

# Chapter IV: Network Layer

UG3 Computer Communications & Networks  
(COMN)

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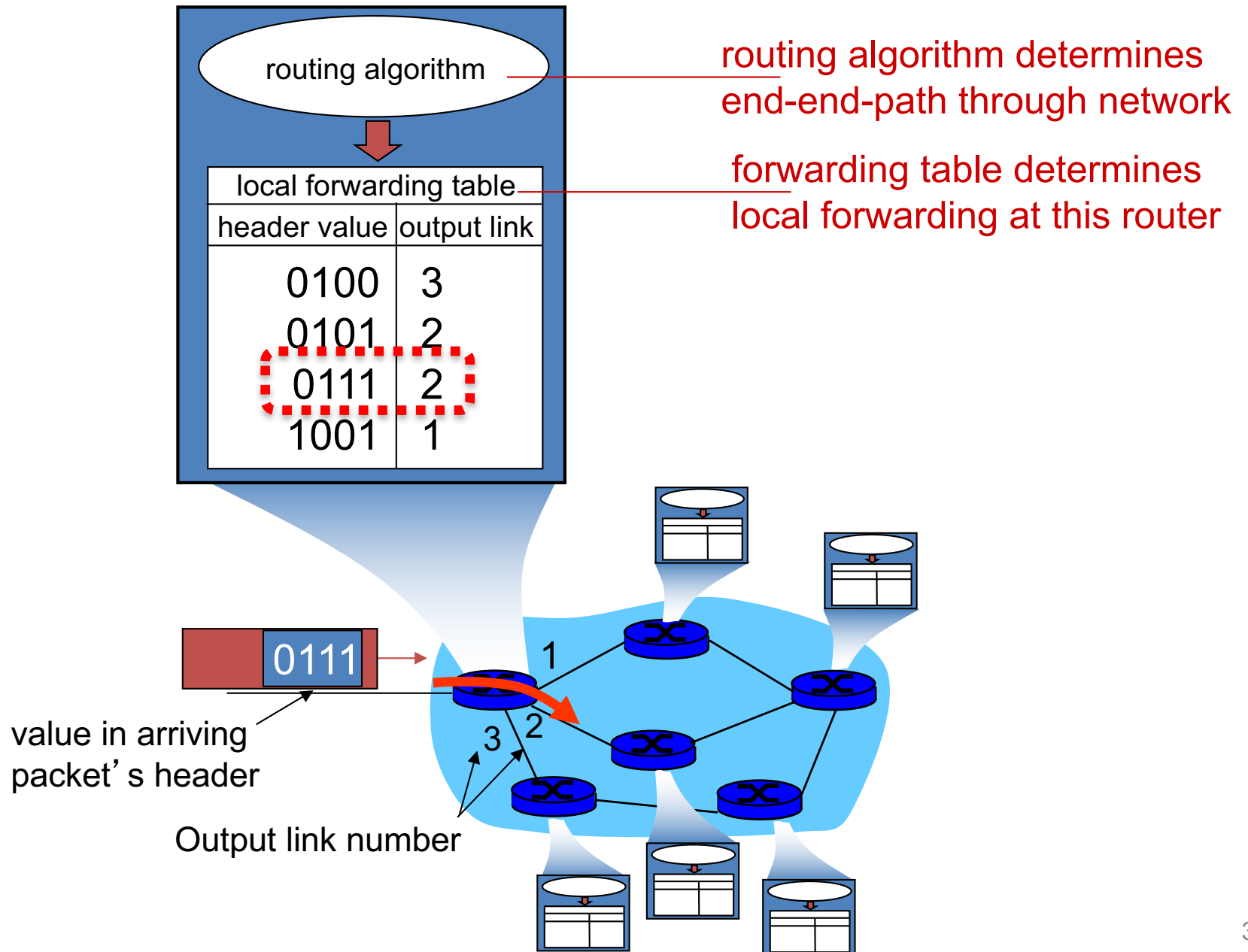
# Two key network-layer functions

- *forwarding*: move packets from router's input to appropriate router output
- *routing*: determine route taken by packets from source to dest.
  - *routing algorithms*

## *analogy:*

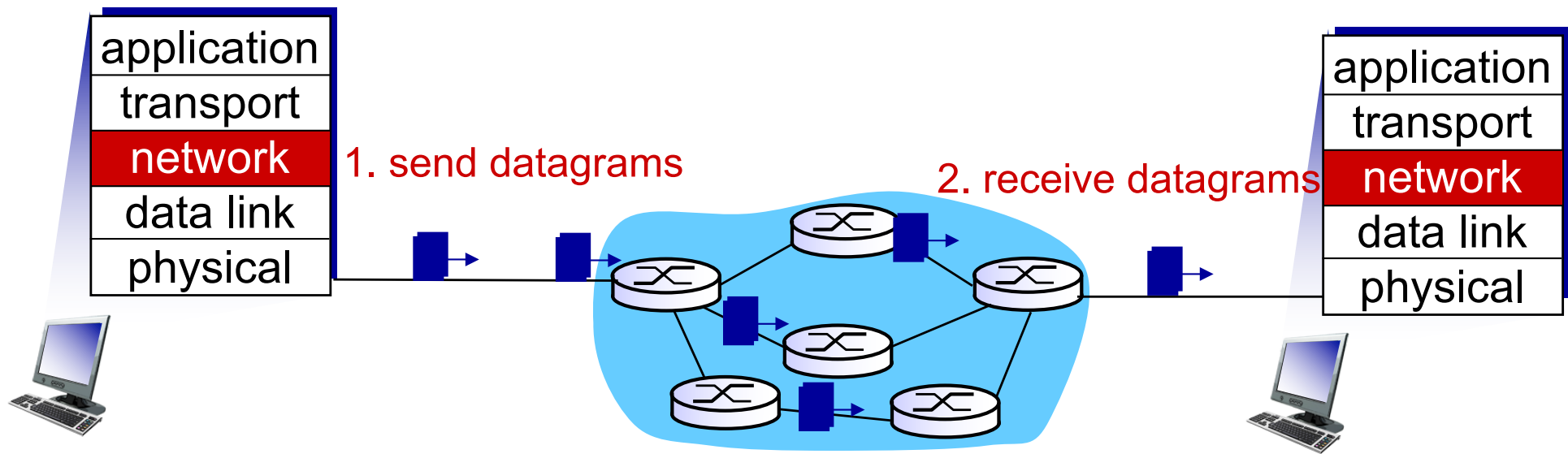
- ❖ *routing*: process of planning trip from source to dest
- ❖ *forwarding*: process of getting through single interchange

