Applied Databases

Lecture 10 Full-Text Search

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University of Edinburgh - February 16th, 2017

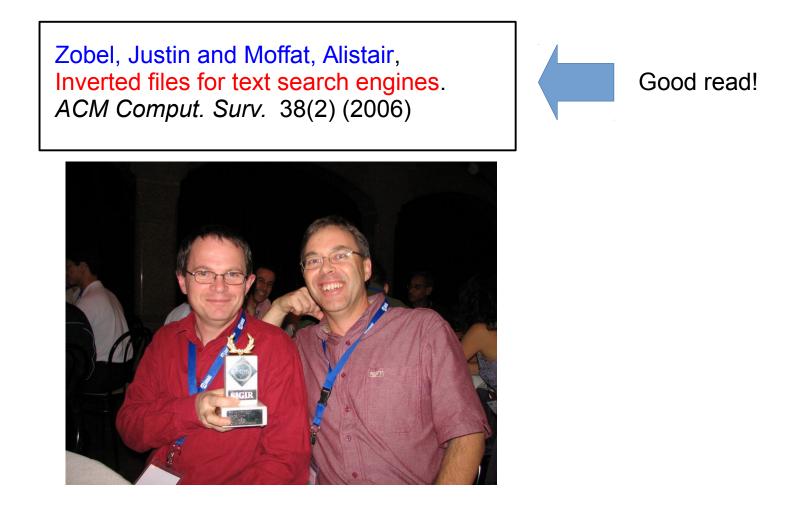
Outline

- 1. Text Search
- 2. Ranking & Similarity Measures
- 3. Inverted Files
- 4. Lucene (outlook)

Extra Reading Material

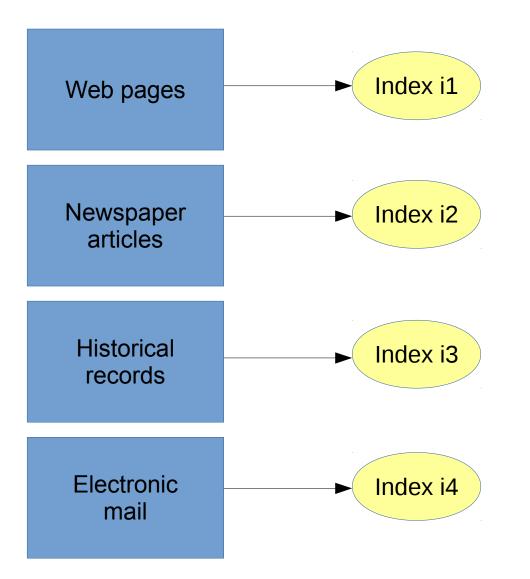
 \rightarrow Please check course web page.

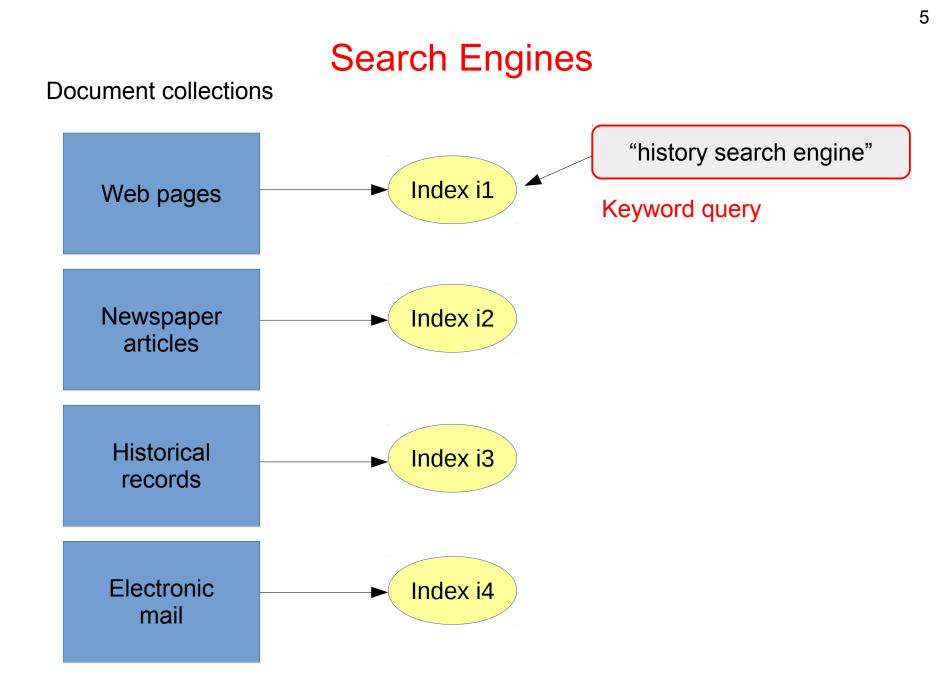
Most of this lecture based on this **article** (PDF linked on course web page)



Search Engines

Document collections





Search Engines Document collections "history search engine" Index i1 Web pages Keyword query Newspaper Index i2 query is executed over index articles engine returns a ranked lists of matching documents **Historical** Index i3 \rightarrow document 1 records \rightarrow document 2 \rightarrow document 3 \rightarrow Electronic Index i4 mail

RDBMS search (e.g. SQL)

