a[i] == sKey;
}

```

► Note how we return true/false ...

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## Measuring running time on a machine

```

#include <time.h>
Bool_t flag = FALSE;
int a[24000000];
clock_t start, stop;
double t;
...
start = clock();
flag = LinearSearch(a, 24000000, -5);
stop = clock();
t = ((double)(stop-start))/CLOCKS_PER_SEC;
printf("Time spent by Linear Search was %lf seconds.\n", t);
...

```

On my laptop:

Time spent by LinearSearch was 0.069064 seconds.

Time spent by BinarySearch was 0.000001 seconds.

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## Running time

The (*worst-case*) *running time* of a function (or *algorithm*) is defined to be the maximum number of steps that might be performed by the program as a function of the *input size*.

- For functions which take an array (of some basic type) as the input, the length of the array ( $n$  in lots of our examples) is usually taken to represent size.
- The running time of Linear Search proportional to  $n$  (i.e., around  $c \cdot n$  for some constant  $c$ ), and the running time of Binary Search is proportional to  $\lg(n)$ .

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## Sorting

Given an array of integers (or any *comparable* type), re-arrange the array so that the items appear in increasing order.

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## Bubble sort

'Pseudo-code'

```
for (i = n - 1; i >= 1; i--) {  
    /* Rearrange the contents of  
     * array elements a[0], ..., a[i],  
     * so that the largest value appears  
     * in element a[i].  
     */  
}
```

'Method':

- ▶ Find the largest item, and move it to the end;
- ▶ **repeat** for 2nd largest item, and so on ...

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## Bubble sort code

```
/* Sorts a[0], a[1], ..., a[n-1] into ascending order. */  
void BubbleSort(int a[], int n) {  
    int i, j;  
    for (i = n - 1; i >= 1; i--) {  
        /* Invariant: The values in locations to the right of  
         * a[i] are in their correct resting places: they are  
         * the (n - i - 1)-largest elements arranged in  
         * positions (i+1), ..., (n-1), in non-descending order  
        */  
        for (j = 0; j < i; j++) {  
            if (a[j] > a[j+1]) {  
                swap(&a[j], &a[j+1]);  
            }  
        }  
    }  
}
```

The swap function used above is the (correct) one from lab 5.

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## Bubble sort (cont'd)

The task of rearranging the contents of array elements  $a[0]$ ,  $a[1]$ , ...,  $a[i]$  so that the largest value appears in element  $a[i]$ , may be handled by the following simple loop:

```
for (j = 0; j < i; j++) {  
    if (a[j] > a[j+1]) {  
        swap(&a[j], &a[j+1]);  
    }  
}
```

(The largest value supposedly 'bubbles' up the array into its appropriate position.)

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## Running time of Bubble Sort

The (worst case) running time of Bubble Sort is proportional to  $n^2$ . **why?**

There are better sorting algorithms ... for example *MergeSort* or *HeapSort* run in time proportional to  $n \lg(n)$ .

For general purpose sorting, often use *QuickSort*, which runs in time around  $n \lg n$  in most cases, though in bad cases (**which?**) it can take  $n^2$ . Standard C systems provide QuickSort as *qsort*. Occasionally you might know that BubbleSort would be quicker in your application, and want to program it. Anything else is probably specialist.

*More about Bubble-Sort can be found in Section 6.7 of 'A Book on C'.*

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## Understanding your loops

These slides are logically small and green: for the mathematically and logically inclined only!

- ▶ How can you show that a program is correct?
- ▶ One way is to show that certain statements are true at all times in the program (*invariants*)
- ▶ In particular, to understand a complex while/for-loop, it's useful to know what remains true every time you go through it.
- ▶ For functions (or other blocks of code) we have *preconditions* (things *assumed* be true before) and *postconditions* (things which *will* be true afterwards given the preconditions).

We'll do a simple example now; then look (in your own time) at the comments in the searching and sorting code, and try to understand what they're saying about invariants.

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Example:  $n = 3, k = 4$ . The answer should be  $3^4 = 81$ .

The computation progresses as follows. Initially,  $i = k$  and  $p = 1$ . Note that  $p \times n^i$  is invariant!

```
/* Invariant before (re-)entering:
   i >= 0 AND p * n^i == n^k */
while (i > 0) {
    p *= n;
    --i;
}
/* After exiting loop: i <= 0 AND p = n^k */
return p;
```

	$i$	$p$	$p \times n^i$
Initial	4	1	$1 \times 3^4 = 81$
Iteration 1	3	3	$3 \times 3^3 = 81$
Iteration 2	2	9	$9 \times 3^2 = 81$
Iteration 3	1	27	$27 \times 3^1 = 81$
Iteration 4	0	81	$81 \times 3^0 = 81$

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## Power of a number

```
int Power(int n, int k)
/* Pre-condition: k >= 0. */
/* On-exit: returns n^k (n raised to the power k). */
{
    int p = 1, i = k;
    /* Invariant before (re-)entering:
       * i >= 0 AND p * n^i == n^k */
    while (i > 0) {
        p *= n;
        --i;
    }
    /* After exiting loop: i <= 0 AND p = n^k */
    return p;
}
```

**Warning:**  $n^k$  in the comments is maths notation, not C notation. In C, the  $\wedge$  symbol is the bitwise exclusive-or operator, something entirely different!

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## Reading material

Sections of 'A Book on C' that are relevant are:

- ▶ A good idea to refresh your memory of arrays (early sections of Chapter 6).
- ▶ Section 6.7 has a discussion of BubbleSort.

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