# Interplay between routing and forwarding

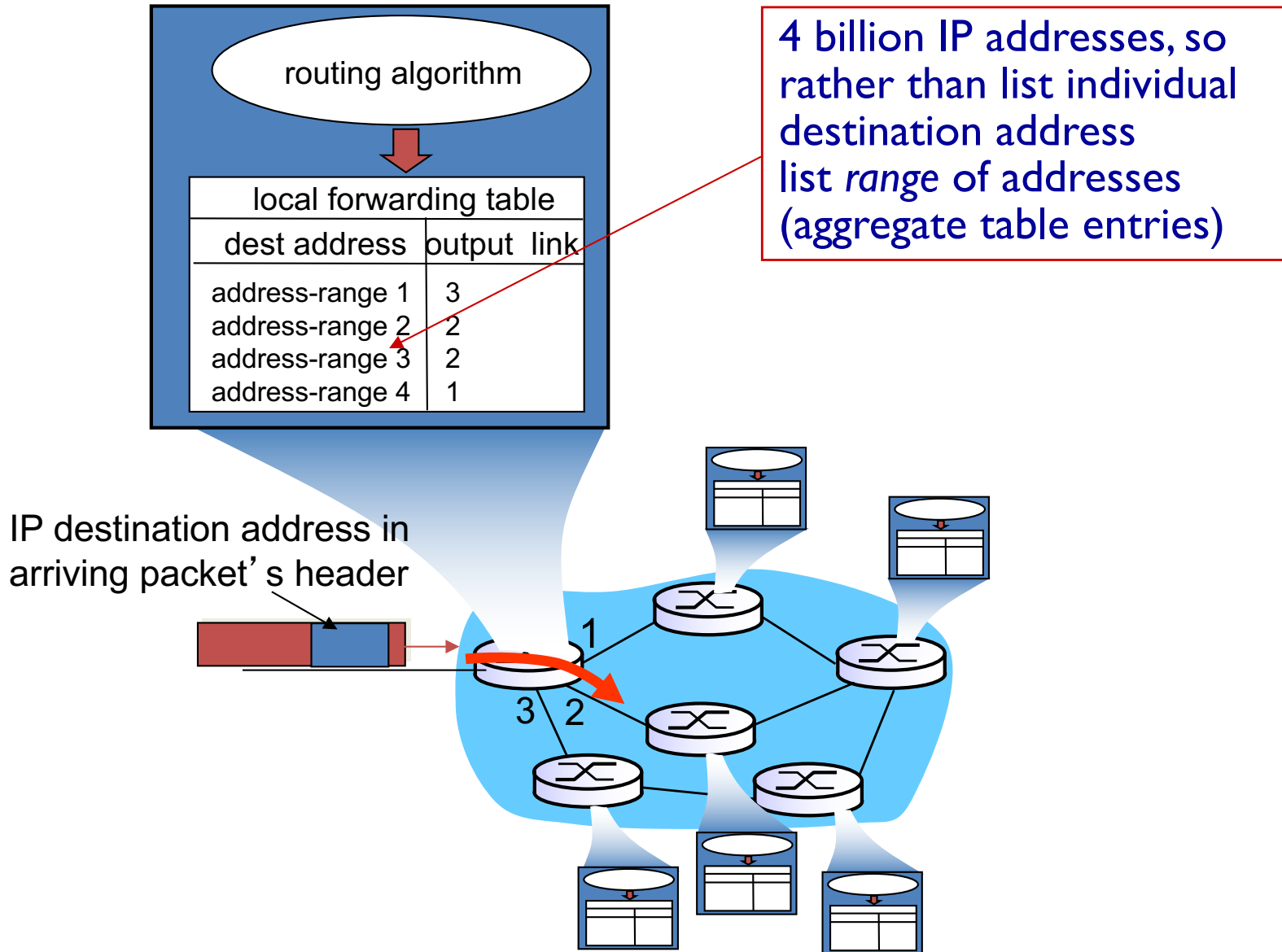


# Datagram networks

- no call setup at network layer
- routers: no state about end-to-end connections
  - no network-level concept of “connection”
- packets forwarded using destination host address



# Datagram forwarding table



# Datagram forwarding table

Destination Address Range	Link Interface
11001000 00010111 00010000 00000000 through 11001000 00010111 00010111 11111111	0
11001000 00010111 00011000 00000000 through 11001000 00010111 00011000 11111111	1
11001000 00010111 00011001 00000000 through 11001000 00010111 00011111 11111111	2
otherwise	3

# Longest prefix matching

## *longest prefix matching*

when looking for forwarding table entry for given destination address, use *longest* address prefix that matches destination address

Destination Address Range	Link interface
11001000 00010111 00010*** *****	0
11001000 00010111 00011000 *****	1
11001000 00010111 00011*** *****	2
otherwise	3

examples:

DA: 11001000 00010111 00010**110** 10100001

which interface?

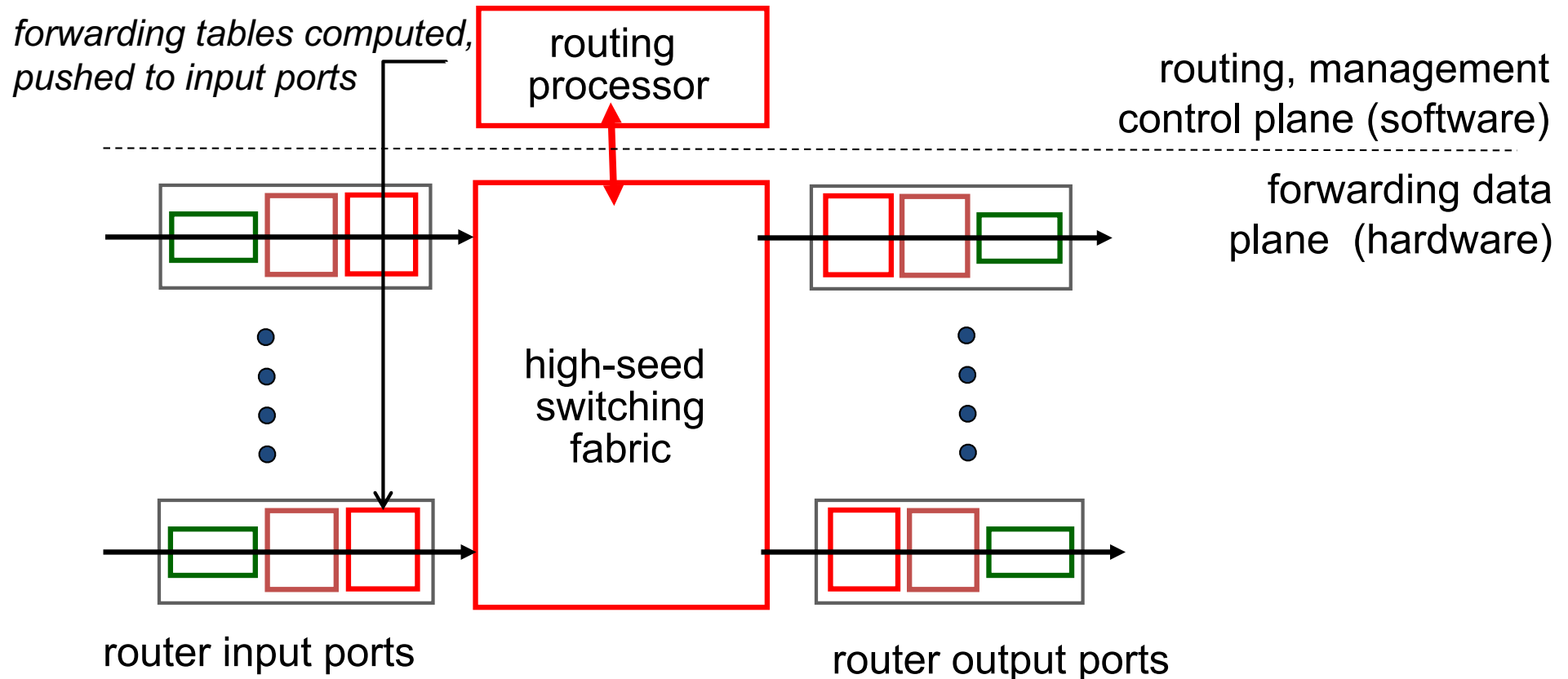
DA: 11001000 00010111 00011**000** 10101010

which interface?