- → DB system must answer arbitrarily complex queries
- → a match is a tuple that meets a specified logical condition

- → DB systems returns all matching tuples
- → each tuple has a unique access key; may search over that key

Text search

- → most queries are simple lists of terms or phrases
- → a match is a document that is appropriate to the query wrt statistical heuristics (it may not even contain all query terms!)
- → search engine returns fixed number of matches ranked by their statistical similarity
- → there may be millions of documents with non-zero similarity

Due to these "cultural" differences, the respective research communities of

- \rightarrow databases
- → information retrieval

have remained separate for many decades (and continue to do so)

For the same reason, we use *separate products* for combined search $(\rightarrow Assignments 1 \& 2 \text{ over ebay data})$:

- \rightarrow MySQL
- \rightarrow Apache Lucene

- \rightarrow databases
- \rightarrow information retrieval

Question for you

 \rightarrow what is the difference between "data" and "information"?

Challenges

- \rightarrow query term may occur in many documents
- \rightarrow each document may contain many terms

New

- \rightarrow representations for text indexes
- \rightarrow index construction techniques
- \rightarrow algorithms for evaluation of text queries

crucial for rapid response of major Web Search Engines (e.g. Google or Yahoo)

- \rightarrow compression and
- \rightarrow careful organization

reduction of

- index sizes
- time
- disk traffic during query evaluation

Search Engine = tool to find documents from a collection that are good matches to a user query

Collections are, e.g., web pages, news articles, emails, etc.

Collections vary dramatically in size

- → 10 years of research papers by a research (plain text) ca. 10 megabytes
- → 10 years of emails of the researcher ca. 100 megabytes
- → books in a small university library ca. 100 gigabytes
- → complete text of the web (year 2006) ca. 100 terabytes

(in 2014, Google has indexed 200TB, which is claimed to be only 0.4% of the Web)

Search Engine = tool to find documents from a collection that are good matches to a user query

Most text querying done

- \rightarrow by content
- \rightarrow satisfies an *information need*

A document matches an information need, if the user perceives it to be relevant.

- \rightarrow relevance is inexact!
- → a document may be relevant, but contain none of the query terms or irrelevant, even though it contains all the query terms.

A system is effective, if a good proportion of the first k search results are relevant.

\rightarrow bag-of-words queries

big old house

_		
-	1	The old night keeper keeps the keep in the town
6	2	In the big old house in the big old gown.
é	3	The house in the town had the big old keep
2	4	Where the old night keeper never did sleep.
ļ	5	The night keeper keeps the keep in the night
(6	And keeps in the dark and sleeps in the light.

Fig. 1. The Keeper database. It consists of six one-line documents.

- \rightarrow docs 2 and 3 contain all query terms
- \rightarrow docs 1 and 4 contain "old"
- \rightarrow only doc 2 contains the *phrase* "big old house"

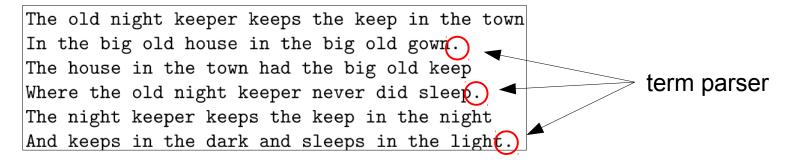
Parsing method for extracting terms from text:

- \rightarrow should HTML markup be indexed?
- \rightarrow or terms that appear within markup?
- \rightarrow hyphenated words, considered as one or two words?

More fundamentally:

- \rightarrow stemming? (= remove variant endings of a word)
- \rightarrow casefolding? (= convert to lowercase)
- \rightarrow stopping? (= remove common / functions words, e.g. "the")

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with casefolding (sorted vocabulary)

and big dark did gown had house in keep keeper keeps light never night old sleep sleeps the town where

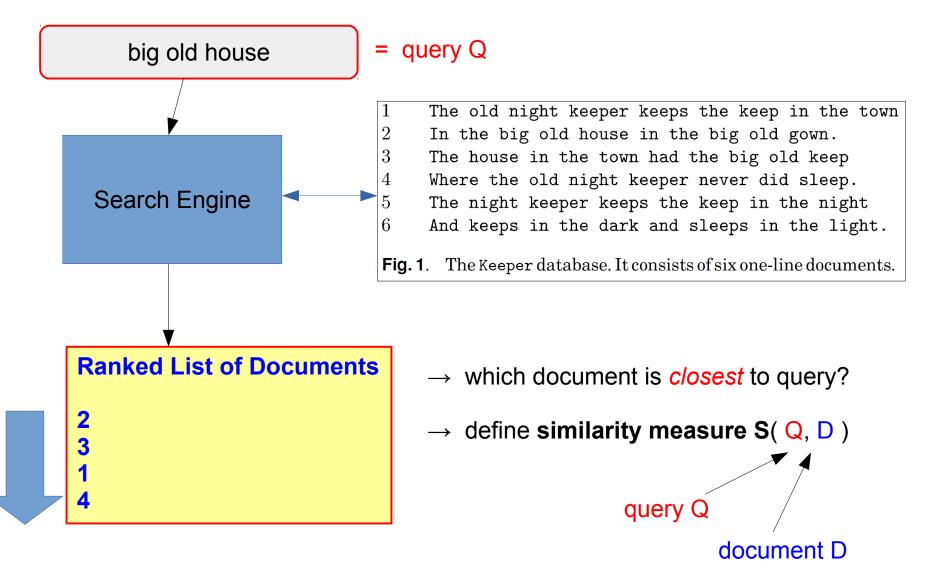
with stemming

and big dark did gown had house in keep light never night old sleep the town where

with stopping

big dark gown house keep light night old sleep town

2. Ranking and Similarity Measures



 \rightarrow how to define a good **similarity measure**?

