Advances in Programming Languages APL14: Polyphonic C#

Ian Stark

School of Informatics The University of Edinburgh

> Monday 3 March 2008 Semester 2 Week 9



This is the second of three lectures presenting some programming-language techniques for managing concurrency.

- Java, Erlang
- Polyphonic C#
- Cautionary Tales

This is the second of three lectures presenting some programming-language techniques for managing concurrency.

- Java, Erlang
- Polyphonic C#
- Cautionary Tales

Concurrent Programming

"It is a truth universally acknowledged that concurrent programming is significantly more difficult than sequential programming"

 $[C\omega | anguage overview]$

Concurrency is useful:

- Efficient use of mixed resources (disk, memory, network)
- Responsiveness (GUI, hardware interrupts, managing mixed resources)
- Speed (multiprocessing, hyperthreading, multicore)
- Multiple clients (database engine, web server)

Concurrency is hard:

- Interference (shared store, simultaneous modification)
- Liveness (deadlock, livelock, lack of progress)
- Fairness (scheduling, prioritization, starvation)
- Safety (correctness, error handling, specification)

Java

... has language facilities for spawning a new java.lang.Thread, with **synchronized** code blocks using per-object locks to protect shared data, and signalling between threads with wait and notify.

C#

... has language facilities for spawning a new System.Threading.Thread, to **lock** code blocks using per-object locks to protect shared data, and signalling between threads with Wait and Pulse methods.

There are many ways to program concurrency, such as the following from functional languages:

Erlang: Multiple share-nothing threads, each with a single mailbox for asynchronous communication by message-passing.

Concurrent ML: Multiple threads, multiple channels for synchronous message-passing communication between them.

Concurrent Haskell: Multiple threads, asynchronous communication through MVar mutable variables.

There are also mathematical models to capture and analyse the essence of these concurrent systems: such as the π -calculus, the join-calculus, and the ambient calculus.

Looking for ways to improve concurrent programming in object-oriented languages, consider the following themes:

- Focus on communication rather than concurrency.
- Unify message passing with method invocation.
- Look not just for individual messages but patterns of messages.

Polyphonic C# is a mild extension of C# which provides novel primitives for writing concurrent programs, based on the join calculus.

"The language presents a simple and powerful model of concurrency which is applicable both to multithreaded applications running on a single machine and to the orchestration of asynchronous, event-based applications communicating over a wide area network."

[Benton, Cardelli, Fournet]

These extensions also appear in $C\omega$, the research programming language we met earlier as the source of LINQ.

Polyphony, n. 1. a. Music. Harmony; esp. the simultaneous and harmonious combination of a number of individual melodic lines. Oxford English Dictionary, draft revision, June 2007

Asynchronous methods

Conventional method invocation in C# is *synchronous*: when code calls a method on an object, it cannot continue until that method completes.

In contrast, when code invokes an *asynchronous* method, it continues at once, and does not have to wait for the method to finish.

Chords

Standard method declarations associate one piece of code (the *body*) to each method name (up to overloading by parameter type and number).

In Polyphonic C#, a *chord* declares code that is to be executed only when a particular combination of methods are invoked.

Example: Storage Cell

The following C# code defines a straightforward storage cell, containing a single String value.

```
public class Cell {
    private String contents = "";
    public String get() { return contents; }
    public void put(String s) { contents=s; }
}
```

Likely to be thread safe, provided put and get methods remain this simple.

Example: Storage Cell

The following C# code defines a straightforward storage cell, containing a single String value.

```
public class Cell {
    private String contents = "";
    public String get() { return contents; }
    public async put(String s) { contents=s; }
}
```

The *asynchronous* put method now returns immediately to its caller, and may use a separate thread to update the cell contents.

public class Buffer {

public String get() & public async put(String s) { return s; }

This still has two methods, get and put, but now jointly defined in a *chord*, with a single return statement in the body.

Consumers call get(): this blocks until a producer invokes put(s), and then the chord is complete so s is returned to the consumer.

Producers call put(s): if a consumer is waiting on get(), then the chord is complete and value is handed on; if not, the call is noted, and control returns to the producer. Either way, the **async** call returns at once.

