

Controlling Backtracking: The Cut

Artificial Intelligence Programming in Prolog

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Lecture 7

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Clearing up equality

- There are various ways to test equality in Prolog.

| | |
|-----------------------------|--|
| $X = Y$ | succeeds if the terms X and Y unify . |
| $X \text{ is } Y$ | succeeds if the arithmetic value of expression Y matches the value of term X. |
| $X ::= Y$ | succeeds if the arithmetic value of two expressions X and Y match. |
| $X \neq Y$ | succeeds if the arithmetic value of two expressions X and Y DO NOT match. |
| $X == Y$ | succeeds if the two terms have literal equality = are structurally identical and all their components have the same name. |
| $X \neq Y$ | succeeds if the two terms are NOT literally identical. |
| $\backslash + \text{ Goal}$ | succeeds if Goal does not true |

Clearing up equality (2)

| ?- 3+4 = 4+3.
no % treats them as terms

| ?- 3+4 = 3+4.
yes

| ?- X = 4+3.
X = 4+3 ?
yes

| ?- X is 4+3.
X = 7 ?
yes

| ?- 3+4 is 4+3.
no % left arg. has to be a term

| ?- 3+4 ::= 4+3.
yes % calculates both values

| ?- 3+4 =\= 4+3.
no

| ?- 3+4 == 4+3.
no

| ?- 3+4 \== 4+3.
yes

| ?- 3+X = 3+4.
X = 4 ? yes

| ?- 3+X == 3+4.
no

| ?- \+ 3+4 == 4+3.
yes

Processing in Prolog

To `call` the goal `G`:

1. Find first clause head that matches `G`:
 1. bind all variables accordingly,
 2. `call` goals in body in order;
 3. if all succeed, `G` succeeds (and `exits`).
2. else try next clause down;
3. if no next clause, `fail` the goal `G`.

When a goal fails:

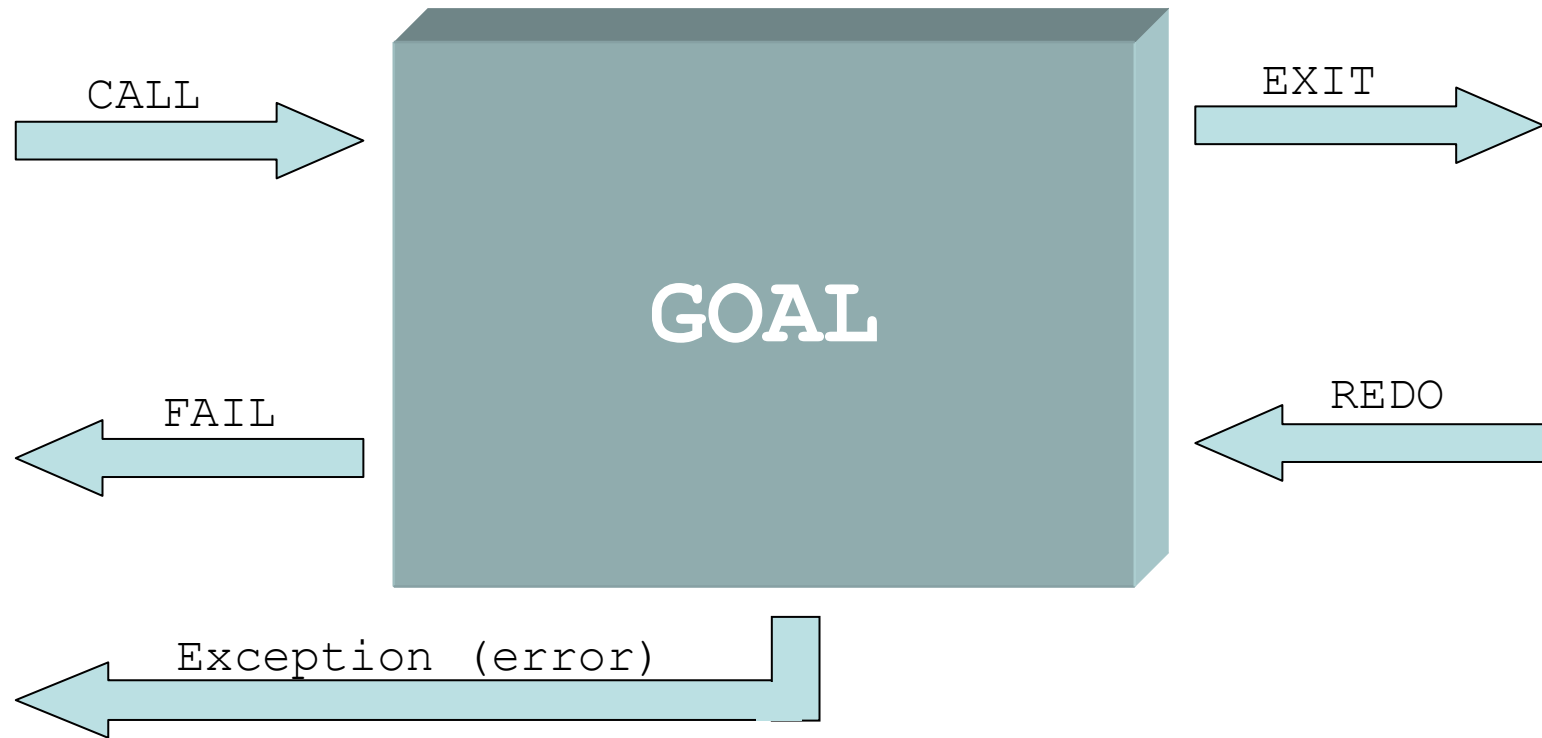
`redo` the most recent successful goal

To `redo` a goal:

1. discard bindings from previous success;
2. try clauses for this goal not so far tried;
3. if none, fail the goal.

Byrd Box model

- This is the model of execution used by the tracer.
- Originally suggested by Lawrence Byrd.



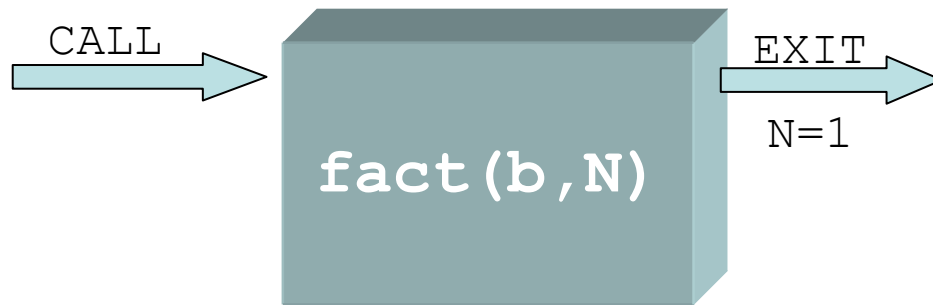
Redo-ing a Goal

`fact(b,1).`

`fact(b,2).`

`a :- fact(b,N), fact(c,N).`

`|?- a.`



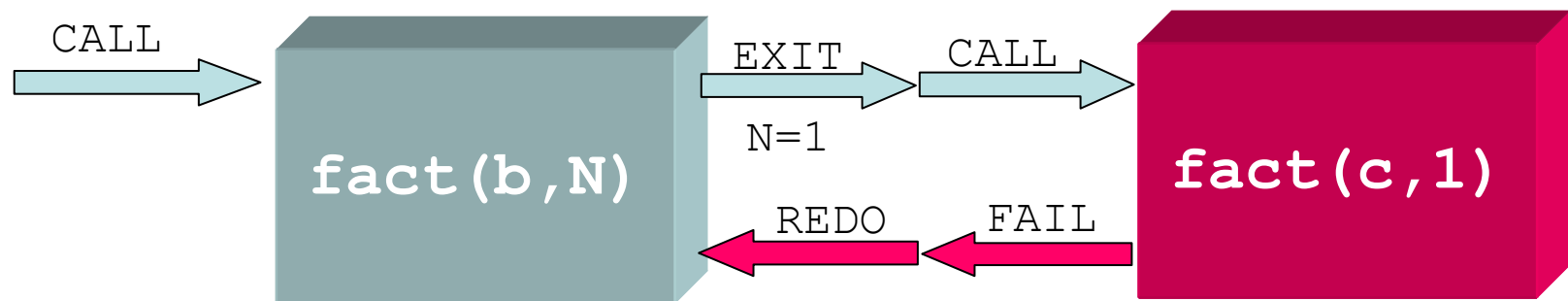
Redo-ing a Goal (2)