(1) give higher score if many query terms appear in the document (many times)

 \rightarrow how to define a good **similarity measure**?

- (1) give higher score if many query terms appear in the document (many times)
- (2) give less weight to terms that appear in many documents
- (3) give more weight to terms that appear many times in a document
- (4) give less weight to documents that contain many terms.

Term Frequency (TF) f(D,T) = how many times does term T appear in document D? Document Frequence (DF)

f(T) = in how many documents of the collection does term T appear?

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```
Inverse Document Frequence (IDF)
1 / f(T)
```

```
TF * IDF = f(D,T) / f(T)
```

```
→ e.g. "old" appears in 4 documents (out of 6)

f(1, "old") = 1, thus TF*IDF = 1 / 4

f(2, "old") = 2, thus TF*IDF = 2 / 4

f(3, "old") = 1, thus TF*IDF = 1 / 4

f(4, "old") = 1, thus TF*IDF = 1 / 4
```

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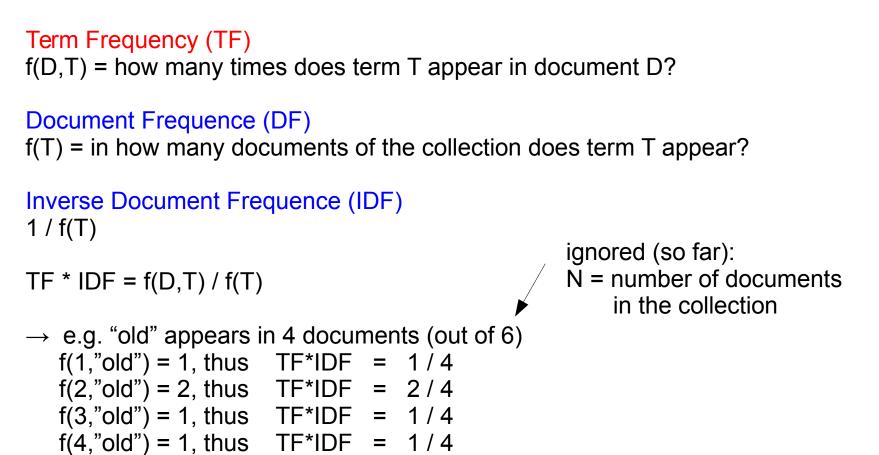
f(1, "old") = 1, thus TF*IDF = 1 / 4

f(2, "old") = 2, thus TF*IDF = 2 / 4

f(3, "old") = 1, thus TF*IDF = 1 / 4

f(4, "old") = 1, thus TF*IDF = 1 / 4
```

could be 300 appearances!



Term Frequency (TF) (non-scaled) scaled: $1 + \ln(f(D,T))$ f(D,T) = how many times does term T appear in document D?

```
Document Frequence (DF)

f(T) = in how many documents of the collection does term T appear?

Inverse Document Frequence (IDF) (non-scaled)

1 / f(T)

TF * IDF = f(D,T) / f(T)

\rightarrow e.g. "old" appears in 4 documents (out of 6)

<math>f(1,"old") = 1, thus TF*IDF = 1 / 4

f(2,"old") = 2, thus TF*IDF = 2 / 4
```

IDF = In (1 + N / DF) - "scaled"

```
Thus, TF*IDF for "old" and doc1: (1+ln(1))*ln(1+6/4) = 0.916
"old" and doc2: (1+ln(2))*ln(1+6/4) = 1.551
```

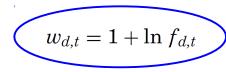
Example (patents)			Here: IDF	= log(N/DF
Term	TF(doc1)	TF(doc2)	TF(doc3)	DF	IDF
method the water bioreact	4,250 50,000 7,600 for 600	3,400 43,000 4,000 0	5,100 55,000 2,000 25	400 25	0.27 0.00 0.54 1.6 is not scaled
term	TF-IDF(do	c1) TF-IDF	(doc2) TF-I	DF(doc3)	
method the water bioreacto	43	148 0 104 960	918 0 2160 0	1377 0 1080 40	

inverse document frequency influences the TF-IDF value:

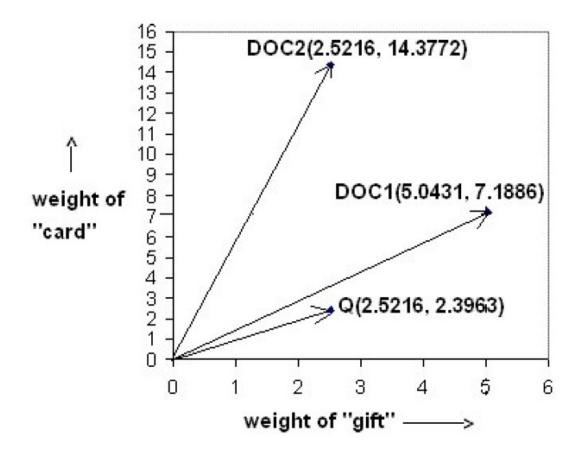
- → "method" occurs nearly as often as "water" in doc2 but TF-IDF value of "water" is more than double that of "method"
- \rightarrow a query "method bioreactor" would assign doc1 a score of 0.15 and doc3 a score of 0.04.