# Router architecture overview

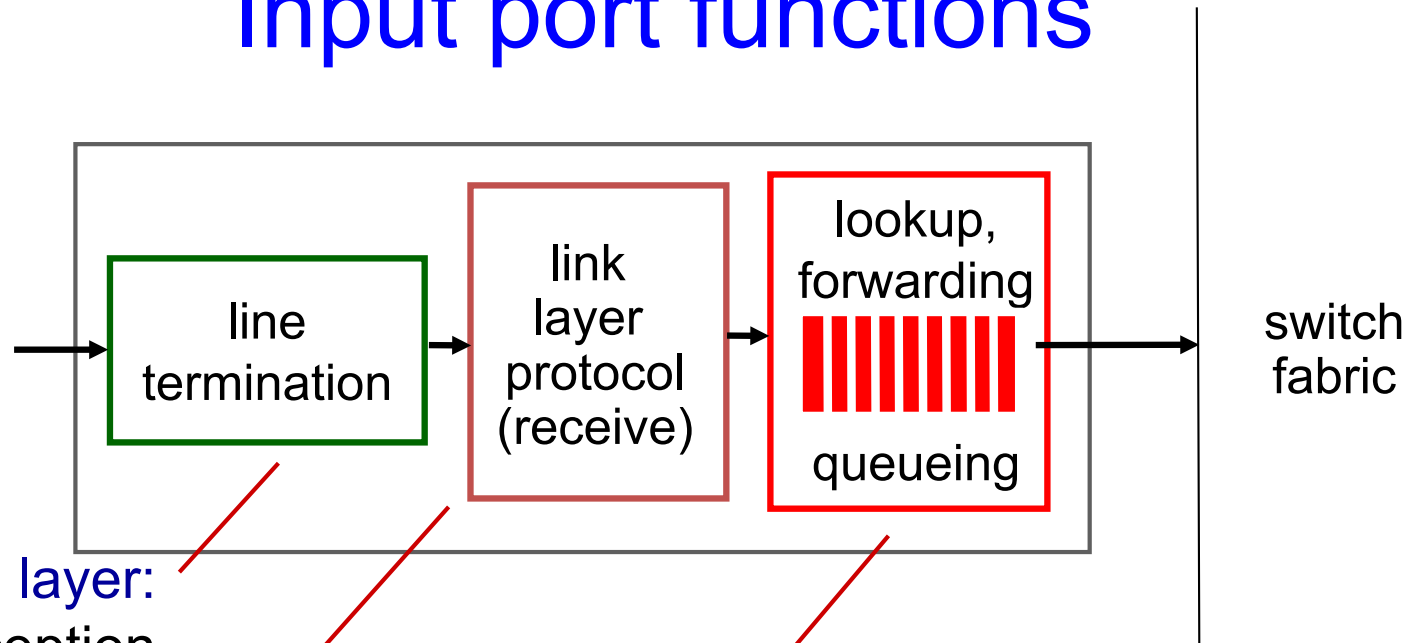
two key router functions:

- ❖ run routing algorithms/protocol (RIP, OSPF, BGP)
- ❖ *forwarding* datagrams from incoming to outgoing link





# Input port functions



physical layer:  
bit-level reception

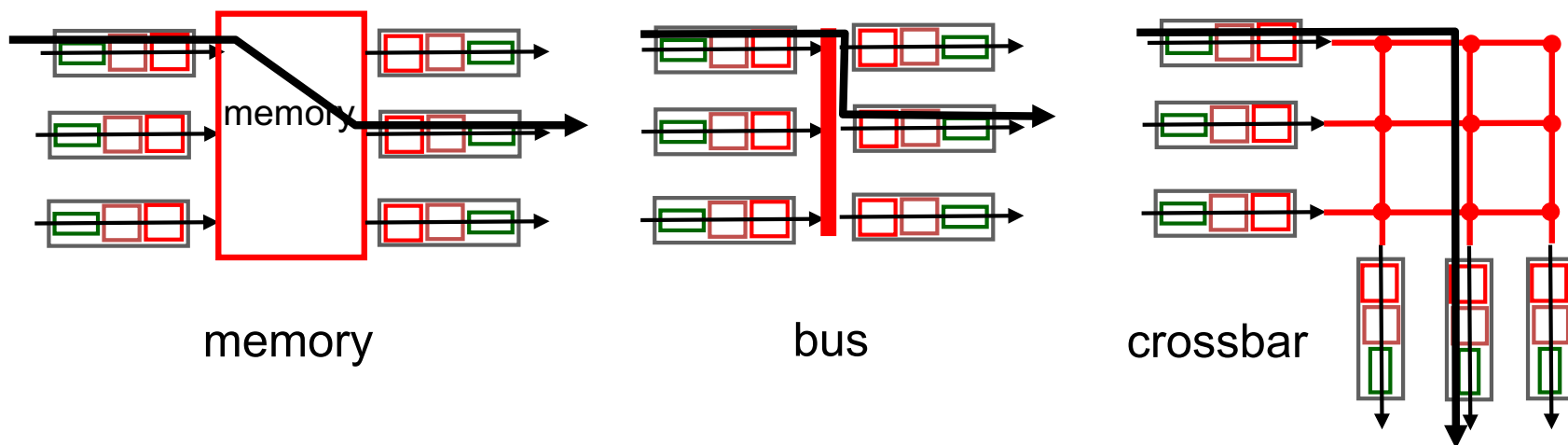
data link layer:  
e.g., Ethernet  
see chapter 5

## decentralized switching:

- given datagram dest., lookup output port using forwarding table in input port memory (*“match plus action”*)
- goal: complete input port processing at ‘line speed’
- queuing: if datagrams arrive faster than forwarding rate into switch fabric

# Switching fabrics

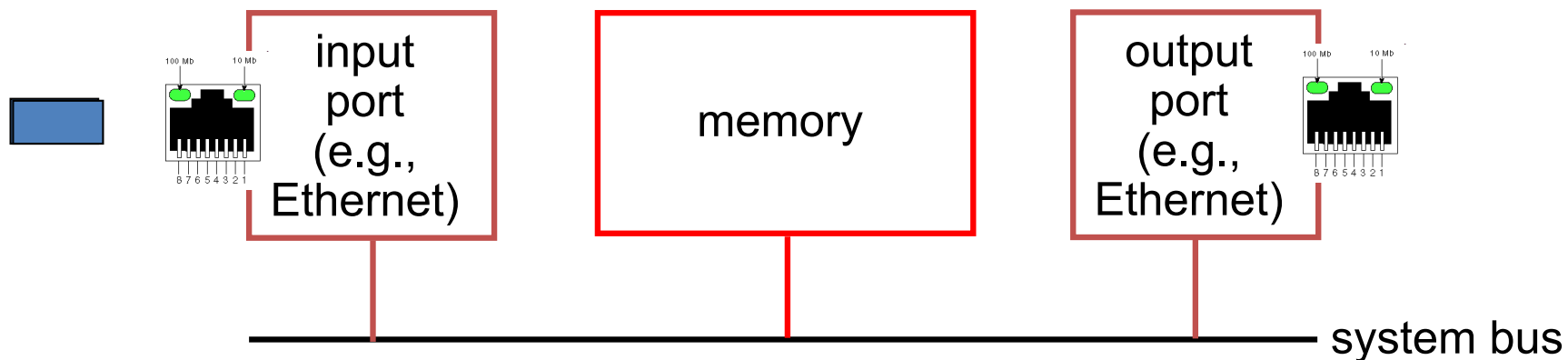
- transfer packet from input buffer to appropriate output buffer
- switching rate: rate at which packets can be transfer from inputs to outputs
  - often measured as multiple of input/output line rate
  - N inputs: switching rate N times line rate desirable
- three types of switching fabrics