`fact(b,1).`

`fact(b,2).`

`a :- fact(b,N), fact(c,N).`

`|?- a.`



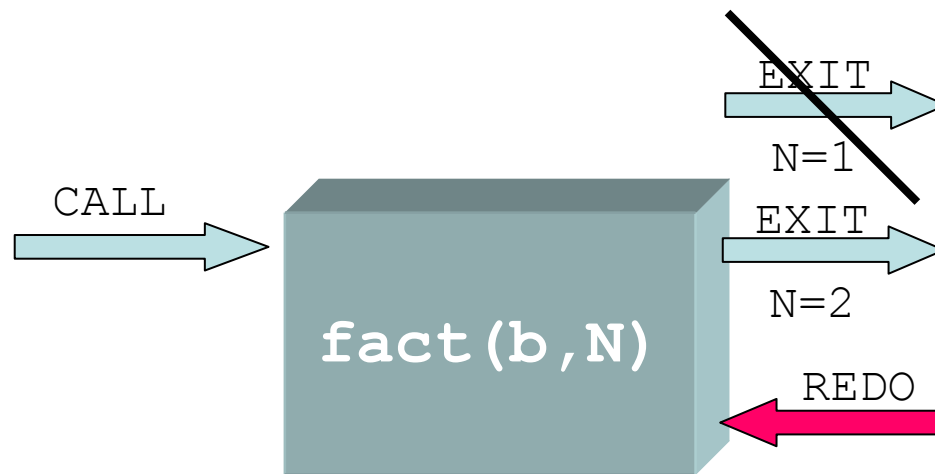
Redo-ing a Goal (3)

`fact(b,1).`

`fact(b,2).`

`a :- fact(b,N), fact(c,N).`

`!?- a.`



Redo-ing a Goal (4)