Given a query (T1, T2, .., Tk), compute for each document D the vector

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right)$$



<TFIDF(T1, D), ..., TFIDF(Tk, D)>

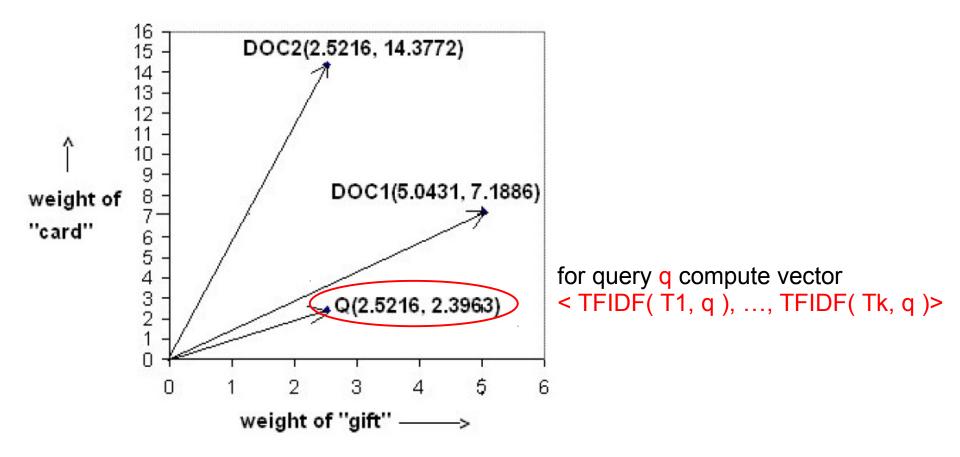


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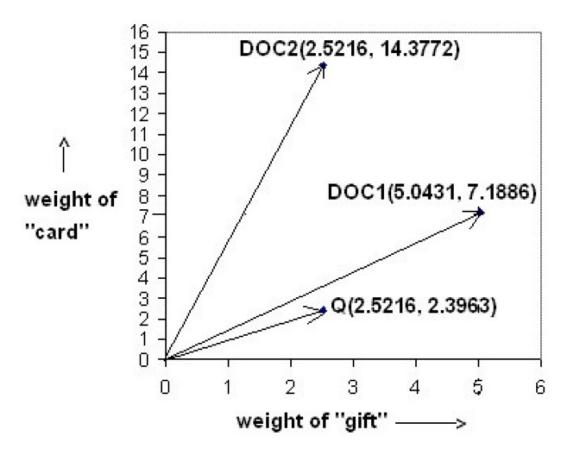
$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right)$$

 $w_{d,t} = 1 + \ln f_{d,t}$

<TFIDF(T1, D), ..., TFIDF(Tk, D) >



- → **angle** between DOC-k and Q determines similarity (length of vector not important)
- \rightarrow "relative closeness" of term weights

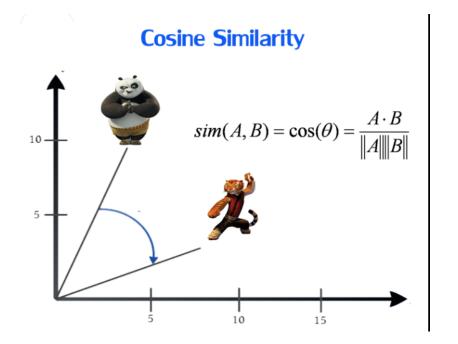


- → the closer angle is to zero, the more similar the documents
- → if angle is >=90 degrees, then documents have no words in common
- \rightarrow DOC1 and Q are very similar!

Given a query (T1, T2, .., Tk), compute for each document D the vector

<TFIDF(T1, D), ..., TFIDF(Tk, D) >

consider cosine similarity between such vector A and vector B for the query



cosine similarity:

cos(angle between A and B)

- → equals "1" if angle is zero (vectors have same direction)
- → equals "0" if orthogonal (90 degree) (means: no words in common)

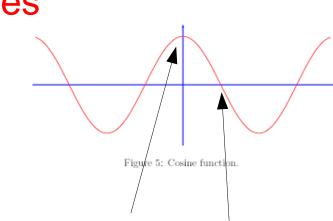
Figure 5: Cosine function.