# Switching via memory

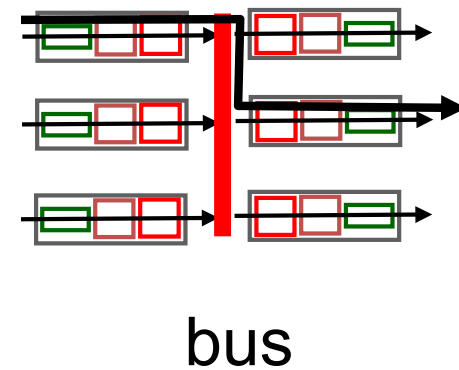
## *first generation routers:*

- traditional computers with switching under direct control of CPU
- packet copied to system's memory
- speed limited by memory bandwidth (2 bus crossings per datagram)



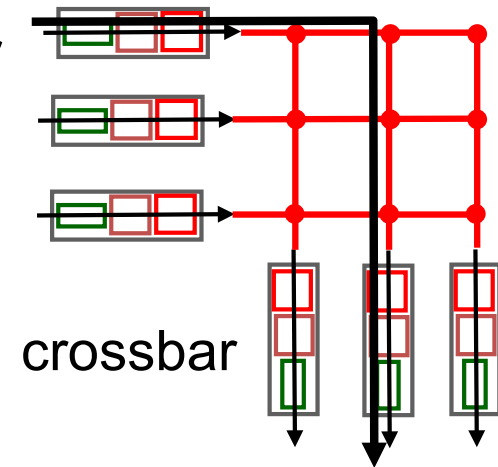
# Switching via a bus

- ❖ datagram from input port memory to output port memory via a shared bus
- ❖ *bus contention*: switching speed limited by bus bandwidth
- ❖ 32 Gbps bus, Cisco 5600: sufficient speed for access and enterprise routers

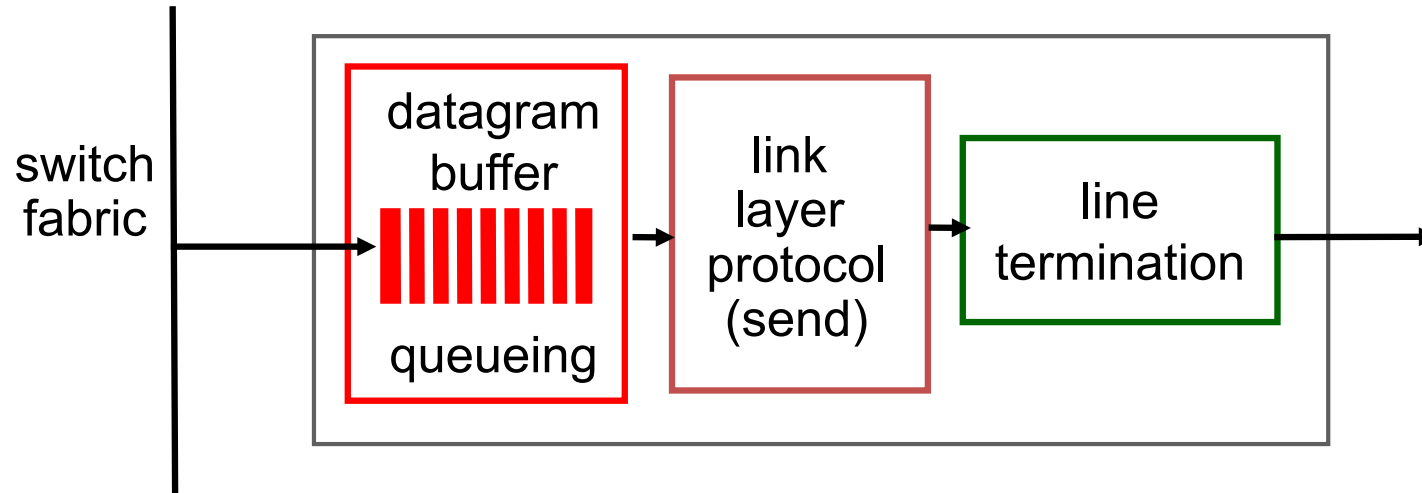


# Switching via interconnection network

- ❖ overcome bus bandwidth limitations
- ❖ banyan networks, crossbar, other interconnection nets initially developed to connect processors in multiprocessor
- ❖ advanced design: fragmenting datagram into fixed length cells, switch cells through the fabric
- ❖ Cisco 12000: switches 60 Gbps through the interconnection network