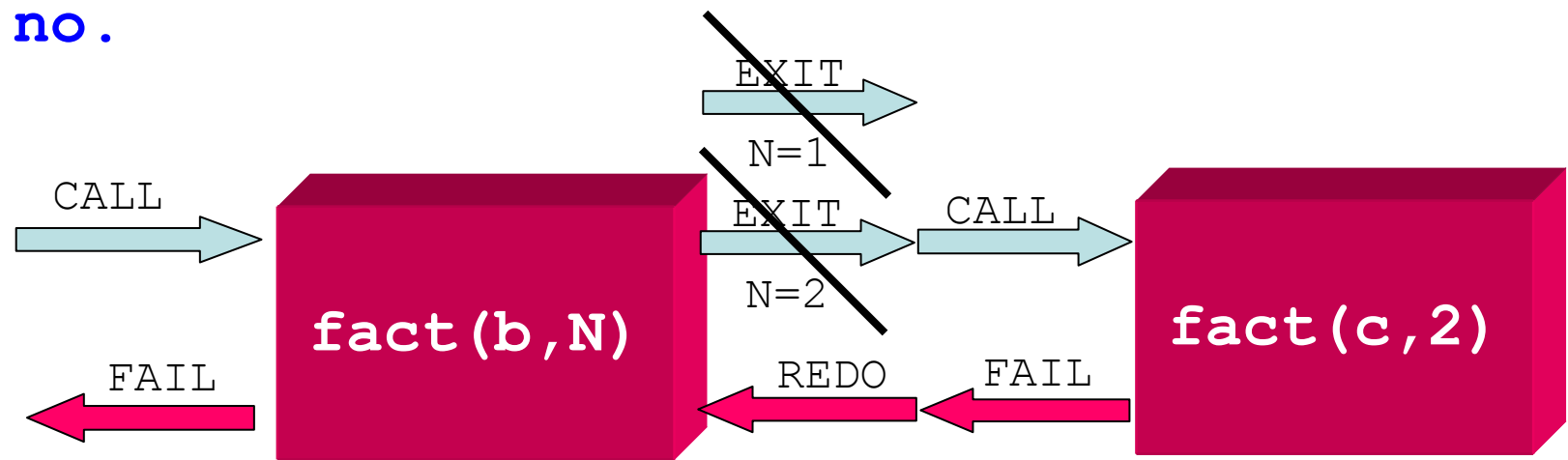
`fact(b,1).`

`fact(b,2).`

`a :- fact(b,N), fact(c,N).`

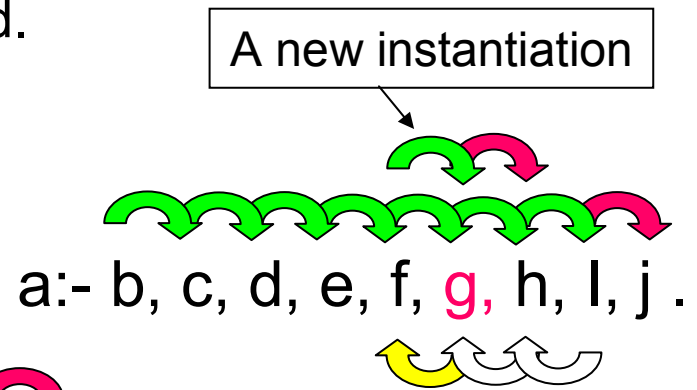
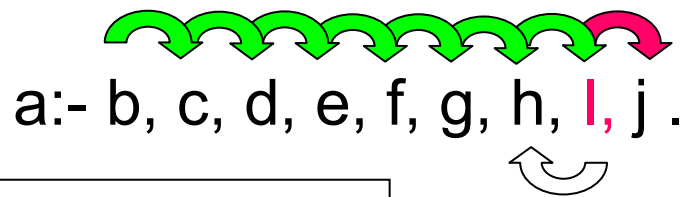
`|?- a.`

`no.`

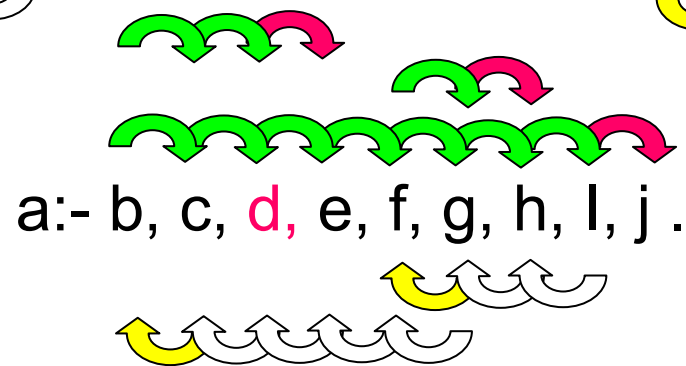


Prolog's Persistence

- When a sub-goal fails, Prolog will backtrack to the most recent successful goal and try to find another match.
- Once there are no more matches for this sub-goal it will backtrack again; retrying every sub-goal before failing the parent goal.
- A `call` can match any clause head.
- A `redo` ignores old matches.