Given a query (T1, T2, .., Tk), compute for each document D the vector

<TFIDF(T1, D), ..., TFIDF(Tk, D) >

consider cosine similarity between such vector A and vector B for the query

Similarity(A,B) =
$$\frac{A \cdot B}{\|A\| \|B\|} = \frac{\sum_{i=1}^{n} A_i B_i}{\sqrt{\sum_{i=1}^{n} A_i^2} \sqrt{\sum_{i=1}^{n} B_i^2}}$$



cosine similarity:

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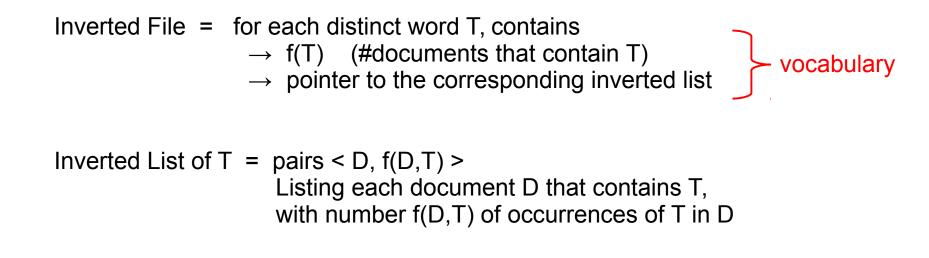
Index = datastructure that maps terms to documents containing them

E.g. consider pages of a book as "documents" Book index: maps words to pages

A concordance is an alphabetical list of the principal words used in a book or body of work, listing every instance of each word with its immediate context. Because of the time, difficulty, and expense involved in creating a concordance in the pre-computer era, only works of special importance, such as the Vedas [1] Bible, Qur'an or the works of Shakespeare or classical Latin had concordances prepared for them.

The first Bible concordance, for the Vulgate Bible, was compiled by Hugh of St Cher (d.**1262**), who employed 500 monks to assist him. In 1448 Rabbi Mordecai Nathan completed a concordance to the Hebrew Bible. It took him ten years.

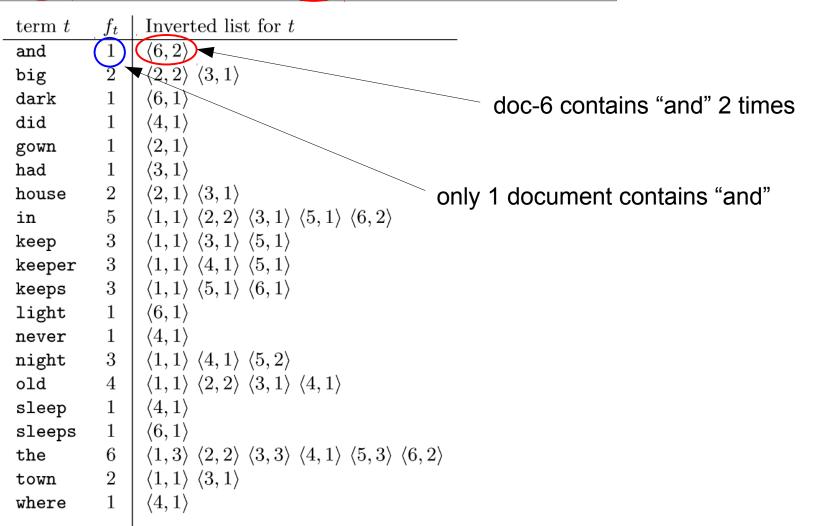




Thumb rule for effective retrieval:

 \rightarrow index **all terms**, even stop words, numbers, etc.

The old night keeper keeps the keep in the town
In the big old house in the big old gown.
The house in the town had the big old keep
Where the old night keeper never did sleep.
The night keeper keeps the keep in the night
And keeps in the dark and sleeps in the light.



1	The old night keeper keeps the keep in the town	
2	In the big old house in the big old gown.	
3	The house in the town had the big old keep	the second station of
4	Where the old night keeper never did sleep.	+ casefolding
5	The night keeper keeps the keep in the night	
6	And keeps in the dark and sleeps in the light.	

term t	f_t	Inverted list for t	
and	1	$\langle 6,2 \rangle$	—
big	2	$\langle 2,2 angle \; \langle 3,1 angle$	
dark	1	$\langle 6,1 angle$	$w(d,t) = 1 + \ln f(D,T)$
did	1	$\langle 4,1 angle$	
gown	1	$\langle 2,1 \rangle$	$\sum2$
had	1	$\langle 3,1 angle$	$W_d = \sum_{t} w_{d,t}^2$
house	2	$\langle 2,1 angle \ \langle 3,1 angle$	c t
in	5	$\langle 1,1 \rangle \langle 2,2 \rangle \langle 3,1 \rangle \langle 5,1 \rangle \langle 6,2 \rangle$	
keep	3	$\langle 1,1 angle \; \langle 3,1 angle \; \langle 5,1 angle$	
keeper	3	$\langle 1,1\rangle \langle 4,1\rangle \langle 5,1\rangle$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
keeps	3	$\langle 1,1 angle\;\langle 5,1 angle\;\langle 6,1 angle$	$W_d \mid 11.4 13.5 11.4 8.0 11.3 12.6$
light	1	$\langle 6,1 angle$	
never	1	$\langle 4,1 angle$	
\mathtt{night}	3	$\langle 1,1 angle \; \langle 4,1 angle \; \langle 5,2 angle$	\rightarrow how do doc-2 and doc-4 differ?
old	4	$\langle 1,1 angle \; \langle 2,2 angle \; \langle 3,1 angle \; \langle 4,1 angle$	
sleep	1	$\langle 4,1 angle$	
sleeps	1	$\langle 6,1 \rangle$	\rightarrow doc-4 is more "specific"
the	6	$\left \begin{array}{c} \langle 1,3 \rangle \ \langle 2,2 \rangle \ \langle 3,3 \rangle \ \langle 4,1 \rangle \ \langle 5,3 \rangle \ \langle 6,2 \rangle \end{array} \right $	
town	2	$\langle 1,1 angle \; \langle 3,1 angle$	
where	1	$\langle 4,1 angle$	