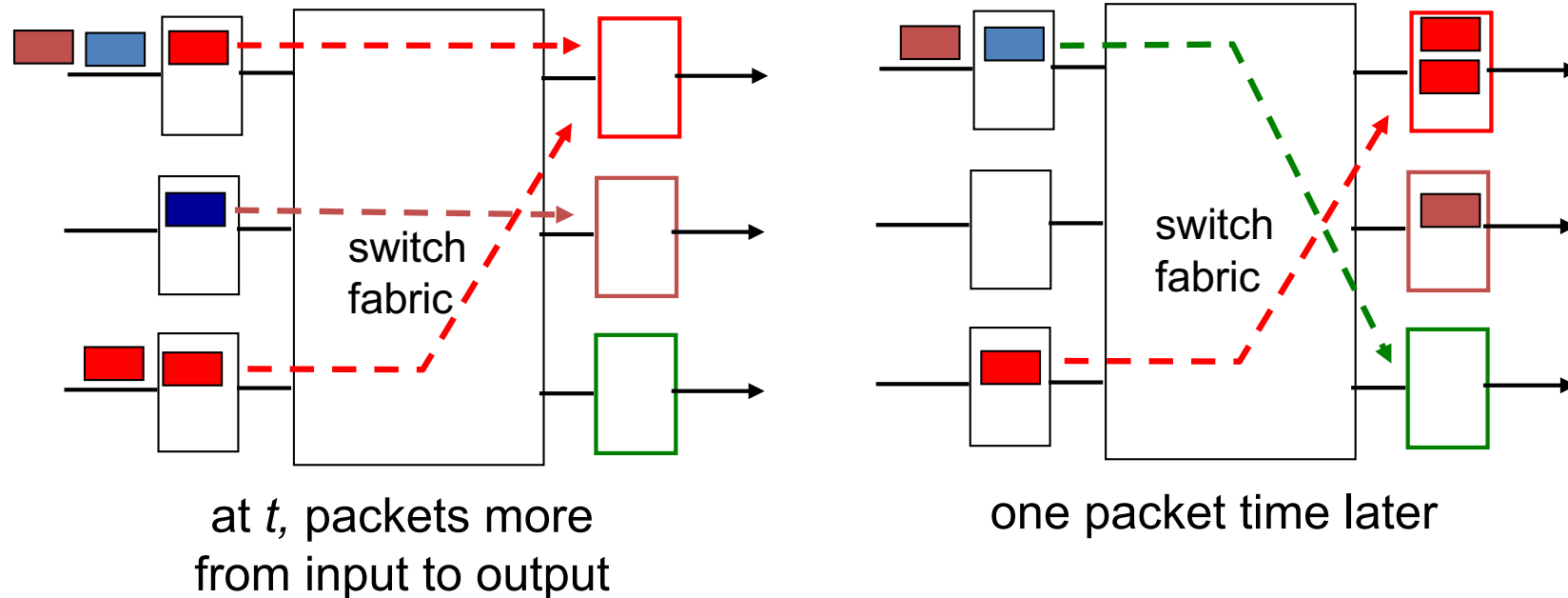


# Output ports



- ❖ *buffering* required when datagrams arrive from fabric faster than the transmission rate
- ❖ *scheduling discipline* chooses among queued datagrams for transmission
  - ❖ FIFO, Weighted Fair Queueing (WFQ)

# Output port queueing



- ❖ buffering when arrival rate via switch exceeds output line speed
- ❖ *queueing (delay) and loss due to output port buffer overflow!*

# How much buffering?

- RFC 3439 rule of thumb: average buffering should be equal to “typical” RTT (say 250 msec) times link capacity  $C$

– e.g.,  $C = 10$  Gpbs link: 2.5 Gbit buffer!!!

- recent recommendation: with  $N$  flows, buffering equal to

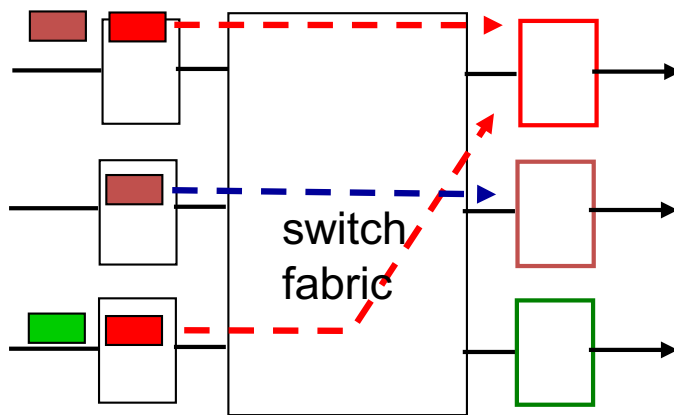
$$\frac{RTT \cdot C}{\sqrt{N}}$$

– e.g.,  $C = 10$  Gpbs link and  $N = 10,000$ : 25 Mbit buffer 😊

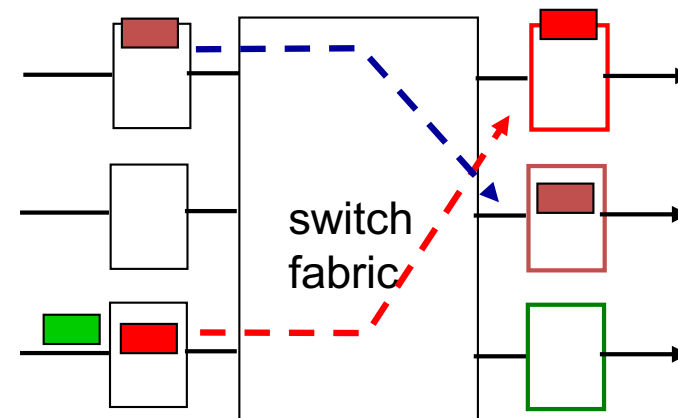


# Input port queuing

- fabric slower than input ports combined -> queueing may occur at input queues
  - *queueing delay and loss due to input buffer overflow!*
- **Head-of-the-Line (HOL) blocking:** queued datagram at front of queue prevents others in queue from moving forward



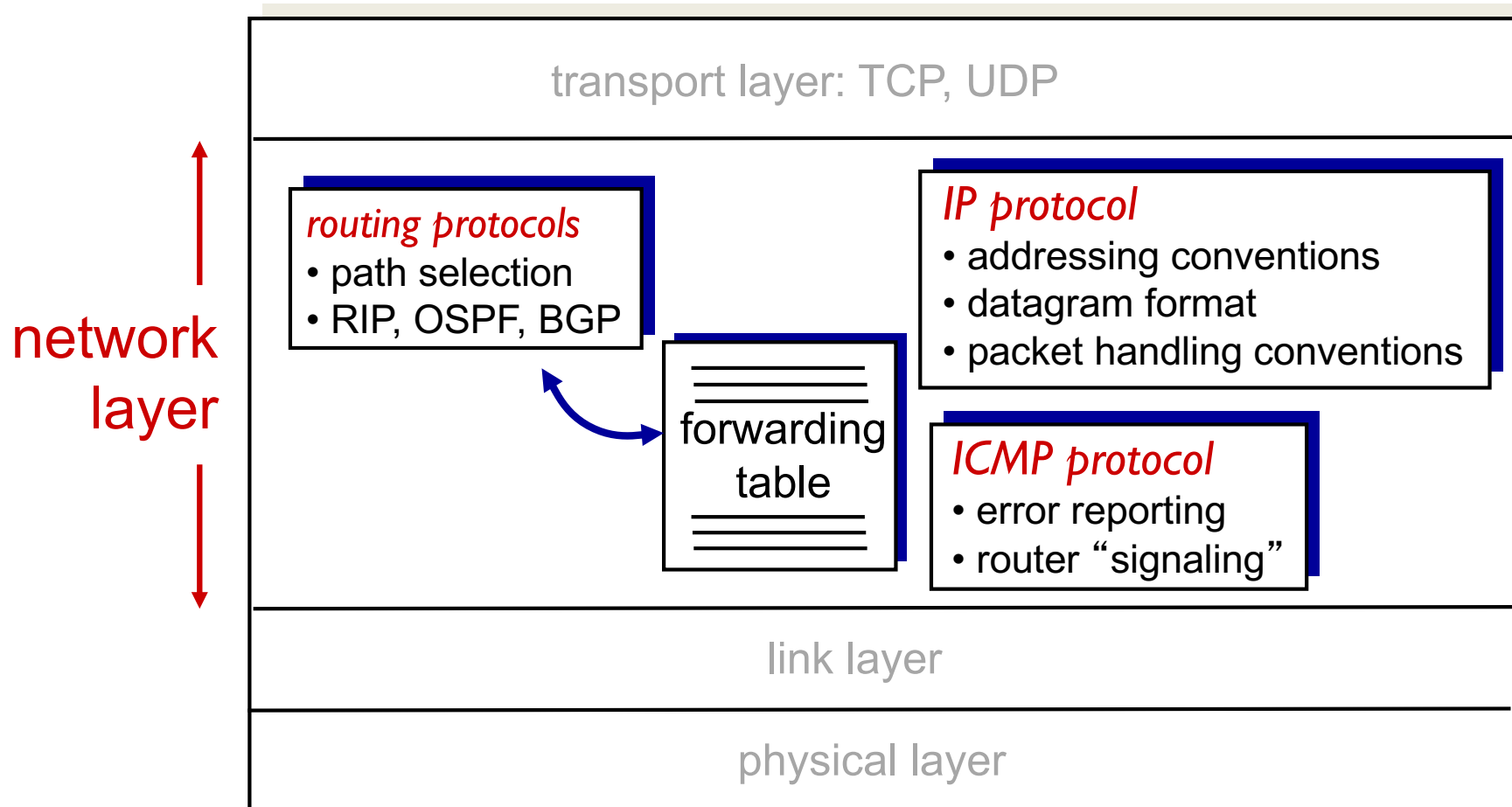
output port contention:  
only one red datagram can be  
transferred.  
*lower red packet is blocked*



