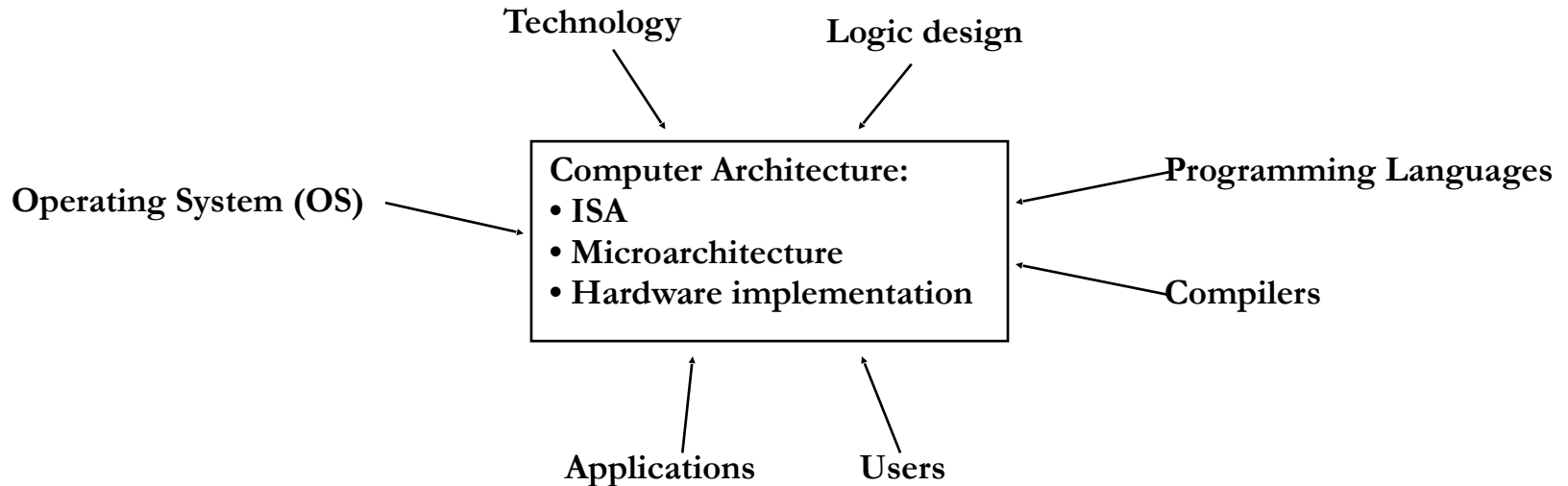


Summary: Computer Architecture

What is Computer Architecture?



Metrics of interest for computer architects

- Performance
- Cost
- Reliability
- Power
- Area



Performance of Computer Architectures

- Metrics:
 - Execution time, response time: overall time for a given computation (e.g. full program execution, one transaction)
 - Latency: time to complete a given task (e.g. memory latency, I/O latency, instruction latency)
 - Bandwidth, throughput: rate of completion of tasks (e.g. memory bandwidth, transactions per second, MIPS, FLOPS)
- MHz, MIPS, FLOPS: must use with caution
- Benchmarking:
 - toy benchmarks, synthetic benchmarks, kernels, real programs,
 - Input sets
 - E.g. SPECint, SPECfp, EEMBC Coremark



CPU Performance Terminology

A is n times faster than B means

$$\frac{\text{Execution time of B}}{\text{Execution time of A}} = n$$

A is m% faster than B means

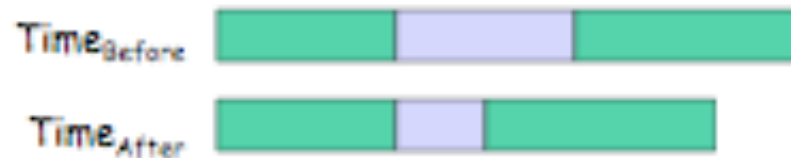
$$\frac{\text{Execution time of B} - \text{Execution time of A}}{\text{Execution time of A}} \times 100 = m$$

Improving Performance: Principles of CA Design

- Take advantage of parallelism
 - System level: multiple processors, multiple disks
 - Processor level: pipelining, superscalar issue
 - Circuit level: carry-lookahead ALU
- Principle of locality
 - Spatial and Temporal Locality
 - 90% of program executing in 10% of code
 - E.g. Caches
- Focus on the common case
 - Amdahl's law, CPU Performance equation
 - E.g. RISC design principle

Amdahl's Law

Speedup due to an enhancement E:



$$\text{Speedup (E)} = \frac{\text{Execution time w/o E (Before)}}{\text{Execution time w E (After)}}$$

Suppose that enhancement E accelerates a fraction F of the task by a factor S, and the remainder of the task is unaffected, what is the *Execution time_{after}* and *Speedup(E)*?