1 The old night keeper keeps the keep in the town
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where $1 \mid \langle 4, 1 \rangle$

term t	f_t	Inverted list for t	
and	1	$\langle 6,2 \rangle$	—
big	2	(2,2)(3,1)	
dark	1	$\langle \overline{6,1} \rangle$	$w(d,t) = 1 + \ln f(D,T)$
did	1	$\langle 4,1 \rangle$	
gown	1	(2,1)	$\sum u^2$
had	1	$\langle \overline{3,1} \rangle$	$W_d = \sum_{t} w_{d,t}^2$
house	2	$\langle 2,1 \rangle \langle 3,1 \rangle$	$z _ t$
in	5	$\langle 1, 1 \rangle \langle 2, 2 \rangle \langle 3, 1 \rangle \langle 5, 1 \rangle \langle 6, 2 \rangle$	$d \mid 1 \mid 2 \mid 3 \mid 4 \mid 5 \mid 6$
keep	3	$\langle 1,1 \rangle \ \overline{\langle 3,1 \rangle} \ \overline{\langle 5,1 \rangle}$ -	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
keeper	3	$\langle 1,1 \rangle \langle 4,1 \rangle \langle 5,1 \rangle$	$W_d \mid 11.4 \mid 13.0 \mid 11.4 \mid 0.0 \mid 11.3 \mid 12.0$
keeps	3	$\langle 1,1 \rangle \ \langle 5,1 \rangle \ \langle 6,1 angle$	
light	1	$\langle 6,1 angle$	$4*(1 \pm \ln 2)$ (A) $\pm 2 = 12$ (666)
never	1	$\langle 4,1 angle$	$4^{*}(1 + \ln 2)^{2} + 2 = 13.4666$
\mathtt{night}	3	$\langle 1,1 \rangle \langle 4,1 \rangle \langle 5,2 \rangle$	
old	4	$\langle 1, 1 (2, 2) \langle 3, 1 \rangle \langle 4, 1 \rangle$	
sleep	1	$\langle 4,1 angle$	
sleeps	1	$\langle 6,1 \rangle$	
the	6	$ \langle 1, 3 (2,2) \langle 3, 3 \rangle \langle 4, 1 \rangle \langle 5, 3 \rangle \langle 6, 2 \rangle $	
town	2	$\langle 1,1 \rangle \overline{\langle 3,1 \rangle}$	

1 The old night keeper keeps the keep in the town
2 In the big old house in the big old gown.
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term t	f_t	Inverted list for t	
and	1	$\langle 6,2 \rangle$	—
big	2	$(2,2)\langle 3,1\rangle$	
dark	1	$\langle 6,1 \rangle$	$w(d,t) = 1 + \ln f(D,T)$
did	1	(4,1)	
gown	1	(2,1)	M d - $\sum u^2$
had	1	$\langle 3,1 \rangle$	W_d = $\sum w_{d,t}^2$
house	2	(2,1)(3,1)	t
in	5	$\langle 1, 1 \rangle \langle 2, 2 \rangle \langle 3, 1 \rangle \langle 5, 1 \rangle \langle 6, 2 \rangle$	$d \mid 1 2 3 4 5 6$
keep	3	$\langle 1,1\rangle$ $\langle \overline{3,1}\rangle$ $\langle 5,1\rangle$	$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
keeper	3	$\langle 1,1\rangle\langle 4,1\rangle\rangle\langle 5,1\rangle$	$W_d \mid 11.4 \mid 15.5 \mid 11.4 \mid 8.0 \mid 11.5 \mid 12.0$
keeps	3	$\langle 1,1 \rangle \ \overline{\langle 5,1 \rangle} \ \overline{\langle 6,1 \rangle}$	
light	1	$\langle 6,1 \rangle$	$4*(1 \pm \ln 2)$ (0.0 $\pm 2 - 12$ (0.0 ± 2)
never	1	(4,1)	$4^{*}(1 + \ln 2)^{2} + 2 = 13.4666$
\mathtt{night}	3	$\langle 1, 1 \rangle \langle 4, 1 \rangle \langle 5, 2 \rangle$	
old	4	$\langle 1, 1, (2, 2) \rangle \langle 3, 1 \rangle \langle 4, 1 \rangle$	8*(1+ ln 1)^2 = 8
sleep	1	(4,1)	
sleeps	1	$\langle \overline{6,1} \rangle$	
the	6	$\left \begin{array}{c} \langle 1, 3 \\ \langle 2, 2 \\ \rangle \\ \langle 3, 3 \\ \langle 4, 1 \\ \rangle \\ \langle 5, 3 \\ \rangle \\ \langle 6, 2 \\ \rangle \\$	
town	2	$\langle 1,1 angle \ \overline{\langle 3,1 angle}$	
where	1	$\overline{\langle 4,1\rangle}$	

The old night keeper keeps the keep in the town
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$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t}$$

$$W_d = \sum_t w_{d,t}^2$$

$$S_{q,d} = \frac{\sum_t w_{d,t} \cdot w_{q,t}}{W_d}$$

$$\frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6 | 6}$$

query score S(q,d) for document d on query q, from [Zobel, Moffat 2006]