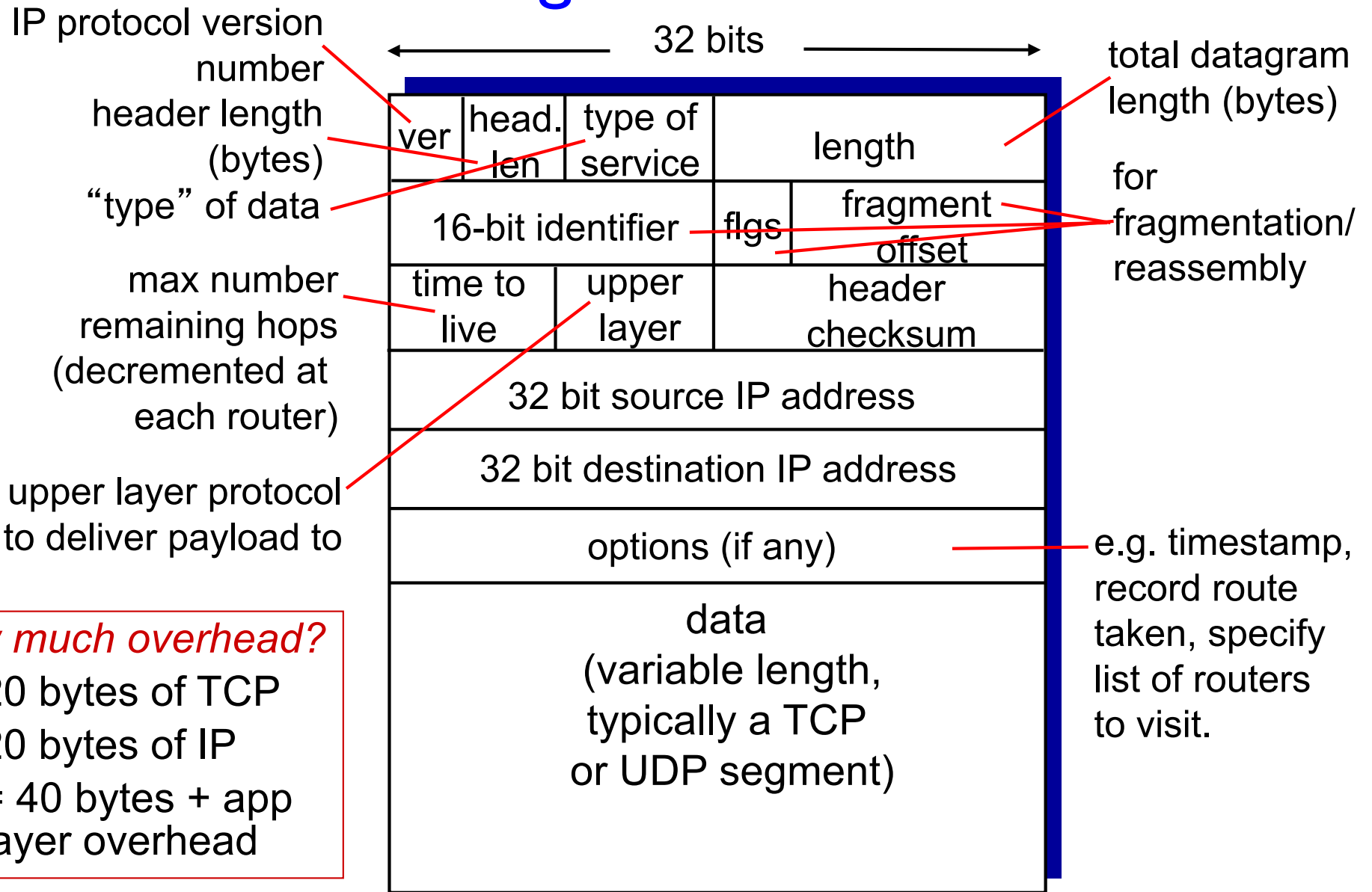
one packet time later:  
green packet  
experiences HOL  
blocking

# The Internet network layer

host, router network layer functions:

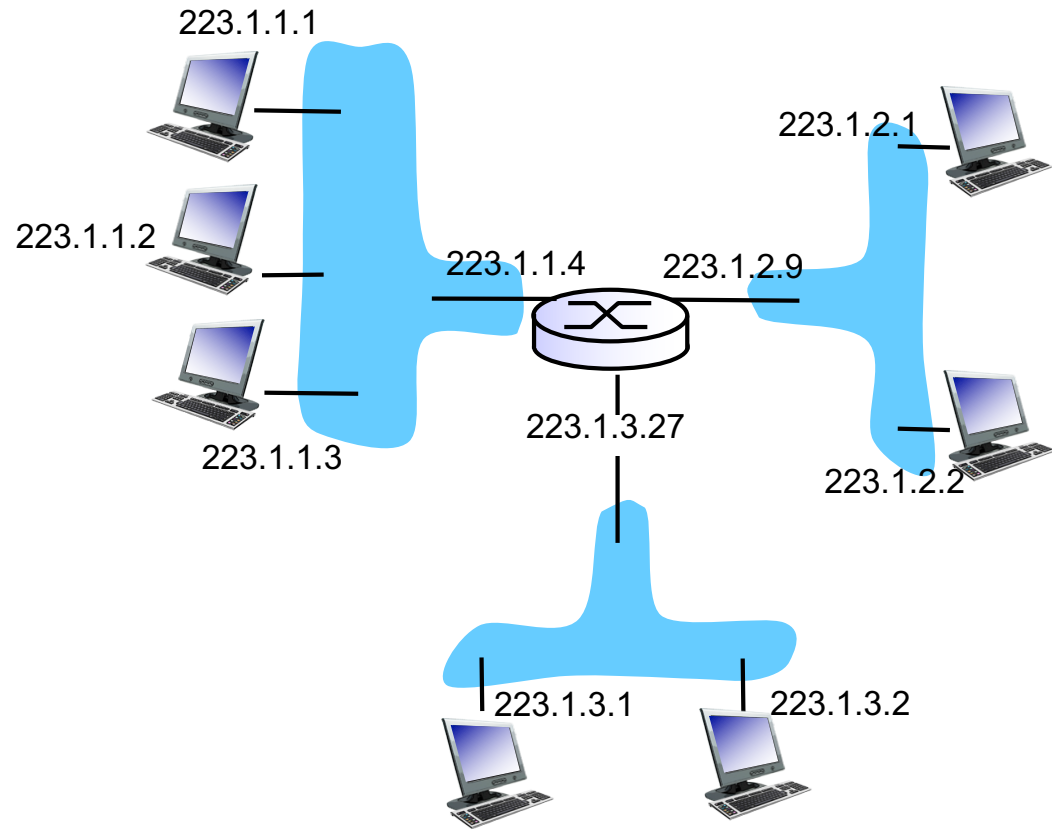


# IP datagram format



# IP addressing: introduction

- **IP address:** 32-bit identifier for host, router interface
- **interface:** connection between host/router and physical link
  - routers typically have multiple interfaces
  - host typically has one or two interfaces (e.g., wired Ethernet, wireless 802.11)
- **IP addresses associated with each interface**



$$223.1.1.1 = \underbrace{11011111}_{223} \underbrace{00000001}_{1} \underbrace{00000001}_{1} \underbrace{00000001}_{1}$$