Amdahl's Law

$$\textit{Execution time}_{\textit{after}} = \textit{ExTime}_{\textit{before}} \times \left[(1-F) + \frac{F}{S} \right]$$

$$\textit{Speedup}(E) = \frac{\textit{ExTime}_{\textit{before}}}{\textit{ExTime}_{\textit{after}}} = \frac{1}{\left[(1-F) + \frac{F}{S} \right]}$$

Amdahl's Law - Example

Q: Floating point instructions improved to run 2X; but only 10% of execution time are FP ops. What is the execution time and speedup after improvement?

Ans:

$$F = 0.1, S = 2$$

$$\text{ExTime}_{\text{after}} = \text{ExTime}_{\text{before}} \times [(1-0.1) + 0.1/2] = 0.95 \text{ ExTime}_{\text{before}}$$

$$\text{Speedup} = \frac{\text{ExTime}_{\text{before}}}{\text{ExTime}_{\text{after}}} = \frac{1}{0.95} = 1.053$$

Factors that affect CPU performance

Instruction count (IC)

- Compiler, ISA

Cycles per instruction (CPI)

- ISA, microarchitecture

Clock time ($1/f$)

- Microarchitecture, technology

**What
determines
these
factors?**

The CPU Performance Equation

$$\text{CPU time} = \text{IC} \times \text{CPI} \times \text{Clock time}$$

where: CPU time = execution time
IC = number of instructions executed (instruction count)
CPI = number of average clock cycles per instruction
Clock time = duration of processor clock

$$\text{CPU time} = \left(\sum_{i=1}^n \text{IC}_i \times \text{CPI}_i \right) \times \text{Clock time}$$

where: IC_i = IC for instruction (instruction group) i
 CPI_i = CPI for instruction (instruction group) i

$$\text{CPI} = \frac{\left(\sum_{i=1}^n \text{IC}_i \times \text{CPI}_i \right)}{\text{IC}} = \sum_{i=1}^n \left(\text{CPI}_i \times \frac{\text{IC}_i}{\text{IC}} \right)$$

Examples

- Example:**
- Branch instructions take 2 cycles, all other instructions take 1 cycle
 - CPU A uses extra compare instruction per branch
 - Clock frequency of CPU A is 1.25 times faster than CPU B
 - On A 20% of instructions are branches (thus other 20% are compare instructions)

Find the CPU time

$$CPI_A = CPI_{\text{branch}} \times \frac{IC_{\text{branch}}}{IC} + CPI_{\text{others}} \times \frac{IC_{\text{others}}}{IC} = 2 \times 0.2 + 1 \times 0.8 = 1.2$$

$$CPI_B = CPI_{\text{branch}} \times \frac{IC_{\text{branch}}}{IC} + CPI_{\text{others}} \times \frac{IC_{\text{others}}}{IC} = 2 \times 0.25 + 1 \times 0.75 = 1.25$$

$$\text{CPU time}_B = IC_B \times CPI_B \times \text{Clock time}_B = 0.8 \times IC_A \times 1.25 \times (1.25 \times \text{Clock time}_A)$$

$$= 1.25 \times IC_A \times \text{Clock time}_A$$

$$\text{CPU time}_A = IC_A \times CPI_A \times \text{Clock time}_A = IC_A \times 1.2 \times \text{Clock time}_A$$

CPU A is faster!

Improving CPU Performance

IC:

- Compiler optimizations (constant folding, constant propagation)
- ISA (More complex instructions)

CPI:

- Microarchitecture (Pipelining, Out-of-order execution, branch prediction)
- Compiler (Instruction scheduling)
- ISA (Simpler instructions)

Clock period:

- Hardware (Smaller transistors – Moore's law)
- ISA (Simple instructions that can be easily decoded)
- Microarchitecture (Simple architecture)