(incorporates cosine-simularity and TF*IDFT)

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$
$$W_d = \sum_t w_{d,t}^2$$
$$\mathbf{q} = \mathbf{old}$$
$$\mathbf{q} = \mathbf{lod}$$

old
$$(4)$$
 $\langle 1,1 \rangle$ $\langle 2,2 \rangle$ $\langle 3,1 \rangle$ $\langle 4,1 \rangle$

inverted file entry for "old"

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$W_d = \sum_{t} w_{d,t}^2 \qquad q = \boxed{\operatorname{old}}$$

$$q = \boxed{\operatorname{old}}$$

S(q,doc-3) = w(doc-3, "old") * w(q, "old") / 11.4 = (1 + ln(1)) * ln(1 + 6/4) / 11.4 = 0.0804

old (4) $\langle 1,1 \rangle \langle 2,2 \rangle \langle 3(1) \langle 4,1 \rangle$

inverted file entry for "old"

$$S(q,doc-3) = w(doc-3, "old") * w(q, "old") / 11.4 = (1 + ln(1)) * ln(1 + 6/4) / 11.4 = 0.0804$$

$$S(q,doc-4) = w(doc-4, "old") * w(q, "old") / 8 = (1 + ln(1)) * ln(1 + 6/4) / 8 = 0.1145$$

old (4) $\langle 1,1 \rangle \langle 2,2 \rangle \langle 3,1 \rangle \langle 4(1) \rangle$

S(q,doc-3) = w(doc-3, "old") * w(q, "old") / 11.4 = (1 + ln(1)) * ln(1 + 6/4) / 11.4 = 0.0804S(q,doc-4) = w(doc-4, "old") * w(q, "old") / 8 = (1 + ln(1)) * ln(1 + 6/4) / 8 = 0.1145

 \rightarrow doc-4 has higher score because of lower W_d value! (it is more 'specific')

- → want to compute score of S(q,d) of query q on document 2
- \rightarrow need to compute:

- → want to compute score of S(q,d) of query q on document 2
- \rightarrow need to compute:

All we need are the **inverted file** entries for "big", "old", and "house"!

w(q,"big") = ln(1 + 6/2) = ln(4)

 $\overline{\langle 2,2\rangle}\ \overline{\langle 3,1\rangle}$

 \bigcirc

big

w(doc-2, "big") = 1 + ln(f(doc-2, "big")) = 1+ ln(2)

old $4 \mid \langle 1,1 \rangle \mid \langle 2,2 \rangle \mid \langle 3,1 \rangle \mid \langle 4,1 \rangle$

 $\langle 2,1 \rangle \ \langle 3,1 \rangle$

house

2

1 The old night keeper keeps the keep in the town
2 In the big old house in the big old gown.
3 The house in the town had the big old keep
4 Where the old night keeper never did sleep.
5 The night keeper keeps the keep in the night
6 And keeps in the dark and sleeps in the light.

 $w(doc-2, "big") = 1 + ln(f(doc-2, "big")) = 1 + ln(\underline{2})$ S(q,doc-2) = w(doc-2, "big") * w(q, "big") / W(doc-2) + .. = (1 + ln(2)) * ln(4) / 13.5 = 2.3472 / 13.5 + ..

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$W_d = \sum_t w_{d,t}^2$$

$$S_{q,d} = \frac{\sum_t w_{d,t} \cdot w_{q,t}}{W_d}$$

$$q = big old house$$

$$w(q, "old") = \ln(1 + 6 / 4) = \ln(2.5)$$

w(doc-2, "old") = 1 + ln(f(doc-2, "old")) = 1 + ln(2)

S(q,doc-2) = (2.3472 + ln(2.5) * (1 + ln(2))) / 13.5 = (2.3472 + 1.5514) / 13.5

 $\langle 2,1\rangle$ $\langle 3,1\rangle$

big2 $\langle 2,2\rangle$ $\langle 3,1\rangle$ old $\langle 4\rangle$ $\langle 1,1\rangle$ $\langle 2,2\rangle$ $\langle 3,1\rangle$ $\langle 4,1\rangle$ house2

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$W_d = \sum_t w_{d,t}^2$$

$$W_{d,t} = \frac{\sum_t w_{d,t} \cdot w_{q,t}}{W_d | 1.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$q = big old house$$

w(q, "house") = ln(1 + 6/2) = ln(4)

w(doc-2, "house") = 1 + ln(f(doc-2, "house")) = 1 + ln(1) = 1

S(q,doc-2) = (3.8986 + ln(4) * 1) / 13.5 = (3.8986 + 1.3863) / 13.5 = 0.3915

 $\overline{\langle 3, 1 \rangle}$

 $\langle 2(1) \rangle$

big $2 \mid \langle 2,2 \rangle \langle 3,1 \rangle \mid$ old $4 \mid \langle 1,1 \rangle \langle 2,2 \rangle \langle 3,1 \rangle \langle 4,1 \rangle \mid$ house (2)