Multiple put or get calls can be outstanding at any time.

public class Buffer {

public String get() & public async put(String s) { return s; }

- No threads are spawned: the body of the chord is executed by the caller of the synchronous get method.
- Where there are multiple threads, it is entirely thread-safe: several producers and consumers can run simultaneously.
- No critical sections, monitors or mutual exclusion: there is no shared storage for interference.
- No explicit locks: the compiler looks after the brief locking required at the moment of chord selection.

public class Buffer {

public String get() & public async put(String s) { return s; }

- Each chord may combine many method names.
- At most one method in a chord can be synchronous.
- Each method can appear in multiple chords.
- A chord may be entirely asynchronous.
- Synchronous calls may block; asynchronous calls always return at once.
- Calls stack up until a chord is matched.

public class OnePlaceBuffer {

```
public OnePlaceBuffer() { empty(); }
```

```
public void put(String s) & private async empty() {
    contains(s);
    return;
}
```

public String get() & private async contains(String s) {
 empty();
 return s;
}

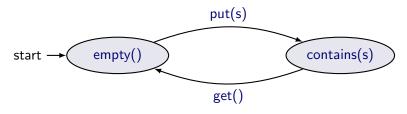
The class has four methods:

- Two public synchronous methods put(s) and get();
- Two private asynchronous methods empty() and contains(s).

There is always exactly one empty() or contains(s) call pending. No threads are needed, but where there is concurrency the code remains safe.

- Method put(s) blocks unless and until there is an empty() call.
- Method get() blocks unless and until there is a contains(s) call.

The code operates a simple state machine:



```
delegate async IntCallback(int value); // Declare function type
```

Client code can dispatch a request to a Service, do some work of its own, then rendezvous to pick up the result when ready.

Compare XMLHttpRequest from Javascript, used in AJAX for asynchronous communication with web services.

```
delegate async IntCallback(int value); // Declare function type
```

Several requests can be dispatched together, and a client might wait until all or any of them are completed.

The Service and client can be on different machines: asynchronous request and callback methods means that they distribute well. Other classic concurrency idioms have versions in Polyphonic C#:

- Combining shared and exclusive access to resources with multiple-readers / single-writer (just five chords).
- Locks, semaphores, condition variables (if you want them).
- Active objects, concurrent objects, Actors.
- Concurrent publish/subscribe, subject/observer pattern.
- Custom schedulers: thread pooling, worker threads.
- (insert your favourite concurrent programming problem)

Some of these are just to show that chords are as expressive as other paradigms: in actual use, the ideal is to raise the level of abstraction and avoid explicit concurrency management.

- Central notion of *asynchronous* computation, directly addressing application responsiveness.
- Concurrency is implicit: no explicit threads, locks, mailboxes, channels,...; although all these could be coded up.
- High-level description of the desired interaction profile through *chords*.
- Declarative presentation means that a compiler can transform and optimize as appropriate: to avoid thread spawning if not required; reusing existing threads; worker threads, thread pools.
- Transformation could in principle also include platform details: multiple processors, multicore, distributed client/server.

The same notions have been applied to other languages in *JoCaml* and *Join Java*, and are also available in the *Joins* library for C# and VB.NET.

- Java and C# use explicit shared-memory concurrency; with threads, locks, monitors and semaphores.
- Erlang has explicit processes but they share no store, instead communicating by mailboxes.
- This can be generalised (Concurrent ML, Concurrent Haskell, ...) to multiple threads and multiple named communication channels.
- Polyphonic C# / Cω provide asynchronous methods and chords to orchestrate implicitly concurrent behaviour.

Homework

Read two pieces describing concurrent programming in Polyphonic C#:

- The three-page online introduction at http://research.microsoft.com/~nick/polyphony/intro.htm
- Nick Benton. Jingle Bells: Solving the Santa Claus Problem in Polyphonic C#. http://research.microsoft.com/~nick/polyphony/santa.pdf

If you are interested, find out more in the concurrency sections of the $C\omega$ web site, and in the following paper:

James Gosling, Bill Joy, Guy Steele.
 The Java Language Specification (First Edition), Chapter 17.
 Addison-Wesley, 1996.
 Modern Concurrency Abstractions for C#.
 ACM Transactions on Programming Languages and Systems 26(5):269–804. September 2004. DOI: 10.1145/1018203.1018205