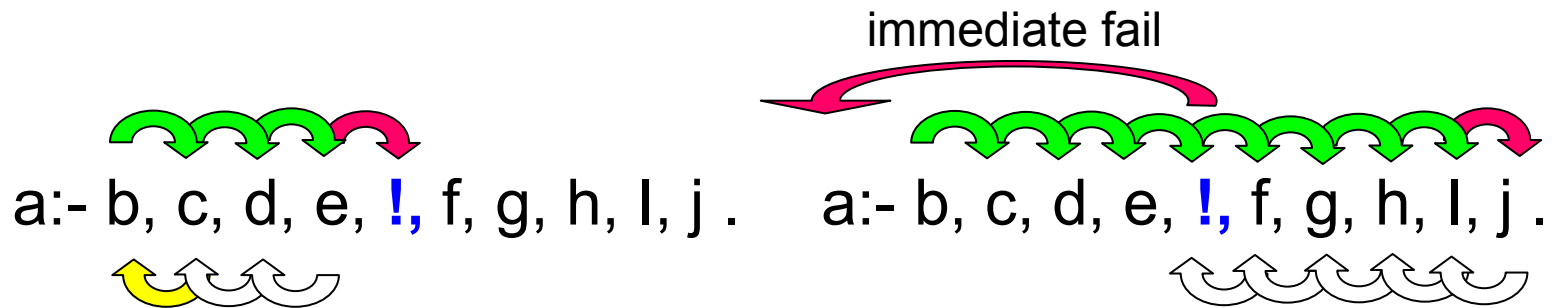


| | |
|--|-----------|
| | Succeed |
| | Fail |
| | Redo |
| | Backtrack |

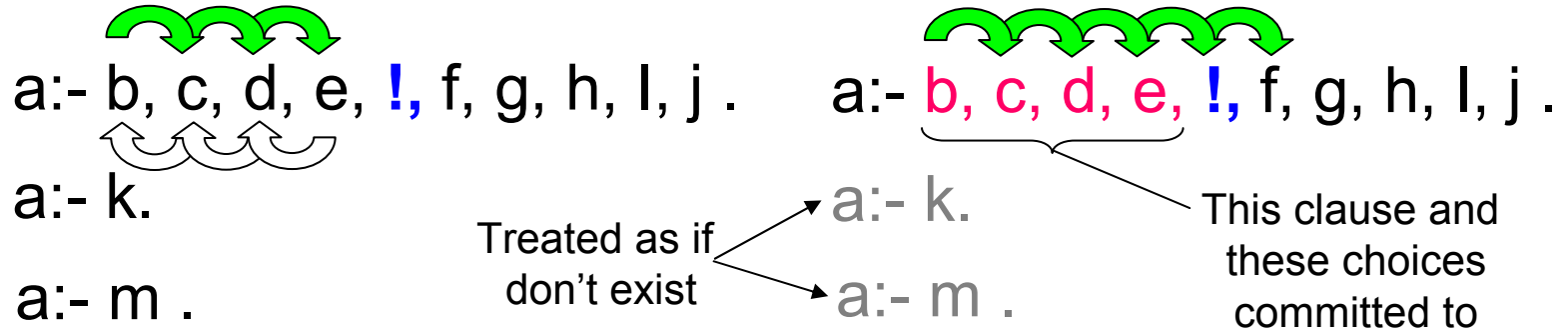


Cut !

- If we want to restrict backtracking we can control which sub-goals can be redone using the cut = !.
- We use it as a goal within the body of clause.
- It succeeds when called, but fails the parent goal (the goal that matched the head of the clause containing the cut) when an attempt is made to redo it on backtracking.
- It commits to the choices made so far in the predicate.
 - unlimited backtracking can occur before and after the cut but no backtracking can go through it.



Failing the parent goal



- The cut succeeds when it is called and commits the system to all choices made between the time the parent goal was invoked and the cut.
- This includes committing to the clause containing the cut.
 = the goal can only succeed if this clause succeeds.
- When an attempt is made to backtrack through the cut
 - the clause is immediately failed, and
 - no alternative clauses are tried.

Mutually Exclusive Clauses

- We should only use a cut if the clauses are mutually exclusive (if one succeeds the others won't).
- If the clauses are mutually exclusive then we don't want Prolog to try the other clauses when the first fails
= **redundant processing**.
- By including a cut in the body of a clause we are committing to that clause.
 - Placing a cut at the start of the body commits to the clause as soon as head unification succeeds.

$$a(1, X) :- !, b(X), c(X).$$
 - Placing a cut somewhere within the body (even at the end) states that we cannot commit to the clause until certain sub-goals have been satisfied.

$$a(_, X) :- b(X), c(X), !.$$

Mutually Exclusive Clauses (2)

```
f(X,0) :- X < 3.
f(X,1) :- 3 =< X, X < 6.
f(X,2) :- 6 =< X.
```

```
l?- trace, f(2,N).
1      1 Call: f(2,_487) ?
2      2 Call: 2<3 ?
2      2 Exit: 2<3 ? ?
1      1 Exit: f(2,0) ?
N = 0 ? ;
1      1 Redo: f(2,0) ?
3      2 Call: 3=<2 ?
3      2 Fail: 3=<2 ?
4      2 Call: 6=<2 ?
4      2 Fail: 6=<2 ?
1      1 Fail: f(2,_487) ?
no
```

Green Cuts !

```
f(x,0) :- x < 3, !.
f(x,1) :- 3 =< x, x < 6, !.
f(x,2) :- 6 =< x.
```

```
|- trace, f(2,N).
1      1 Call: f(2,_487) ?
2      2 Call: 2<3 ?
2      2 Exit: 2<3 ? ?
1      1 Exit: f(2,0) ?
N = 0 ? ;
no
```

If you reach this point don't bother trying any other clause.

- Notice that the answer is still the same, with or without the cut.
 - This is because the cut does not alter the logical behaviour of the program.
 - It only alters the procedural behaviour: specifying which goals get checked when.
- This is called a **green cut**. It is the correct usage of a cut.
- Be careful to ensure that your clauses are actually mutually exclusive when using green cuts!