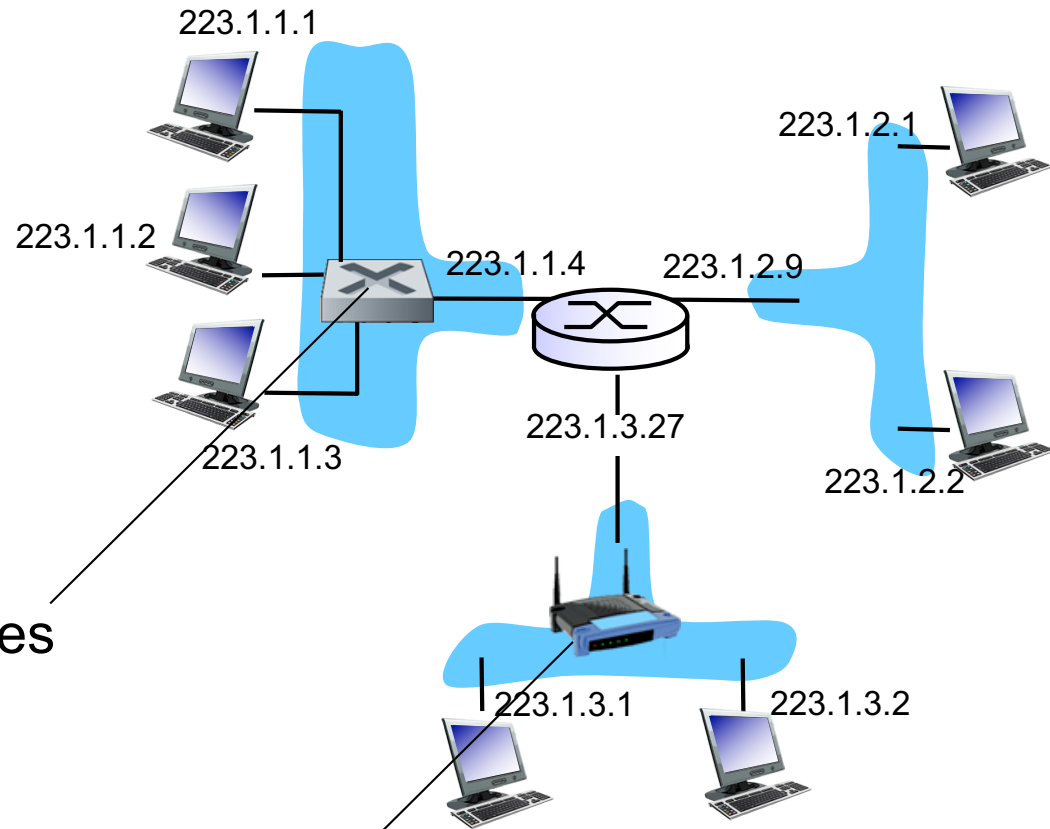
# IP addressing: introduction

**Q:** how are interfaces actually connected?

**A:** we'll learn about that in the link layer chapter

**A:** wired Ethernet interfaces connected by Ethernet switches

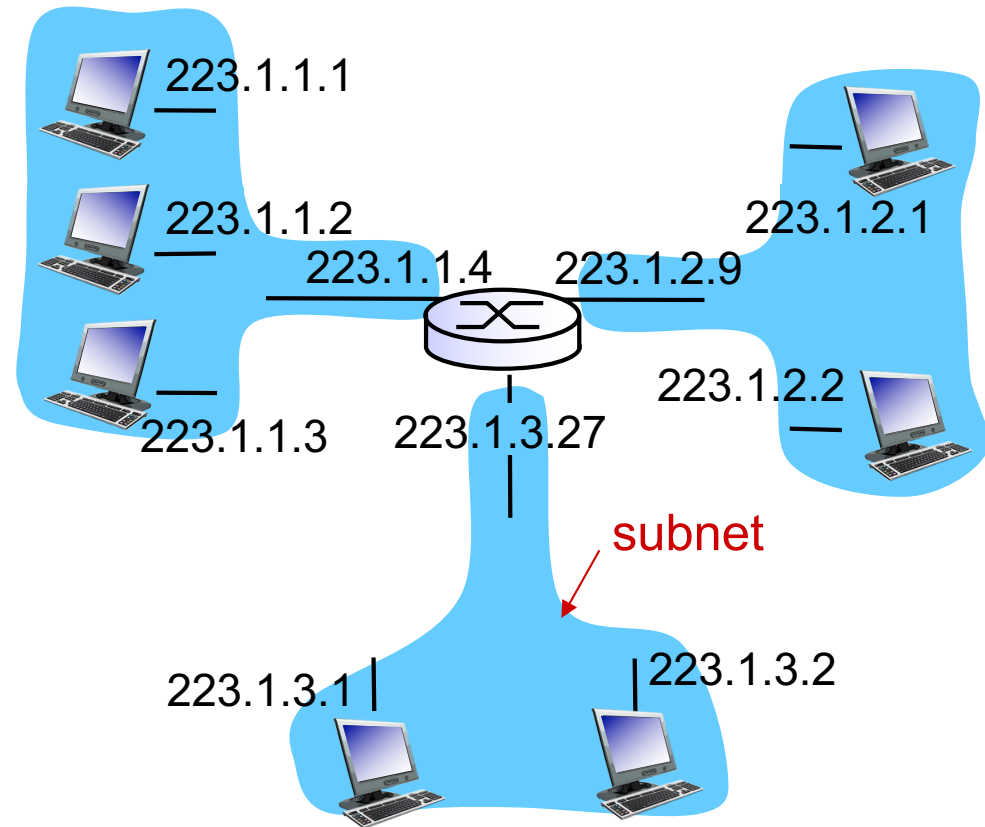
**For now:** don't need to worry about how one interface is connected to another (with no intervening router)



**A:** wireless WiFi interfaces connected by WiFi base station

# Subnets

- IP address:
  - subnet part - high order bits
  - host part - low order bits
- *what's a subnet?*
  - device interfaces with same subnet part of IP address
  - can physically reach each other *without intervening router*