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$W_d = \sum_t w_{d,t}^2$$

$$g = big old house$$

$$w(q,"big") = \ln(1 + 6 / 2) = \ln(4)$$

$$w(doc-3, "big") = 1 + \ln(f(doc-3, "big")) = 1 + \ln(1) = 1$$

1.3863 / 11.4

 $\fbox{big} \qquad \fbox{2} \quad \langle 2,2 \rangle \ \langle 3, \fbox{1} \qquad \texttt{old} \qquad 4 \quad \langle 1,1 \rangle \ \langle 2,2 \rangle \ \langle 3,1 \rangle \ \langle 4,1 \rangle \qquad \texttt{house} \qquad 2 \quad \langle 2,1 \rangle \ \langle 3,1 \rangle \\ \end{cases}$

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d}{1} \quad \frac{1}{2} \quad \frac{3}{3} \quad \frac{4}{5} \quad \frac{5}{6} \\ W_d \quad 11.4 \quad 13.5 \quad 11.4 \quad 8.0 \quad 11.3 \quad 12.6 \\ W_d \quad 11.4 \quad 13.5 \quad 11.4 \quad 8.0 \quad 11.3 \quad 12.6 \\ Q = \qquad big \text{ old house} \\ W(q, \text{"old"}) = \ln(1 + 6 / \underline{4}) = \ln(2.5)$$

w(doc-3, "old") = 1 + ln(f(doc-3, "old")) = 1 + ln(1) = 1

S(q,doc-3) = (1.3862 + ln(2.5) + ...) / 11.4 = (1.3863 + 0.9163) / 11.4

$$w_{q,t} = \ln\left(1 + \frac{N}{f_t}\right) \qquad w_{d,t} = 1 + \ln f_{d,t} \qquad \frac{d | 1 | 2 | 3 | 4 | 5 | 6}{W_d | 11.4 | 13.5 | 11.4 | 8.0 | 11.3 | 12.6}$$

$$W_d = \sum_t w_{d,t}^2$$

$$Q = \qquad big old house$$

$$q = big old house$$

w(q, "old") = ln(1 + 6 / 2) = ln(4)

w(doc-3, "old") = 1 + ln(f(doc-3, "old")) = 1 + ln(1) = 1

S(q,doc-3) = (1.3863 + 0.9163 + ln(4)) / 11.4 = 0.3236

big2 $\langle 2,2\rangle$ $\langle 3,1\rangle$ old4 $\langle 1,1\rangle$ $\langle 2,2\rangle$ $\langle 3,1\rangle$ $\langle 4,1\rangle$ houseQ

 $\langle 2,1\rangle$ $\langle 3,1\rangle$

1 The old night keeper keeps the keep in the town
2 In the big old house in the big old gown.
3 The house in the town had the big old keep
4 Where the old night keeper never did sleep.
5 The night keeper keeps the keep in the night
6 And keeps in the dark and sleeps in the light.

3. Inverted Indexes / Files

- \rightarrow inverted lists are stored contiguously
- → vocabulary stored in simple extensible structure (e.g., B-tree) (may be preprocessed by stemming and stopping)
- → inverted lists consist of doc numbers with #occurrences (possibly augmented by word positions)
- → ranking involves a set of accumulators and term-by-term processing of inverted lists

4. Lucene (outlook)

Lucene allows you to take care of everything mentioned today:

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- \rightarrow you can choose different Analyzers to do
 - casefolding
 - stemming (wrt a given language)
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- → you can insert documents into a collection and let Lucene generate inverted files for you (= "indexing" – very efficient!)

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 - stemming (wrt a given language)
 - stopping (wrt a given language)
- → you can insert documents into a collection and let Lucene generate inverted files for you (= "indexing" very efficient!)
- → you can then (very efficiently) retrieve the k top-most relevant documents in your collection!
- \rightarrow ranking function is a bit more sophisticated

 $score(q,d) = \sum [tf(t_d) \times idf(t) \times boost(t.field_d) \times lengthNorm(t.field_d)] \times coord(q,d) \times qNorm(q)$

1) sizes of inverted files?

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	NewsWire	Web
Size (gigabytes)	1	100
Documents	400,000	12,000,000
Word occurrences (without markup)	180,000,000	11,000,000,000
Distinct words (after stemming),	400,000	16,000,000
per document, totaled	70,000,000	3,500,000,000

Size of Inverted Index for **NewsWire (1 GB):** 435 MB

- **12MB** for 400,000 words, pointers, and counts
- **1.6MB** for 400,000 W(D)-values
- **280MB** for 70,000,000 document identifiers (four bytes each)
- **140MB** for 70,000,000 document frequencies (two bytes each)

- 1) sizes of inverted files?
- 2) limits of inverted files
 - \rightarrow imagine **substring search** (e.g. in DNA strands)
 - \rightarrow number of substrings is quadratic, cannot possibly generate/store them!

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 - \rightarrow imagine **substring search** (e.g. in **DNA** strands)
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SOLUTION:

- 3) in-memory indexes (for substring search, like DNA)
 - \rightarrow occupy a fraction of 435MB (40% of NewsWire)
 - \rightarrow run much faster :-)

Google's speed = (in-memory + MANY machines)

- 1) sizes of inverted files?
- 2) limits of inverted files
 - \rightarrow imagine **substring search** (e.g. in DNA strands)
 - → number of substrings is quadratic, cannot possibly generate/store them!
- 3) in-memory indexes for substring search

4) online substring search (without indexes)

END Lecture 10