Red Cuts !

```
f(x,0):- x < 3, !.
f(x,1):- 3 =< x, x < 6, !.
f(x,2):- 6 =< x.
```

Redundant?

```
| ?- f(7,N).
1      1 Call: f(7,_475) ?
2      2 Call: 7<3 ?
2      2 Fail: 7<3 ?
3      2 Call: 3=<7 ?
3      2 Exit: 3=<7 ?
4      2 Call: 7<6 ?
4      2 Fail: 7<6 ?
5      2 Call: 6=<7 ?
5      2 Exit: 6=<7 ?
1      1 Exit: f(7,2) ?
```

N = 2 ?

yes

- Because the clauses are mutually exclusive and ordered we know that once the clause above fails certain conditions must hold.
- We might want to make our code more efficient by removing superfluous tests.

Red Cuts !

```
f(X,0):- X < 3, !.
f(X,1):- X < 6, !.
f(X,2).
```

```
f(X,0):- X < 3.
f(X,1):- X < 6.
f(X,2).
```

```
| ?- f(7,N).
1      1 Call: f(7,_475) ?
2      2 Call: 7<3 ?
2      2 Fail: 7<3 ?
3      2 Call: 7<6 ?
3      2 Fail: 7<6 ?
1      1 Exit: f(7,2) ?
```

N = 2 ?
yes

```
| ?- f(1,Y).
1      1 Call: f(1,_475) ?
2      2 Call: 1<3 ?
2      2 Exit: 1<3 ? ?
1      1 Exit: f(1,0) ?
Y = 0 ? ;
1      1 Redo: f(1,0) ?
3      2 Call: 1<6 ?
3      2 Exit: 1<6 ? ?
1      1 Exit: f(1,1) ?
Y = 1 ? ;
1      1 Redo: f(1,1) ?
1      1 Exit: f(1,2) ?
Y = 2 ?
yes
```

Using the cut

- *Red cuts* change the logical behaviour of a predicate.
- TRY NOT TO USE RED CUTS!
- Red cuts make your code hard to read and are dependent on the specific ordering of clauses (which may change once you start writing to the database).
- If you want to improve the efficiency of a program use *green cuts* to control backtracking.
- Do not use cuts in place of tests.

To ensure a logic friendly cut either:

```
p(X) :- test1(X), !, call1(X).
p(X) :- test2(X), !, call2(X).
p(X) :- testN(X), !, callN(X).
```

```
p(1,X) :- !, call1(X).
p(2,X) :- !, call2(X).
p(3,X) :- !, callN(X).
```

testI predicates are mutually exclusive.

The mutually exclusive tests are in the head of the clause.

Cut - fail

- As well as specifying conditions under which a goal can succeed sometimes we also want to specify when it should fail.
- We can use the built-in predicate `fail` in combination with a cut to achieve this: “ `!, fail.` ”
 = if you reach this point, fail regardless of other clauses.
- e.g. If we want to represent the fact that ‘*Mary likes all animals except snakes*’.

```
likes(mary,X) :-
    snake(X), !, fail.

likes(mary,X) :-
    \+ snake(X),
    animal(X).
```

We need to combine a cut with the fail to stop the redundant call to the second clause on backtracking.

Cut – fail: why?

- However, using a cut-fail can make your code hard to follow.
- It is generally clearer and easier to define the conditions under which a fact is true rather than when it is false.

```
likes(mary,X) :-
    \+ snake(X),
    animal(X).
```

} This is sufficient to represent the fact.

- However, sometimes it can be much simpler to specify when something is false rather than true so cut-fail can make your code more efficient.
- As with all cuts; be careful how you use it.

Summary

- Clearing up equality: `=`, `is`, `==`, `=\=`, `==`, `\==`, `\+`
- REDO vs. CALL
- Controlling backtracking: the cut !
 - Efficiency: avoids needless REDO-ing which cannot succeed.
 - Simpler programs: conditions for choosing clauses can be simpler.
 - Robust predicates: definitions behave properly when forced to REDO.
- Green cut = cut doesn't change the predicate logic = **good**
- Red cut = without the cut the logic is different = **bad**
- Cut – fail: when it is easier to prove something is false than true.