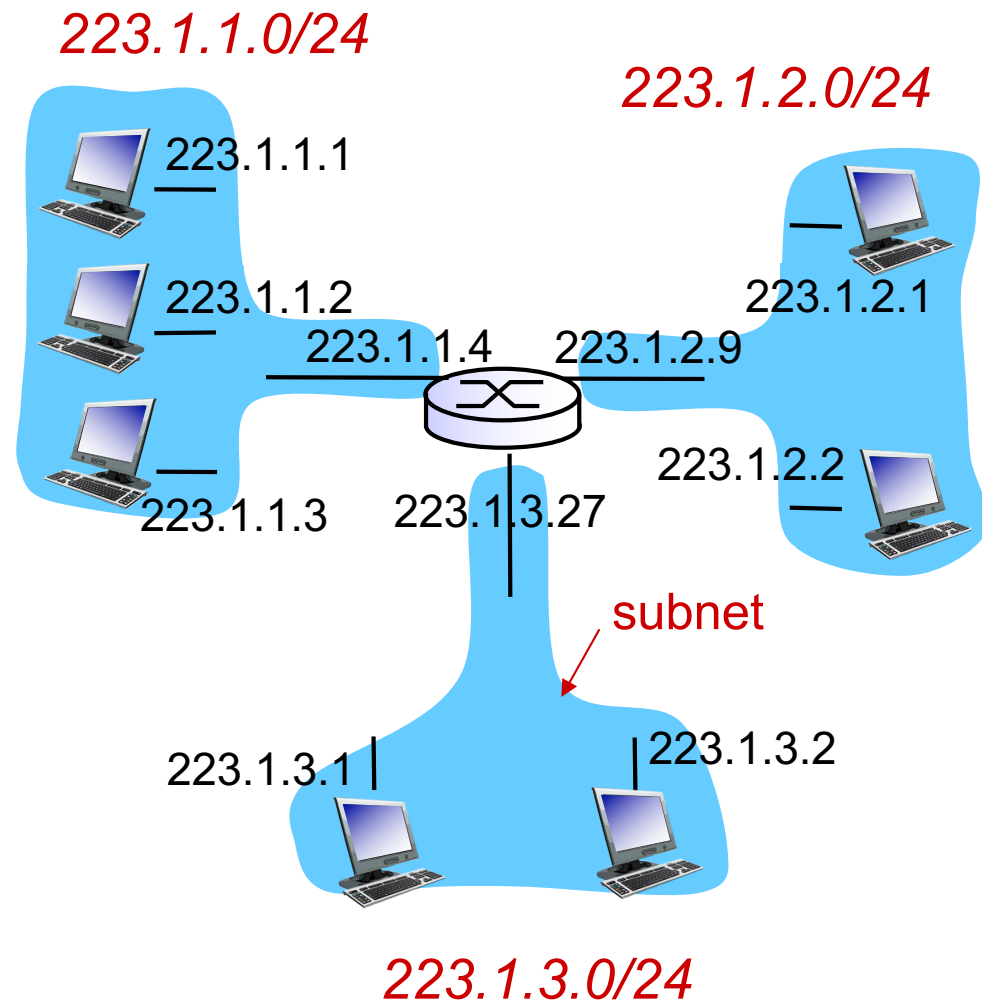


network consisting of 3 subnets

# Subnets

## *recipe*

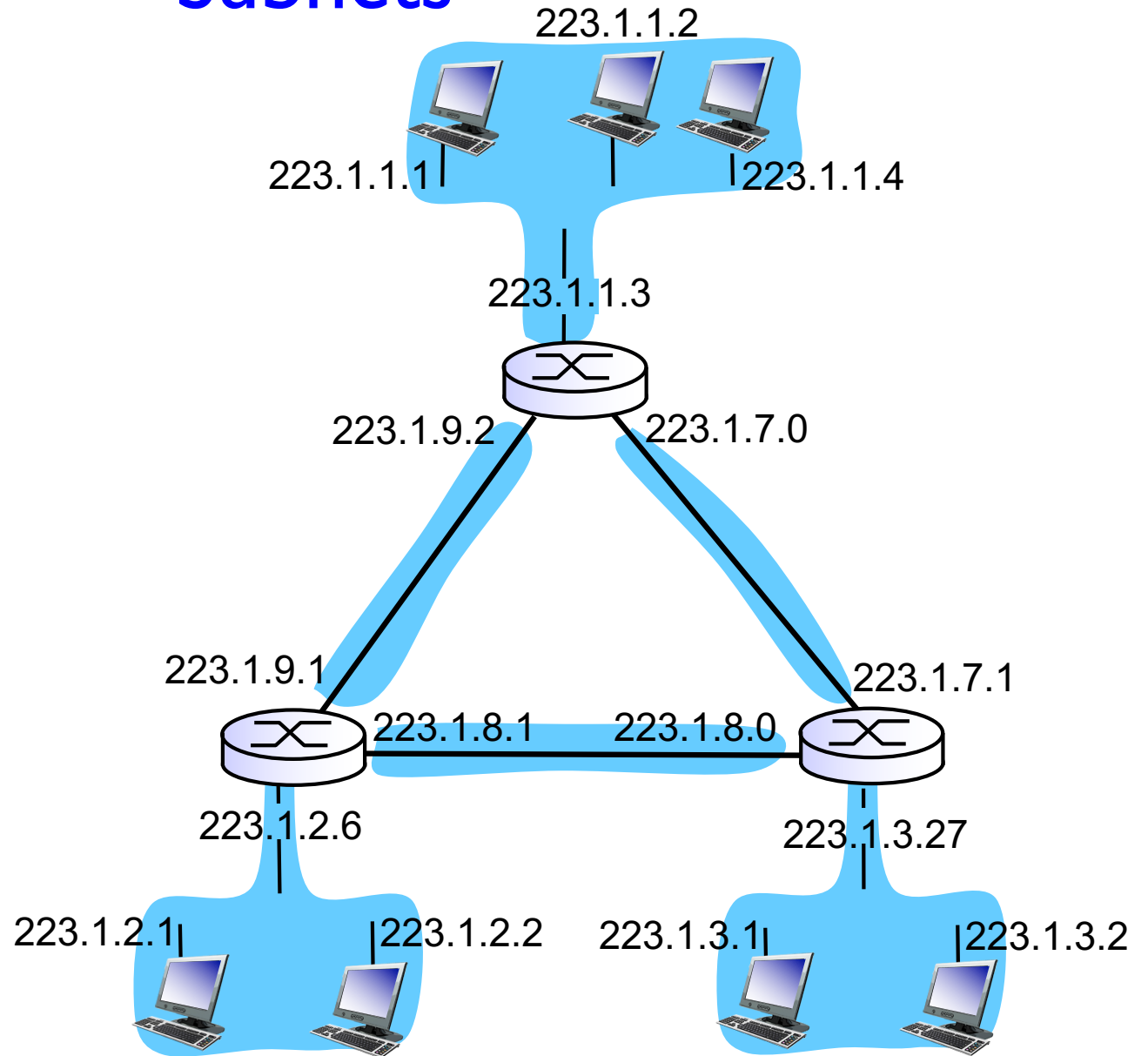
- to determine the subnets, detach each interface from its host or router, creating islands of isolated networks
- each isolated network is called a *subnet*



subnet mask: /24

# Subnets

how many?





# IP addressing: CIDR

## CIDR: Classless InterDomain Routing

- subnet portion of address of arbitrary length
- address format: **a.b.c.d/x**, where x is # bits in subnet portion of address

